

# The FAMOOS Object-Oriented Reengineering Handbook

http://www.iam.unibe.ch/~famoos/handbook/

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Version: October 15, 1999 (As Released to the general Public)

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This work has been funded by the European Union under the ESPRIT program Project no. 21975 (FAMOOS) as well as by the Swiss Government under Project no. NFS-2000-46947.96 and BBW-96.0015.

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## Preface

### How to Read this Book

The book is organized into five parts:

**I. Techniques.** This first part describes various techniques that help during the reengineering life cycle. Source code metrics automatically measure properties of the software product and help in focusing the reengineering task by pin-pointing key classes of the system. Then the annotation of basic graphs with metrics reveals semantic structures that are hidden in plain source code and are generally helpful when presenting large amounts of data. Grouping is a way of building more abstract views that are based on the elementary, often overly detailed views from the source code in order to reveal higher-level problems. Finally, an approach to reorganisation is presented based on refactorings and design patterns.

#### **II. Reverse Engineering**

- and III. Reengineering. The second and third parts form the core of the book. They consists of so-called *reengineering patterns*. Reengineering patterns capture tacit knowledge about when and how to apply reverse and reengineering tools and techniques as well as their implications. A reengineering pattern is somewhat like a design pattern. However, while a design pattern presents a solution to a design problem, a reengineering pattern relates two solutions (an existing solution and a target solution) via a process which transforms the one into the other. The reengineering patterns in the book tackle well known reverse and reengineering techniques often encountered in object-oriented programming.
- **IV. Tools.** The third part contains a description of some of the tool prototypes that have been developed in the context of the Famoos project. Without tools, reengineering is an almost impossible task because of the huge amount of information comprising most legacy systems. To be able to exchange reengineering data between different tools, an information exchange model called FAMIX is proposed.
- **Background.** The last part contains an introduction into software metrics and some discussions about the use of metrics when dealing with object oriented software.

## **Annotated Bibliography**

Several good books exist today on how to improve the development of applications at several levels. We invite the reader to read these books in order to have a better overview of the field. This sections contains an annotated bibliography of material which is relevant to OO Reengineering. We did not aim for completeness, but rather selected information sources we have found interesting. Omission of any work does not imply that the work is less significant than those annotated here.

#### **Software Engineering In General**

- Both [Som96] and [Pre94] provide a broad overview of software engineering. Their books cover issues like reengineering and reverse engineering, CASE tools and metrics.
- [Dav95] provides lots of good practical advice on how to tackle software projects, some of which motivates work on reengineering.

#### **Object-Oriented Engineering**

• [GR95] provides a decision framework for managing OO projects. Rather then imposing a particular software process or method, it tells you how you can built your own.

#### **Conferences, Journals and Special Issues**

- [CT97] is a special issue on object-oriented reengineering.
- [Arn92] is a book collecting various early papers on reengineering.
- [WC94], [WN96] are more recent special issues on reverse and reengineering.
- Since 1994, there is a yearly conference on reverse engineering. It is called WCRE (*Working Conference on Reengineering*). The proceedings from 1995 onwards are published by IEEE Computer Society Press.
- ICSM (*International Conference on Software Maintenance*) and EuroMICRO (*Software Maintenance and Reengineering*) are other conferences focussing more on reengineering and maintenance. Their proceedings are also published by IEEE.
- The *Journal of Software Maintenance Research and Practicve* is a journal dedicated to software maintenance and published bi-mothly by Wiley and Sons.

#### **Metrics**

- [FP97] is the seminal work on metrics but does cover very little on object-oriented metrics.
- [HS96a] provides an overview of the state of the art in object- oriented metrics.
- [LK94] is a pragmatic handbook on how to use metrics to check object-oriented source code.

#### **Object-Oriented Design**

- [Rie96] presents a list of object-oriented design heuristics using C++.
- [Mey97] elaborates on Design by Contract.
- Design patterns are discussed in many books, most notably in [GHJV95] and in proceedings of the different PLoP conferences.
- [Lak96] describes issues in building large scale (C++) systems including design considerations such as layering and more practical issues such as finding efficient include structures.

#### **Information Exchange (Meta-Meta Models)**

- CDIF (CASE data interchange format). See http://www.eigroup.org/
- MOF (Meta-Object Facility) and XMI. See http://www.omg.org/

#### **Idioms**

The books that follow contain practical information on exploiting programming language features to write good code.

- [Bec97] contains a set of idioms related to the SMALLTALK language. The main focus of the book is to show how to write code that communicate its intent.
- [Mey98, Mey96] Meyer on the one hand focusses very much on the specific issues of C<sup>++</sup> and explains complicated concepts such as the const mechanism in detail, or how to replace the default memory manager. On the other hand, Meyer also explains general OO concepts, like multiple inheritance, and its pitfalls in C<sup>++</sup>.
- [Cop92] Coplien strives to teach C<sup>++</sup> *fluency* by well known idioms like the orthodox canonical class form. He shows examples of how C<sup>++</sup> can be used in a functional style. Some of the desing idioms presented in this book have been later rewritten into a pattern language.

#### **UML**, Object-Oriented Documentation

- [Fow97b] provides a fast introduction to UML including the notion of "perspectives" which is quite interesting from a reverse engineering point of view because it is a way to specify how a certain UML diagram should be interpreted (i.e., on a Conceptual, Specification or Implementation level).
- [BRJ98], [RJB99] provide a good user reference and language reference for UML.
- [Joh92],[OQC97] present how patterns can support the documentation of a frameworks.
- [Bro96], [Wuy98], [PK98] present some possible approaches to support design patterns extraction.
- [FMvW97] shows how design patterns can be supported at the development environment level.
- [SLMD96] presents Reuse Contracts as a way to document frameworks for evolution.
- [WCH87] presents some discussion about variety of composition relationships.

#### **Refactoring and Code Smells**

- [Fow99] summarises practical experience with refactorings and code smells.
- The Ph.D. work of Opdyke [Opd92] on Refactoring resulted in a number of papers describing incremental redesign performed by humans supported by refactoring tools [OJ93], [JO93].
- [RBJ97b] describes the Refactoring Browser, a SMALLTALK tool that represents the state of the art in the field is described in and can be obtained from http://st-www.cs.uiuc.edu/.
- Both Casais [Cas91], [Cas92], [Cas94], [Cas95a] and Moore ([Moo96]) report on tools that optimise class hierarchies without human intervention. Schulz et al. illustrate the feasibility of refactorings on a subset of C++ [SGMZ98b].
- There exists a web-page discussing "code smells", i.e. suspicious symptoms in source code that might provide targets for refactoring. See http://c2.com/cgi/ wiki?CodeSmells

#### **Reverse and Reengineering Taxonomy**

• [CC190] (reappeared in [CC92]) provides a reverse and reengineering taxonomy. Unfortunately, it does not cover OO specific issues like refactoring. http://www.tcse.org/ revengr/ taxonomy.html

#### **Organisations**

- IEEE Computer Society's Technical Committee on Reverse Engineering. See <a href="http://www.tcse.org/revengr">http://www.tcse.org/revengr</a>
- The Reengineering Forum (an industry association). See http://www.reengineer.org/

## **Chapter 1**

# The Need for Object-Oriented Reengineering

Reengineering legacy systems has become a vital matter in today's software industry. In the past few years, most of the reengineering efforts were focussed on systems written in traditional programming languages such as COBOL, Fortran and C. But recently an increasing demand for reengineering object-based systems has emerged. This recent evolution is not caused by failure of the object-oriented paradigm. Rather, it illustrates that the mere application of object-oriented techniques is not sufficient to deliver flexible and adaptable systems. This is due to a number of obvious problems:

- *Lack of experience*. It requires several years of experience to fully exploit the potential of the objectoriented paradigm. Such experience is often built up during the initial stages of a project, at the time when the most crucial parts of the system are implemented.
- *Hybrid programming languages*. The use of hybrid languages –like C++ and Ada–, combined with a "learn on the job" approach, prevents programmers from making the necessary paradigm shift.
- *Technology expansion*. Legacy systems could not benefit from emerging standards (e.g., UML, CORBA), technological advancements (e.g., design patterns, architectural styles) and extra language features (e.g., C++ templates, Ada inheritance).

These problems are accidental in nature: given proper training and sufficient tool support they will eventually be resolved. So why should one worry about object-oriented reengineering, since within a few years there won't be any more object-oriented legacy systems? In fact there a more fundamental problem.

The law of *software entropy* dictates that even when a system starts off in a well-designed state, requirements evolve and customers demand new functionality, frequently in ways the original design did not anticipate. A complete redesign may not be practical, and a system is bound to gradually lose its original clean structure and deform into a bowl of "object-oriented spaghetti" [WH92], [Cas98], [BMMM98].

Many of the early adopters of the object-oriented paradigm have experienced such software entropy effects. Their systems are developed using object-oriented design methods and languages of the late 80s and exhibit a range of problems that prohibits them meeting the evolving requirements imposed by their customers. Their systems have become overly rigid thus compromising thier competitive advantage and as a consequence object-oriented reengineering technology has become vital to their business.

## 1.1 The FAMOOS Project

The need for object-oriented reengineering technology has been recognised by two of the leading European companies, namely Daimler-Benz and Nokia. Together with the University of Berne, Forshungszentrum

Informatik, SEMA Spain and Take5 they started a research project –named FAMOOS  $^{1}$  – to investigate tools and techniques for dealing with object-oriented legacy systems.

The handbook you are reading right now is one of the main results of the FAMOOS project. It collects techniques and knowledge on the problem of software evolution with a special emphasis on object-oriented software. Most of the subject matter is not "new" in the sense that it represents new discoveries. Rather the handbook regroups much of the knowledge about redesign, metrics and heuristics into a single work that is focused on object-oriented reengineering.

#### 1.1.1 Case Studies

All the techniques described in the handbook have been verified on six industrial case-studies, ranging from 50.000 lines of C++ up until 2,5 million lines of Ada. Figure 1.1 provides a quick overview of all of the FAMOOS case studies.

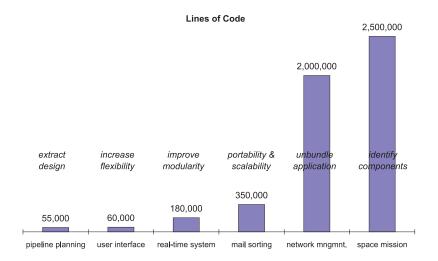


Figure 1.1: Overview of the FAMOOS Case Studies

- *Pipeline Planning.* The system supports the planning of liquid flow in a pipeline between multiple stations. The reengineering goal was to extract design from source-code, in order to reduce the cost of implementing similar systems, probably in other languages. The system is written in C++ and is a candidate for being rewritten in Java or Smalltalk.
- *User Interface.* This software provides graphical representations of telecommunication networks to telecom operators. The reengineering goal was to increase the flexibility of the software, i.e. improve its portability, facilitate addition of functionality and enhance tailorability towards customers. The system is written in C++.
- *Real-time System*. This software provide operating system features for embedded real-time controlling of hardware. The reengineering goal was to improve modularity for gaining shorter editcompile-run cycles. The system is written in a mixture of C and C++.
- *Mail Sorting*. A control system for machines sorting mail envelopes. The software is highly configurable, to deal with the different ways countries over the world handle letters. The software itself is based on an internally developed distributed architecture which hindered the future evolution. The

<sup>&</sup>lt;sup>1</sup>If you want to read more about the FAMOOS project and its results, we suggest to browse the web-sites offered by the respective project partners: http://dis.sema.es/ projects/ famoos/; http://www.iam.unibe.ch/ famoos/; http://www.fzi.de/ prost/

reengineering goal was to investigate how new technology could improve the portability and scalability (e.g. CORBA, Java, HTML). The system is written in a mixture of C and C++.

- *Cellular Network Management.* This case-study concerned a management system for digital networks. The main goal of the reengineering project was to unbundle the application, i.e. split the system into sub-products that can be developed and sold separately. The system is written in a mixture of C and C++.
- *Space Mission Management.* A set of applications that in different combinations form systems to support the planning and execution of space missions. The reengineering goal was identify components in order to improve reliability and facilitate system maintenance. The system is written in Ada.

#### 1.1.2 Reengineering Goals

From this list of case studies some interesting information can be learned. First of all, the goals and motivations for reengineering the software systems are quite diverse, yet some common themes emerge.

- *Unbundling*. Unbundle the software system into subsystems that can be tested, delivered and marketed separately.
- *Performance*. Improving performance is sometimes a goal and sometimes considered as a potential problem once the system is reengineered.
- *Port to other Platform.* Porting to other (user-interface) platforms, sometimes requiring overall changes to the system.
- *Design Extraction*. Always a necessary step in understanding the system; sometimes even an explicit reengineering goal.
- *Exploitation of New Technology*. This may range from new features of the programming language up until upcoming standards (CORBA and UML).

#### 1.1.3 Architectural Problems

Besides the motivations for reengineering problems, the case studies experience recurrent problems that are perceived as key obstacles for achieving the stated reengineering goals. Solving these problems requires significant human intervention since it involves an intimate understanding of and considerable changes to the architecture of the legacy system.

- *Insufficient Documentation*. All of the case studies face the problem of non-existent, unsatisfactory or inconsistent documentation. Tools to document module interfaces, maintain existing documentation and visualise the static structure and dynamic behaviour are required.
- *Lack of Modularity.* Most of the case studies suffer from a high degree of coupling between classes / modules / sub-systems that hampers further software development (compilation, maintenance, versioning, testing). A solution will involve metrics to help detect such dependencies and refactoring tools to help in resolving them.
- *Duplicated Functionality*. In many of the case studies several modules implement similar functionality in a slightly different way. This common functionality should be factored out in separate classes / components, but tools are missing which help in recognising similarities and in restructuring the source code.

• *Improper Layering*. In a few case studies the user-interface code is mixed in with the "basic" functionality, creating problems in porting to other user-interface platforms. A general lack of separation, or layering, is observed with regard to other aspects (distribution, database, operating system) in other case studies. In contrast to a lack of layering, one case study suffers from unnecessary layers. Overly layered modules resulted from each successive developer encapsulating the module with a new concept instead of revising it. This problem needs tool support for defining layers and subsequent correction of broken layers.

### 1.1.4 Code Clean Up

There are quite a number of problems that have to do with "code clean up". Many of these problems arise from the lack of familiarity with the new object-oriented paradigm. But several years of development with sometimes geographically dispersed programming teams that change over time exacerbate these problems. Since they involve behaviour preserving restructuring of code only, those problems could be identified and repaired almost mechanically.

- *Misuse of Inheritance*. Inheritance is used as a way to add missing behaviour to one superclass. This is a often a result of having a method in a subclass being a modified clone of the method in the super-class.
- *Missing Inheritance*. In some cases, programmers have duplicated code instead of creating a subclass. In other parts, long case statements that discriminate on the value of a variable are used instead of method dispatching on a type.
- *Misplaced Operations*. Operations on objects were defined outside the corresponding class. Sometimes this was necessary in order to patch "frozen" designs.
- *Violation of Encapsulation.* This was observed in extensive use of the C++ friend mechanism. Also, software engineers rely on the strong typing of the compiler to ensure certain constraints, and afterwards use typecasts to circumvent the safety-net. In some cases this leads to redundant type definitions which contaminate the name space.
- *Class Misuse*. This problem has been named "C style C++", although it is observed in Ada as well. It refers to the usage of the classes as a structuring mechanism for namespaces only. Sometimes this is necessary to interface with external non object-oriented systems.

#### 1.1.5 Requirements

Last but not least, the case studies impose a number of constraints on the techniques and heuristics presented in this book.

- *Language Independent*. All material in this handbook is applicable on all major object-oriented languages, in particular C++ and Ada, Java and Smalltalk.
- *Scalable*. Some techniques and heuristics scale better than others. Rather than restricting ourselves to those techniques that can deal with small as well as large systems, we choose to specify for every technique the scale of systems it can be applied upon.
- *Tool Support.* There is a heavy emphasis on available tool support for all techniques covered in the book.

## **1.2 Basic Terminology**

Before diving in the specific solutions for object-oriented reengineering, it is useful to agree on some terminology. We rely on the taxonomy of Chikofsky and Cross which is well-accepted within the reengineering community [CCI90]. For terminology specific to the object-oriented paradigm, we draw upon the design pattern book [GHJV95].

- *Reverse engineering*. Originally used for the process of analysing hardware to discover its design, the term refers to the process of recovering information from an existing software system. In general reverse engineering seeks to recover information at a higher level of abstraction such as design information from code. Reverse engineering does not involve modifying the software system: it may be done as a stage in the reengineering process (model capture), as part of an effort to document the system, or as an attempt to extract reusable components from the software.
- *Forward engineering.* Refers to the usual process of software engineering: moving from requirements to high-level design, to progressively lower design levels and to implementation. While it may seem unnecessary to introduce a new term, the adjective "forward" has come to be used where it is necessary to distinguish from reverse engineering and reengineering.
- *Reengineering*. Reengineering is the modification of a software system which in general requires some reverse engineering to be done. That is, reengineering requires that we first recover a view of the system at a higher level of abstraction than the code itself, then make changes to this view and implement these changes at the code level again. Simplistically, reengineering thus involves moving from code to model (reverse engineering), making modifications to the model, and then moving to "better" code (forward engineering).

There is some discussion as to whether or not reengineering involves a change in the functionality of the system – what it does for the user – since practically speaking reengineering almost always modifies the existing behaviour of the system, and indeed is usually motivated by a need to meet new requirements.

- *Restructuring*. Restructuring refers to transforming a system from one representation to another while remaining at the same abstraction level. At implementation level, this usually means changing the code structure without changing the semantics. However, even if the semantics are not changed at implementation level, restructuring might affect higher levels of abstraction (changing design vocabulary without affecting the implementation).
- *Refactoring*. Refactoring is restructuring within an object-oriented context. Refactoring involves tearing apart classes into special and general purpose components and rationalising class interfaces. The principle behind refactorings is that some relatively simple transformations (e.g., renaming a class, renaming a method, moving a method or attribute to another class) are combined into quite powerful semantic preserving transformations (i.e., componentise parts of a class, introduce a bridge design pattern).

## **1.3 The Reengineering Life-cycle**

In this section we present the FAMOOS reengineering life-cycle. We regard reengineering as an evolutionary process consisting of the following six stages (see also figure 1.2):

- 1. Requirements analysis: identifying the concrete reengineering goals.
- 2. Model capture: documenting and understanding the design of a legacy system.
- 3. Problem detection: identifying violations of flexibility and quality criteria.

- 4. Problem analysis: selecting a software structure that solves a design defect.
- 5. Reorganisation: selecting the optimal transformation of the legacy system.
- 6. Change propagation: ensuring the transition between different software versions.

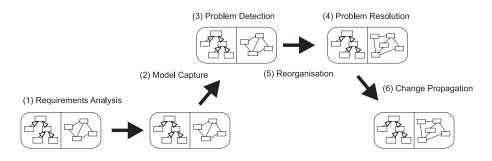


Figure 1.2: The Reengineering Life-cycle

Several iterations of these re-engineering stages might be needed before achieving a stable system with the desired degree of generality and adaptability.

**Requirements Analysis.** The specification of the criteria that the new, reengineered software must fulfill (for example, faster network performance).

**Model Capture.** In order to understand and to manipulate an object-oriented legacy system, it is necessary to capture its design, its architecture and the relationships between different elements of its implementation. A common problem in legacy systems is the lack of documentation. As a consequence, a preliminary model capture is often unavoidable, in order to document the system and the rationale behind its design. This requires reverse-engineering the legacy system to extract design information from the code.

**Problem Detection.** According to the reengineering requirements, problem areas within the legacy systems need to be detected. This requires methods and tools to inspect, measure, rank and visualise software structures. The problem areas have typically properties that deviate strongly form the properties as they are defined in the requirements. Detecting the problems with respect to flexibility requires a definition of these deviations (for example through thresholds on metrics). Problem detection can be based on a static analysis of the legacy system (i.e. analysing its source code or its design structure), but it can also rely on a dynamic usage analysis of the system (i.e. an investigation of how programs behave at run-time).

**Problem Analysis.** Upon detection of possible defects in the legacy system, software developers have to analyse them; that is, match detected problems against unmet requirements and understand how they concretely affect the software. Because applications are organised as intricate webs where classes, objects and methods may participate in various interactions, a detected problem may have to be decomposed into elementary sub-problems. A selection follows of appropriate target software structures - such as design patterns - that impart the software with the desired flexibility and functionality. A combination of such structures may be necessary to handle the particular design defect at hand. A prerequisite for problem analysis is an identification and specification of software structures to serve as the targets of reengineering, and a classification that allows to look for target structures corresponding to particular flexibility criteria or functional requirements.

**Reorganisation.** This phase of reengineering consists in physically transforming software structures according to the operations selected previously. This requires methods and tools to manipulate and edit software systems, to reorganise and recompile them automatically, to debug them and check their consistency, and to manage versions of software.

**Change Propagation.** The process of establishing a revised system throughout a corporate software environment. This might involve reengineering methodology that supports dissemination of improvements in more than one step.

The Need for Object-Oriented Reengineering

Part I

Techniques

## Chapter 2

## Techniques

To reengineer and reverse engineer a software system one needs a range of techniques. This chapter provides a summary of the techniques investigated within the FAMOOS project.

- **Metrics** are definitively appealing as support for understanding huge systems [DD99], [DDL99], [Bau99]. Section 2.1 presents a list of the principal metrics and discussed their possible use and applicability.
- **Program Visualization** is well suited to help to understand huge systems [DDL99], [Lan99]. Section 2.2 presents how the combination of simple graph layouts and metrics gives a quick means to understand and analyse an application.
- Abstracting from the code contributes to the understanding of the system and can help to detect certain flaws [Ciu99]. Section 2.3 presents how grouping entities at another level of abstraction supports a first analysis of the system.
- **Refactoring** is now a well-known technique that helps behavior transforming code change [JO93], [RBJ97b], [Fow99]. Section 2.4 will focus more on advanced techniques based on design pattern based transformations.

## 2.1 Metrics

#### **Author: Markus Bauer**

### 2.1.1 Introduction

*What are software metrics?* – Formally, they measure certain properties of a software project by mapping them to numbers (or other symbols) according to well-defined, objective measurement rules. The measurement results are then used to describe, judge or predict characteristics of the software project with respect to the property that has been measured. Usually, measurements are made to provide a foundation of information upon which decisions about software engineering tasks can be both planned and performed better.

Although software metrics can be used to measure properties of the software development processes as well as organisations that produce software, we will only deal with software product metrics. These metrics measure properties of the source code of a software project and are the most interesting ones within the context of reengineering.<sup>1</sup>

To illustrate the concept of a software product metric, consider one of the most famous software metrics, *Lines of Code*. This metric measures the size of a piece of source code. We use this example to introduce the format which we will use to describe the metrics in this text: Each metric is presented with an acronym and its full name; a scope, explaining what entities of the software system are being measured (the system as a whole, a class,...); a category (the metrics in this text can be grouped into certain categories, see below); a (detailed) description, defining the metric and thus giving the measurement rule; related metrics that measure the same properties and references pointing to the original paper, where the metric has been defined first.

#### LOC - Lines Of Code.

Scope System, Class, Method

Category Complexity

Description *Measures the size of a piece of source code by counting its lines.* Since the size of some source code can be seen as an indicator of its complexity, *LOC* is often used as a complexity metric or as an indicator on how much effort required to implement that piece of code.

The line counting is usually done with respect to a certain coding standard which defines precisely what constitutes a line of code in a particular programming language. This is necessary for obtaining comparable, well-defined measurement results.

See also -

References [Hum97] provides a good discussion of all aspects related to LOC.

Why are software metrics important when reengineering (object oriented) legacy systems? – Software metrics support numerous reengineering tasks, because they help to focus reengineering efforts. They aid in forming an initial understanding of the legacy system and can often uncover hints about design flaws that that are obstructing the modification and extension of the system. Metrics lend themselves we to automatization and with appropriate tools they can provide easy access to meaningful information about the source code without requiring you to read through all the source code by hand. Instead you can use the information to make a more efficient study of the source code based on the points of interest indicated by the metrics results.

<sup>&</sup>lt;sup>1</sup>Note, however, that process metrics or metrics that measure resources of an organisation should still be applied in reengineering projects to support the project management, but this is beyond the scope of this text.

The next section of this text gives an overview over some important object-oriented software metrics and explains some basic properties that can be measured by them. This provides the background needed to present how metrics can be used during reengineering tasks through some some typical usage scenarios based on some of these metrics. We believe that learning how metrics are applied in these usage scenarios will illustrate ideas on how to use metrics in your own reengineering projects.

#### 2.1.2 Some Important Metrics

In this section, we present some object oriented software metrics, that have proven to be particularly use- $ful^2$ . These metrics fall into several categories depending on thes aspects of a system they measure. We have identified the following categories: *complexity* metrics, *coupling* metrics, *cohesion* metrics and *inheritance tree* metrics.

#### 2.1.2.1 Complexity Metrics

Complexity metrics measure the *complexity* of an entity of a system. The metrics presented here measure the complexity of a class. By the term complexity of a piece of software or source code, we usually try to describe how much effort has to be spent by a software engineer to understand, write or modify that piece of software – code that is difficult to read and understand is considered as complex. Since measuring the complexity directly is not possible (since we would have engineers read the code and check how much time they needed to understand it), we use some metrics to estimate that complexity. The metric *LOC* (p. 22), mentioned above, is an example of such a complexity metric.

One of the well established metrics to measure the complexity of a class is:

#### WMC - Weighted Method Count.

Scope Class

Category Complexity

Description Measures the complexity of a class by adding up the complexities of the methods defined in the class. Thus,

$$WMC = \sum_{i=1}^{n} c_i$$

where  $c_i$  denotes a complexity measurement of method *i*.

Complexity measurements for methods are usually given by code complexity metrics like *LOC* (p. 22) or the *McCabe cyclomatic complexity*. The McCabe cyclomatic complexity measures the complexity of some code by taking into account the decision structure of the code, i.e. code that contains a lot of loops or *if-then-else*-constructs is considered more complex.

See also *NOM* (p. 24) is a special case of this metric – all method complexities are assumed to be 1. References [CK94], [CS95], [EBD99]

<sup>&</sup>lt;sup>2</sup>Chapter **??** contains a survey and critiques on a large amount of metrics, we used that survey to select the metrics presented here. Additionally, we recommend [LK94] as a good textbook on metrics.

A special case of WMC (which is very simple to compute) is:

#### NOM – Number Of Methods.

Scope	Class
Category	Complexity
Description	Measures the complexity of a class by counting the number of methods defined in that class.
See also	<i>WMC</i> (p. 23)
References	[HS96a]

Obviously, complexity metrics play an important role when reengineering software systems: classes with high complexity measurements are difficult to understand and consequently difficult to change. For details, check the scenarios in sections 2.1.3.1-2.1.3.4.

#### 2.1.2.2 Coupling Metrics

Another important aspect when dealing with a legacy system is the coupling between classes. A class is *coupled* to another class, if it depends on (or knows) that class, for example by accessing variables of that class, or by invoking methods from that class.

#### DAC – Data Abstract Coupling.

Scope	Class
Category	Coupling
Description	Measures coupling between classes that results from attribute declarations.
	DAC counts the number of abstract data types defined in a class. Essentially, a class is an abstract data type, therefore $DAC$ reflects the number of declarations of complex attributes, i.e. attributes that have another class of the system as a type.
See also	<i>RFC</i> (p. 24), <i>CBO</i> – <i>Coupling Between Objects</i> [CK94], <i>NIV</i> – <i>Number of Instance Variables</i> [LK94]
References	[LH93], [HM95], [HM96]

The following metric is a coupling metric as well; however, the complexity of the class affects the measurement, thus it cannot be considered as a pure coupling metric.

#### RFC – Response Set For A Class.

Scope	Class
Category	Complexity, Coupling
Description	Measures complexity and coupling properties of a class by evaluating the size of the response set of the class, i.e. how many methods (local to the class and methods from other classes) can be potentially invoked by invoking methods from the class.
	More formally, $RFC$ for a class $C$ is defined as $RFC =  RS $ , where the response set $RS$ is given by

$$RS = M \cup \bigcup_{m \in M} R_m$$

M is the set of methods defined in C and  $R_m$  is the set of methods called by method  $m \in M$ . See also DAC (p. 24), MPC – Message Passing Complexity [LH93] References [CK94]

Why are we interested in the coupling between classes? – Classes, that are tightly coupled cannot be seen as isolated parts of the system. Understanding or modifying them requires that other parts of the system must be inspected as well. Conversely, if other parts of a system get changed, classes with high coupling measurements are more likely to be affected by these changes (see scenario in section 2.1.3.3). Additionally, classes with high coupling tend to play key roles in the system, making them a good starting point when trying to understand an unfamiliar legacy system (see scenario in section 2.1.3.1).

#### 2.1.2.3 Cohesion Metrics

The *cohesion* of a class describes how closely the entities of a class (such as attributes and methods) are related. Often, cohesion is measured by establishing relationships between methods of the class in the case where the same instance variables are accessed. A useful metric measuring this property is:

#### TCC - Tight Class Cohesion.

Scope Class

Category Cohesion

Description *Measures the cohesion of a class as the relative number of directly connected methods*, where methods are considered to be connected when they use at least one common instance variable.

More formally TCC for a class C is defined as follows: Let

 $NDC = |\{(m, n) \mid \text{ methods } m, n \text{ access a common instance variable}\}|$ 

be the number of connected methods and  $NPC = \frac{n(n-1)}{2}$  the possible number of connected methods, then

$$TCC = \frac{NDC}{NPC}$$

See also *LCOM – Lack of Cohesion On Methods* [CK94] References [BK95], [HM95], [HM96], [EDL98]

Cohesion is an important concept: good object oriented design styles usually require that classes have high cohesion, since they should encapsulate concepts that belong together. Classes with low cohesion often represent violations to a flexible, extensible or reusable design. All of these are issues that must be dealt with during reengineering projects. The scenario in section 2.1.3.2 further deals with this issues.

#### 2.1.2.4 Inheritance Tree Metrics

A basic concept that is central to object-oriented systems is *inheritance*. Through inheritance, relationships between objects can be modelled (for example the *is-a* relationship). Moreover, inheritance is often used to allow for some reuse of existing classes. Accordingly, measuring properties of the inheritance tree of a system often gives interesting results. Some simple metrics relating to the inheritance tree and its layout are:

#### **DIT – Depth in Inheritance Tree.**

ScopeClassCategoryClass HierarchyDescriptionMeasures the depth of a class in the system's inheritance tree.The DIT-value for a class C is defined as length of the longest path in the inheritance tree from the root class of the system to C.See also—References[CK94]

#### NOC - Number Of Children.

Scope Class Category Class Hierarchy Description *Counts the number of children (direct subclasses) of a class.* See also *NOD* (p. 26) References [CK94]

#### NOD - Number Of Descendants.

Scope	Class
Category	Class Hierarchy
Description	Counts the number of descendants (direct and indirect subclasses) of a class.
See also	<i>NOC</i> (p. 26)
References	[TSM95]

Obviously, the inheritance metrics presented above may be used to measure a special case of coupling – the usage of classes through inheritance relations. For example, classes with low *DIT* values and high *NOC*- or *NOD*-values are classes that affect a lot of other classes, because they are (direct or indirect) super classes to them. Changes to such classes are likely to require changes in the subclasses. We believe, however, that inheritance based properties of a system are more easily understood through visualising the inheritance tree.

Recently, more sophisticated metrics have been defined that measure the amount of reuse in an inheritance tree. Because of their complexity, these metrics are outside the scope of this text, we refer you to chapter **??** for a description of these metrics.

#### 2.1.3 Usage Scenarios

In this section, we present some scenarios illustrating how metrics can be applied successfully in software reengineering projects.

#### 2.1.3.1 Get a Basic Understanding of the System

Usually, one of the initial steps when reengineering a legacy system is to aquire a basic understanding on how the system works and how it is structured. However, the documentation alone is typically insufficient. Therefore some analysis of the system's source code is required. This is (*model capture*) and metrics can provide valuable help during this task.

A good way to start model capture is to find out which parts (i.e. classes) implement the key concepts of the system. A technique to do this is described in [Bau99]: Usually, the most important concepts of a system are implemented by very few *key classes*<sup>3</sup>, which can be characterised by the following properties: Key classes *manage* a large amount of other classes or *use* them in order to implement their functionality, thus they are tightly *coupled* with other parts of the system. Additionally, they tend to be rather *complex*, since they implement much of the system's functionality.

Based on this observation, it is straightforward to identify these key classes by using both a complexity metric like *WMC* (p. 23) and a coupling metric like *DAC* (p. 24). Figure 2.1 shows a diagram that can be used for such an analysis – the classes of the legacy system are placed into a coordinate system according to their complexity and coupling measurements. Classes that are complex and tightly coupled with the rest of the system fall into the upper right corner and are good candidates for these key classes. To understand how the legacy system works we should thus concentrate on understanding these classes and their interactions by studying their source code.

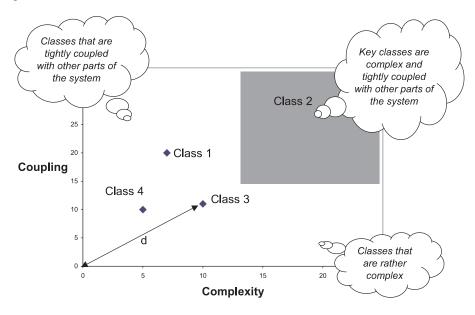


Figure 2.1: Finding the key classes in a legacy system

#### 2.1.3.2 Find Violations of "Good" OO Design

Though reengineering projects are usually started in order to make the legacy system satisfy additional functional or non-functional requirements, general improvements to the software are often desired as well.

<sup>&</sup>lt;sup>3</sup>Case studies in [Bau99] have shown that about 10 % of the classes of a system can be considered as key classes

One way to achieve such general improvements consists in finding violations of a "good" object oriented software design (*problem detection*). Unfortunately, there is no consensus on what "good" design really is, however, some guidelines and principles exist that are considered helpful to achieve an understandable, flexible and extensible software design. Metrics are particularly suitable to check, whether the legacy system adheres to such design principles or to find violations of them.

A (very basic) principle in object oriented software engineering states that a class should implement *one single concept* (of the application domain). Some violations of that principle can be detected by using software metrics if we make the following assumptions:

- A class that implements more than one concept, has probably *low cohesion* measurements, since these concepts can be implemented separately.
- A class that by itself does not implement one concept, i.e. the implementation of the concept is distributed among many classes, is probably *tightly coupled* to other classes.

Therefore, by applying cohesion metrics like *TCC* (p. 25) and coupling metrics like *DAC* (p. 24) or *RFC* (p. 24) to the legacy system, possible violations of the principle "one class – one concept" can be found. These classes tend to have either low *TCC*-values or high *DAC-*, *RFC*-values.

For example, classes that have very low *TCC*-values, can often be split. Sometimes this leads to a more flexible design, since the two separate classes are easier to understand and are more reusable. Low *TCC*-measurements can as well point to classes that have not been designed in an object-oriented flavour at all – these classes are not implementing a self contained object from the application domain, they just group methods together, acting as a module. In a similar manner other design principles can be checked and violations can be detected by using metrics [Rie96].

However, we should be aware of some difficult issues, when applying metrics for such problem detections: It is difficult to specify thresholds for the measurements, i.e. values, which classes adhering to a "good" design should fulfill. Additionally metrics can produce "false alarms". They can label classes as being problematic, but there may well be a reason that these classes present untypical measurement values. Measurement results must always be taken with a grain of salt and problematic results should always be checked against the source code.

#### 2.1.3.3 Identify Change Sensitive Parts

Whenever we make changes to an existing software system, it is likely that these changes will require further changes throughout the system since the entities of the system are interdependent. Changes in some parts of the system can produce a lot of work, if a lot of other parts depend on them, and, inversely, some parts probably change often during the evolution of a system, because they depend on lots of other parts and changes to the system are likely to affect them.

To make sure that the system does not misbehave after making some changes, we would be interested in localizing these *change sensitive* classes, i.e. classes that are most likely to be affected by changes to a system because the depend on lots of other parts. To do this, we can use coupling metrics like the *DAC* (p. 24). Classes with high *DAC*-values "know" a lot of other classes and are therefore change sensitive. These classes should then be carefully examined and tested after modifications of the system.

#### 2.1.3.4 Track the Evolution of the System

Most software systems *evolve* over the time, i.e. new functionality is added, extensions are made,.... This poses an important question: Does the quality of our system decrease during the evolution of the system? Do some reorganisation raise the quality of the system?

Metrics can be used to answer these question and to control the quality of the software. A lot of research work has been done in this area, see for example [LK94], [EL96] or [DD99].

The basic steps of using metrics for quality control are:

- 1. Establish quality goals for your software.
- 2. Decide on a set of metrics to check your software with respect to the quality goals.
- 3. Use the metrics to constantly monitor the quality while the system evolves.

A simple example: a very high level quality goal for a software system could be maintainability, thus, coupling measurements should not be high in order to ensure that changes to the system do not trigger changes throughout the system (see 2.1.3.3). Therefore, monitoring *DAC*-values can be promising. When a significant number of classes evolves to higher *DAC*-measurements, some refactorings of the system could be appropriate, to reduce coupling.

Another application of metrics when tracking the evolution of a system is to identify stable and unstable parts of your system (for details, see [DDN99]). Often, this can be interesting information: Stable parts can be declared as "frozen" and can often be reused in other projects (i.e. factored out into frameworks), whereas unstable parts should be tested thoroughly.

#### 2.1.4 Summary

In the previous sections we have seen what object oriented software metrics are, and how they can be applied in reengineering projects. We have illustrated that metrics are able to support model capture and problem detection phases. Figure 2.2 sums up our experiences with the metrics mentioned in this text and gives hints on how well these metrics are suited for model capture or problem detection.

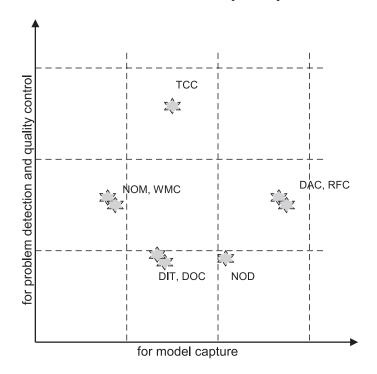


Figure 2.2: Applicability of the metrics

Reengineering projects can benefit from metrics in the following ways: Applied with well documented scenarios (as given in section 2.1.3), they make reengineering tasks more organised and focused. They provide an abstraction mechanism from the huge amount of source code of the legacy system, thus allowing you to concentrate your work on the important or critical parts of the system that have been identified by the measurements.

However, metrics can fail. They can point to wrong places in your source code, or you can even miss important classes of your system, because the measurements just do not "catch" them. Therefore, metrics should be used with care, and you should always check with the system's source code or make additional analyses to avoid drawing wrong conclusions.

## 2.2 Program Visualisation and Metrics

#### Author: Michele Lanza

"Continuous visual displays allow users to assimilate information rapidly and to readily identify trends and anomalies. The essential idea is that visual representations can help make understanding software easier." [BE96]

Although the object-oriented paradigm lets programmers work at higher levels of abstraction than procedural models, the tasks of understanding, debugging, and tuning large systems remain difficult. This has numerous causes: the dichotomy between the code structure as hierarchies of classes and the execution structure as networks of objects; the atomisation of functionality - small chunks of functionality dispersed across multiple classes; and the sheer numbers of classes and complexity of relationships in applications and frameworks. The fields of scientific visualisation and program visualisation have demonstrated repeatedly that the most effective way to present large volumes of data to users is with a continuous visual fashion [PHKV93].

In this chapter we list some properties that a graphical representation of source code should possess to be useful for reverse engineering. We then see in what respect our approach fulfils those requirements and include a short scenario to explain our approach. We also list some problems concerning the visualisation of metrics, colors and issues concerning interactivity.

The central point of this chapter is to show how we merge the concepts of program visualisation, metrics and interactivity. These three aspects are the cornerstones of this work. The concepts that are explained here have been implemented in a single tool called CodeCrawler, which we present in the next chapter.

#### 2.2.1 Graphs for Reverse Engineering

In this section we list some features that in our eyes graphs for reverse engineering should have. We emphasise that we use the term graph in a very broad sense: often we mean its picture or graphical representation on screen or on paper and not necessarily its scientific definition.

- Simplicity and Quality. The first important prerequisite is that the generated pictures of a graph have to be relatively simple and easy to grasp. The main reason for that is that too much displayed information overloads the viewer's perception. This tends to backfire and causes an unwanted information loss. A secondary aspect is that simple graphs are also easily reproducible, while complex techniques like hyperbolic trees [LRP95] are affected by a considerable complexity which is hard to grasp and reproduce. Many approaches have been discussed as to how a software entity could be represented for program visualisation ([BE96, PHKV93, KG88] to name but a few). We think that a graphical representation of an object oriented entity should be easy to grasp and not make use of a specific dictionary of shapes which has first to be learned. A graph should be able to transmit useful information to the viewer at first sight.
- *Quantity*. We have to be able to select how much of the subject system we want to display and at what level of granularity. Thus, we should be able to zoom in and out of such a graph and reduce the amount of displayed information at will.
- Colors. Program visualisation can be supported by colors, because they can attract the eye to interesting hot spots, while other parts of the graph which look less colorful can be ignored by the viewer. Colors have often been used in program visualisation [Riv98]. While colors are a good way to attract the attention of the eye, the usage of too many colors in a graph is not advised, since this results in an optical overload for the viewer of the graph. We also advise against the use of color conventions which have first to be learned by the viewer, as this lessens the impact of the colors.

- *Scalability*. As reverse engineering is especially crucial in very large systems, a visualisation should be scalable and work if possible at any level of granularity. The number of displayed entities should not affect the quality of the graph.
- *Interactivity*. A very important aspect of graphs is not only their layout algorithm but also that they can provide interactivity to the user through direct-manipulation interfaces. Making a static display of nodes and trying to extract information from the graph has clearly defined limits, which we discuss below in Section 2.2.2.
- *Metrics*. Although intangible in the physical sense, software *has* size. It can be measured, especially in object-oriented code we can assign numerical values (metric measurements) to its entities. Although the concept of software is abstract and often exists only in the head of the programmer, we can measure it. Once we can measure it, we can assign a size to it and represent this size graphically. We think that metrics enrich the semantic value of a graphical representation of a software entity, and discuss this below in Section 2.2.4.

#### 2.2.2 Interactivity

A graph which lacks interactivity has certain drawbacks:

- 1. The user can't produce new views starting from a part of the graph.
- 2. The user can't find out secondary information (e.g. he can't inspect the nodes or browse through their source code).
- 3. The user can't reduce the amount of displayed data by either removing nodes by hand or by filtering out nodes through algorithms.

Those limits can be overridden if the graph is interactive:

- If we produce a view on a system and one particular node is drawing our attention, we'd like to know more about this node and the entity that it is representing. So we should be able to know its name, to have a look at its properties, to zoom in into the node, to have a list of all nodes that have a relationship with this node, or even to have a look at the source code behind the node (suppose the node is a method).
- Starting from a part of the graph or from one single node we'd like to be able to generate new views without having to go through the whole graph generation procedure again. The viewer should be able to 'navigate' around the code travelling from one point of interest to the next.
- Sometimes the relationship edges in a graph make the whole graph look like a cobweb. We should be able to switch off edges and switch them on again on demand depending on nodes we selected, etc.
- Suppose we have displayed a graph with a lot of nodes and edges. One particular node is of interest to us. But since there are too many edges in the graph it's hard to see how many times and to which other nodes the node in question is connected. So the graph should also be able to provide a 'highlighting feature' where we can display on top of all edges and nodes the connections of the node in question. It is important to note here that compared to the previous point we don't want to reduce the complexity of the displayed graph. We just want to have a better view on it.

It is an important point we are stating: The interactivity of a graph is *not just a nice feature* but one of its *most important aspects*.

#### 2.2.3 The Use of Layout Algorithms

Perhaps the most difficult aspect of showing software through graphs involves the graph layout problem. The nodes and edges of the graph must be positioned in a pleasing and informative layout that clearly shows the underlying graph's structure. Many techniques have been proposed for laying out arbitrary graphs. Unfortunately, in practice, drawing informative graphs is exceedingly difficult, particularly for large systems. The resulting graphs, even when drawn carefully, are often too busy and cluttered to interpret [BE96].

The opposite case can also be true: sometimes elaborate layout algorithms can't ameliorate the user's perception or can do that only at the cost of algorithm complexity: There are various (and sometimes very complex) techniques to display a tree graph, but in the end it's still just a tree.

However, we don't want to minimise the importance of complex layout algorithms, on the contrary: we believe they could bring many more benefits than drawbacks. Good layout algorithms just were not part of the constraints of this work. But it is certainly a very promising field of research in this context. The layout algorithms used in this work are discussed in detail in Section 2.2.6.

#### 2.2.4 The Use of Metrics in Graphs

In [BE96] the following statement is made: "Software is intangible, having no physical shape or size. Software visualisation tools use graphical techniques to make software visible by displaying programs, program artifacts, and program behaviour."

It is obvious that everything regarding metrics possible through their graphical display is also possible by just calculating and analysing the metric measurements. So the question arises why we should have a graphical display of them, since the information sought is in the metrics themselves. But in the same way one could think to listen to music by just reading the partiture of a song instead of using the sense normally designed for that (the hearing)<sup>4</sup>. *What changes is the perception and the impact of what is perceived.* 

**Our Idea.** The whole concept is fairly easy: we map metric measurements of software entities on their graphical representation on the screen. As we said before we chose the entities to be represented by rectangles. Rectangles have a certain width and a certain height. They can be filled with a color. Their position can also bear a certain amount of information.

With this approach, in a two-dimensional graph consisting of nodes and eventually edges between the nodes, up to five metrics can be assigned to a node and rendered visually at the same time. These are:

- 1. The X coordinate of the position
- 2. The Y coordinate of the position
- 3. The width
- 4. The height
- 5. The color shade

This concept is rendered clearly in Figure 2.3, where we see where the metrics can be applied on a node.

Not every graph can make use of five the metrics at the same time. In a graph that does not have an origin (which defines an absolute coordinate system) it makes no sense using two metrics for the position of the

 $<sup>^{4}</sup>$ A short comment on perception: the size of software can be seen through other means: if we scroll through the source code of a very large class, we probably have to either move the mouse or press some keys on the keyboard to scroll on. This physical act of scrolling can also transmit size and complexity to the viewer.

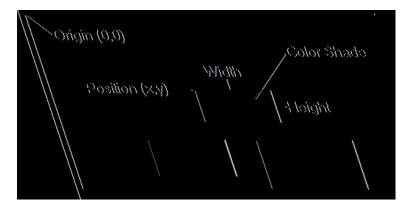


Figure 2.3: An example of nodes and their possible metrics.

nodes. A good example for such a graph is a tree graph, where the position on the nodes is implicitly defined by the logical position of the nodes in the tree. Another property which came up during our experiments was that sometimes the multiple use of the same metric (for example if we choose the same metric to reflect width and height) can emphasise certain parts of the graph and render them more clearly for the viewer.

#### 2.2.5 A Concrete Graph Specification.

In our approach a *concrete graph*, this means the resulting displayed graph, is the combination of four factors :

- 1. **The Graph Type**. Its purpose is to render a certain aspect of a system: a tree graph is good for displaying hierarchical information, a circle for communication, a confrontation graph for dependencies, etc.
- 2. **The Layout Algorithm**. Starting from the original idea of the graph, variations refine the concrete display. The layout takes into account the following issues:
  - Display concerns (i.e. the fact that the complete graph should or not: fit into the screen, minimise the space used, sort the nodes according to certain criteria, etc.).
  - The entities and their relationships. This implies the choice of the represented entities (class, attribute and/or method) to be rendered as graphic elements and the logical link between the graphical elements and the metrics. For example in some graphs the position of the nodes reflects the size of the entities whereas in others this is the size of the node.
- 3. **The Metric Selection**. Once the layout algorithm stands, metrics are associated to the graph. This application depends on the specification of the previous step.
- 4. **The Interaction**. Since the goal of a graph is to support the reverse engineering of the application, the interaction that a user can perform should be specified. All the graphs support basic navigation functionality which allows one to access code elements. However, the interaction is refined for specific graphs, for example to walk through it, to highlight the edges, to zoom in/out, etc.

#### 2.2.6 A Short Example

To make the whole idea of visualising software structures with the help of metrics a bit more understandable we included here a short example of our approach.



Figure 2.4: A simple inheritance tree.

Suppose we want to understand the inheritance hierarchy of a small system. The idea that comes up is to display the graph as a tree. The nodes in tree represent classes, the edges represent inheritance relationships.

The layout algorithm for displaying a tree is arbitrary, for reason of simplicity we chose a very simple one, which sometimes can make edges cross nodes, but it renders the whole concept nonetheless. Keep in mind that this layout part can also make use of very complicated algorithms for space optimisations, edge crossing reduction, etc.

Once we have displayed the tree as we see in Figure 2.4 we apply size and color metrics to the nodes. The use of position metrics is not possible here, as the position of the nodes is intrinsically defined by their logical position in the tree. As the nodes represent classes, we use class size metrics. The width and height of the nodes render the number of methods (NOM) while the color renders the number of attributes (instance variables).

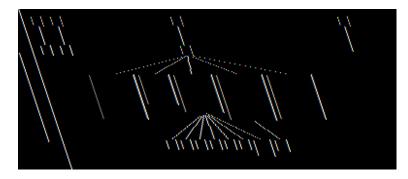


Figure 2.5: An inheritance tree that makes use of size and color metrics.

Once the tree is rendered as in Figure 2.5 we can start interacting with the graph. We can freely move nodes around, delete them, inspect them (i.e. browse the underlying classes), filter out parts, etc.

In fact, if we left out the interactive part, the amount of useful information that we could extract would be limited to the display in Figure 2.5.

## **Useful Graphs**

This section is dedicated to the graphs which prove to be useful when it comes to the understanding of software systems and the detection of design problems using the approach discussed in this work. Although this may seem a little confusing, what in this chapter is called a 'useful graph' is not only its layout, but primarily the *combination of a layout with object-oriented metrics and the consequent extraction of information made by the viewer through interaction with the graph.* 

The following is structured as follow:

- Graph Structure. Section 2.2.7 presents the structure we adopted to describe a useful graph. Every useful graph is presented using this format.
- Case Studies. Section 2.2.8 is a short presentation of the two case studies we chose to apply each useful graph on.
- Layout Algorithms. Section 2.2.9 presents the layout algorithms we selected.
- Useful Graphs. In section 2.2.10 we present the useful graphs divided into the 4 distinct groups: *class, method, attribute* and *class internal*. The names indicate which kind of entities are displayed in the graphs. Class internal treats the special case where methods and attributes are displayed at the same time.

#### 2.2.7 Graph Structure

For each graph which we treat in this chapter, we discuss the following properties:

- **Graph:** Indicates what type of graph and layout has to be chosen, and whether a sorting of the nodes has to precede the display.
- **Scope:** At what granularity level the graph can be applied. We differentiate between *full system*, *subsystem* and *single class*. Sometimes the subsystems are indicated as a single inheritance hierarchy. We also indicate if the graph is language specific.
- **Metrics:** We list five metrics in the following order: width metric, height metric, color metric, horizontal position metric, vertical position metrics. When we write a dash (-), this means that the metric should not be set. In case we write an asterisk (\*) this means that the metric can be set freely. In the case of class internal graphs we repeat the five metrics twice, once for the method nodes and once for the attribute nodes.
- **General idea:** We write what the graph is all about and what ideas lie underneath it. We also indicate what the user should be searching for in the graph.
- Results: Here we present the results obtained after applying the graph on our case studies.
- **Possible Alternatives:** We list a few alterations that could be made regarding the metrics, so as to obtain slightly different graphs, and list also some eventual interactions that could be applied on the graph to increase its usefulness.

Evaluation: Some statements about the advantages and drawbacks of the graph.

Application	Refactoring Browser	Duploc
Classes	166	123
Methods	2365	2382
Attributes	365	386

Table 2.1: An overview of the size of our case studies.

# 2.2.8 Case Studies

This section contains a short overview of the systems we used as case studies for this work. Basically we use them to test the graphs listed in the remainder of this chapter. We chose these two case studies for the following reasons:

- Availability. Both case studies are public domain and can be downloaded freely. With this point we can ensure that the results are reproducible.
- Size. We chose two case studies which can be termed as being of an *average size* and are representative of medium-sized standalone applications. We think that very small applications can't reflect results properly because the purpose of most graphs is coping with complexity, which in such cases is not necessary. On the other hand, if we had chosen very big applications, it would have been hard to present results in a concise manner, because many graphs can be applied in various areas and at various levels of granularity.
- Level of maturity. We chose one very mature application which has gone through some refactorings and redesigns, and another one which has been developed in a rush and which has yet to undergo its first redesign. We did this to see if the results of our experiments would differ and in what way they would do that.

**Refactoring Browser.** The Refactoring Browser is a widely used tool for the implementation of Smalltalk programs [RBJ97b]. We took it as a case study because it is an application which has gone through several refactorings and redesigns and has been written by some very experienced programmers. This quality of implementation should thus be visible in such a system. It is a medium sized application as we can see in table 2.1.

**Duploc.** Duploc is a tool for the detection of duplicated code [RD98]. Duploc was the first application written in Smalltalk by its developer, Matthias Rieger and has yet to undergo its first major redesign. Thus we expect it to have some of the flaws which new systems tend have, like oversized classes and methods, obsolete attributes, etc.

# 2.2.9 Graphs

This section is dedicated to the graphs and layouts we have selected to implement in CodeCrawler. We discuss the properties, advantages and drawbacks of each one of them. We include this here because they are mentioned throughout the remainder of this chapter.

We discuss the original idea of a graph and the scope of its applicability. Each graph has at least one possible kind of layout and we discuss it with a regard for the metrics that can be applied for that special layout. Sometimes a sorting of the nodes has an influence on the usefulness of a graph and we discuss that as well as the general pros and contras for each graph.

In Table 2.2 we have an overview of all graphs and their properties supported by CodeCrawler.

Graph Type	Metrics	Entities	Sort	Scope
Tree	3	С		Global
Correlation	5	CMA		Global- Local
Histogram	3	CMA	Х	Global- Local
Checker	3	CMA	Х	Global- Local
Stapled	3	CMA	XX	Global- Local
Confrontation	3 + 3	MA	Х	Local

Table 2.2: CodeCrawler's graph layouts.

The 'Metrics' column specifies how many metrics can be rendered by the graph. 5 means that the a single node can render 5 metrics at the same time. 3 + 3 means that two separate groups of entities and metrics can be defined. The 'Entities' column refers to the kind of entities the graph can be applied upon: C for class, M for method and A for attribute<sup>5</sup>. The 'Scope' column specifies if the graph can be applied to the complete (sub)system or only to some entities like a class or a method. The 'Sort' column indicates if a sorting of the nodes according to a certain metric measurement can enhance the usefulness of the graph in question.

<sup>&</sup>lt;sup>5</sup>The limitation to these three types of entity is due to the current implementation of the Moose model. Future implementations of it may include supplemental entities.

### 2.2.9.1 The Tree Graph

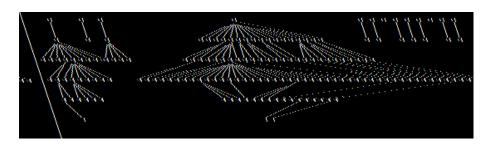


Figure 2.6: A tree graph of a system.

**Overall Idea.** A tree graph is useful for the display of hierarchical structures like inheritance hierarchies containing classes. The nodes represent classes, while the edges between the nodes represent inheritance relationships.

**Scope.** The scope of this graph ranges from very large systems to subsystems consisting of few classes. A requirement is that there is some usage of inheritance in the system. Otherwise the graph gets very flat and wide.

**Layouts.** We implemented three slightly different layout algorithms, which we simply called left, centered and right. Each one of them is based on recursion.

**Metrics.** The number of possible metrics that can be applied is 3. The two position metrics cannot be used, as the position of the nodes is defined by the layout algorithm. However, a virtual fourth metric is present, HNL. It is rendered by the layout algorithm through the vertical position of the nodes.

**Sort influence.** This graph is one of the few cases where a sorting of the nodes is not advised, as it disturbs the recursive layout algorithm.

**Pro et contra.** The advantage of this graph is that it can render a complex system in a very simple manner. Its only drawback is that because the position of the nodes is defined by the layout algorithm, this graph tends to get very large for big systems and will sometimes not fit on one single screen. The use of node shrinking can alleviate this problem.

### 2.2.9.2 The Correlation Graph

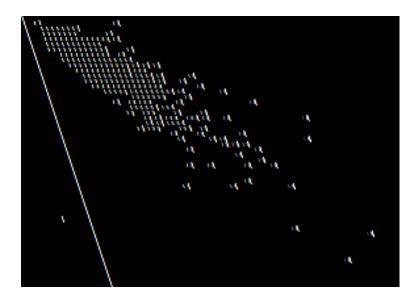


Figure 2.7: A correlation graph of method nodes using LOC and NOS as position metrics.

**Overall Idea.** This graph can render the relationship between two metrics when they are applied to entities. The two metrics are directly mapped onto the position coordinates of the nodes. This graph needs an absolute origin within a coordinate system, which in our case is the upper left corner of the graph. If the chosen metrics are in close relation to each other, the nodes are positioned along a certain correlation axis, which is defined by the metrics. If a node finds itself far away from this correlation axis, it means that its metric measurements are somehow abnormal compared to the other nodes and should be inspected. Very large measurements put a node far away from the origin, if one of the two position metric measurements is very small, the node finds itself near the left or top border of the graph.

**Scope.** This graph can be applied to any type of entity. The maximum number of displayable nodes is very big, as the expansion of the graph drawing depends on the outliers in the system and not on the number of displayed nodes. This involves an overlapping of nodes, which however is not negative, because we are mainly interested in the outliers (i.e. the extreme values).

**Layouts.** There is only one possible layout in this case, which directly maps the position metrics to the position of the nodes.

**Metrics.** The number of possible metrics that can be applied is 5. Indeed, each metric can be applied in this case. However, if we choose to select size metrics this involves that the nodes overlap, while without size metrics the nodes will either be positioned next to each other or cover up other nodes entirely. The overlap problem is especially acute when the chosen size metrics tend to have big values, like LOC.

Sort influence. A sort has no influence on the layout.

**Pro et contra.** The main advantage of this graph is its scalability. Another advantage is that we can pick out the outliers at one glance. The drawback is a certain loss of overview, because the nodes overlap. However, as we often do not make use of size metrics for this graph, we can circumnavigate this problem.

### 2.2.9.3 The Histogram



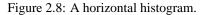




Figure 2.9: A horizontal histogram using the size addition layout

**Overall Idea.** A histogram provides a representation of the distribution of entities related to a certain metric. The distribution of the nodes can in turn give us general information about a system. For example if we use as vertical position metric LOC of methods, we are able to gather if the methods tend to be overlong or not, and if there are any significant outliers.

**Scope.** This graph can be applied to any type of entity, class, method or attribute. The number of displayable nodes is also very large. However, since a large part of the nodes distribute around a certain value, a few of the rows of this graph can get very large and eventually get bigger than the screen. This problem is sometimes acute if we use the size addition layout described below. One of the fields where its use is advised, is to make a distribution of the methods of single classes or of attributes of subsystems.

**Layouts.** There are two possible layouts. The first, called *horizontal*, ignores size metrics and displays every node with the same size. The second one, called *size addition*, makes use of the width metric, and puts the nodes next to each other, while taking their size in consideration. Only the horizontal layout can be considered to be a real histogram, the kind which is used in the field of statistics.

**Metrics.** The number of possible metrics depends on the used layout. The horizontal layout can make use of 2 metrics, namely the color and the vertical position. The size addition layout can also make use of the width metric.

**Sort influence.** In the case of the horizontal layout, a sort has a positive effect if we take the color metric as sort criterion. It makes the detection of color metric outliers easier. In the case of the size addition layout, a sort according to the width metric also has some positive effect for the detection of width metric outliers.

**Pro et contra.** This graph shows a good behaviour in terms of scalability. Its major drawback is that the vertical position metric needs to have a rather large measurement interval, otherwise the nodes will be distributed all near the same vertical position.

## 2.2.9.4 The Checker Graph

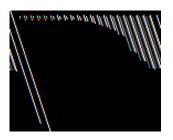


Figure 2.10: A checker graph using a sorted horizontal layout.

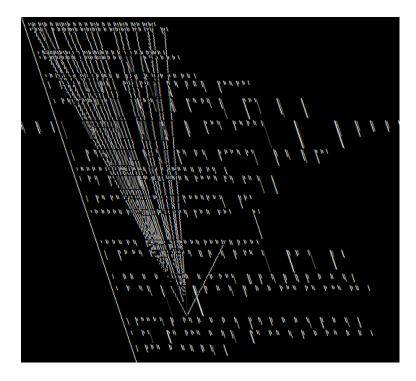


Figure 2.11: A checker graph using a quadratic layout with method nodes and invocation edges.

**Overall Idea.** The base idea for this kind of graph is simplicity. We want to lay out nodes without a special algorithm, we just place them one next to each other, to prevent them from overlapping.

**Scope.** This graph scales up quite well (especially if node shrinking is applied). Therefore it can be used for any kind on entity. However, it's not advisable to use edges in this graph, because it looks very chaotic, as they will cross the nodes.

**Layouts.** The first layout kind is called *horizontal* and *vertical*. We just place the nodes next to each other. We see such a layout in Figure  $2.10^6$ . Because this wastes a lot of space, we introduced the *quadratic* layout which tries to lay out the nodes to make them form a rectangle, whose width is dependent of the number of

<sup>&</sup>lt;sup>6</sup>This figure suggests that a histogram is a special case of a checker graph. This is not true: a histogram makes use of a more complex layout algorithm which makes use of position metrics, as we see in the following subsections.

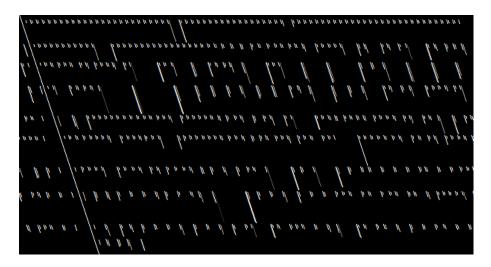


Figure 2.12: A checker graph using a maximal space usage layout.

displayed nodes. The graph which makes the best use of space is called *maximal space usage*, which tries to put as many nodes on the visible part of the drawing as possible.

Metrics. As the position metrics can't be used in this graph, we can only use size and color metrics.

**Sort influence.** The sort is essential for this graph. Indeed, if we don't make use of it, the nodes are placed randomly on the screen and it will be very hard to discern significant nodes. If we do make use of a sort according to a metric (especially the width metric), the detection of outliers will be very easy.

**Pro et contra.** The advantage is that we end up with a very easy to analyse layout. If the nodes are sorted, the detection of outliers is very easy, and the detection of suspicious node shapes is easy as well. This graph scales up well and several hundreds of nodes can be displayed at the same time without overlapping.

# 2.2.9.5 The Stapled Graph

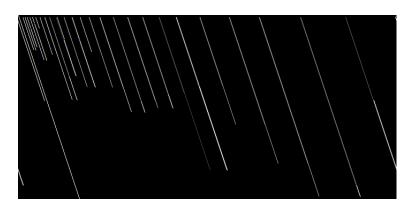


Figure 2.13: A stapled graph of class nodes.

**Overall Idea.** The idea for this graph came up when we tried to cure a small flaw in the horizontal checker layout: The width of the whole graph is defined by the summed widths of the nodes and cannot be influenced by the user. In such cases it often happens that the checker graph is wider than the screen. The stapled graph is thus a derivate: the user can indicate the maximum width of the graph he'd like to have, and all the node are accordingly shrunk in their width to make the graph fit the indicated space.

Scope. This graph can also display any kind of entity.

Layouts. A this time there is only one possible layout, which displays the nodes horizontally.

Metrics. The size and color metrics can be used, while this is not possible for the position metrics.

**Sort influence.** The sorting of nodes is essential for this graph to get some meaningful results. In fact it can be used for the detection of outliers regarding the height metric, if the nodes are sorted according to the width metric. If the two metrics are in close relation we often get a "staircase effect" because the nodes tend to get equally bigger in width and height. If this is not the case, the staircase effect breaks and we'll be able to easily detect those cases.

**Pro et contra.** One major drawback is that the width of a node will not directly reflect its metric, because it's being distorted by the graph width mapping function. Another drawback is that if the summed undistorted node widths of all nodes is bigger than the desired graph width, the nodes are shrunk in their width (otherwise they will be enlarged). If this shrinking is heavy, many small nodes will somehow disappear because they get very narrow, often only one pixel wide. The pro is obviously the intuitive detection of abnormal nodes which *don't* have to be outliers, but which stand out because two normally related metrics are not closely related in their case. Another pro is also that the graph will always fit the screen.

#### 2.2.9.6 The Confrontation Graph

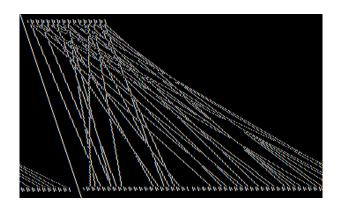


Figure 2.14: A confrontation graph using an horizontal layout

**Overall Idea.** This graph grew out of the necessity to display the access relationships between methods and attributes. An access is the only type of relationship between two entities of a different type.

**Scope.** This graph can only be applied on methods and attributes at the same time with accesses as edges. It's best used with the methods and attributes of one class.

**Layouts.** There are two possible displays. The first, called either *horizontal* or *vertical* displays on one row (column) the attributes and on the other one the methods. We can see such a layout in Figure 2.14. However, since in a class often the number of methods is much greater than the number of attributes, and the graph very soon gets larger than the screen, we introduced the *three row* layout. In this case the attributes are in the middle row, while the methods are in the upper and lower row.

Metrics. The size and color metrics can be used, while this is not possible for the position metrics.

**Sort influence.** A sort is advised for this graph. In the case of the method nodes it's especially useful according to the metrics LOC, NOS and NMAA. In case of the attribute nodes it's best to use NAA. If such a sort is applied, the number of edge crossings tends to drop and makes the graph look less cluttered.

**Pro et contra.** The major contra for this graph is that there is no special ordering of the nodes like clustering, except for a possible sort. However, it's the best graph to look at the internals of a class.

# 2.2.10 Useful Graphs: Class Graphs

In this section we list all graphs which display class nodes. We have noticed that the following graphs can be separated in two distinct groups. The graphs in the second group are normally applied after those in the first group, because they address more precise issues. We distinguish the following groups:

- 1. Those which serve primarily for system understanding. They work at a higher abstraction level, and in some cases can only return a general statement about the system. Problem detection is secondary in such graphs and in some cases not even possible. The following graphs fall under this category:
  - SYSTEM COMPLEXITY, Section 2.2.10.1.
  - SYSTEM HOT SPOTS, Section 2.2.10.2.
  - WEIGHT DISTRIBUTION, Section 2.2.10.3.
  - ROOT CLASS DETECTION, Section 2.2.10.4.
- 2. Those which primarily address problem detection, and secondarily program understanding. They must be applied on subsystems, rather than full systems. We list the following:
  - SERVICE CLASS DETECTION, Section 2.2.10.5.
  - COHESION OVERVIEW, Section 2.2.10.6.
  - SPINOFF HIERARCHY, Section 2.2.10.7.
  - INHERITANCE IMPACT, Section 2.2.10.8.
  - INTERMEDIATE ABSTRACT, Section 2.2.10.9.

## 2.2.10.1 System Complexity

Graph	Inheritance tree, without sort.	
Scope	Full system.	
Metrics		
Size	NIV (number of instance variables)	NOM (number of methods)
Color	WLOC (lines of code)	
Position	-	-

**General Idea:** This is one of the first graphs that should be applied to a system. It is an overview of the inheritance hierarchies of a whole system. This graph can give clues on the complexity and structure of the system (how many classes are present?), as well as information on the use of inheritance in the system (how deep do the hierarchies go and is the system in general flat or deep?). If we furthermore apply some class complexity metrics we can extract some more information. In this case we use as size metrics NIV and NOM, while for the color we choose WLOC. The detection of aberrant classes is now made easy: we can see if there are *very large classes, small classes* or even *empty classes*.

**Results with the Refactoring Browser:** In Figure 2.15 we see the SYSTEM COMPLEXITY graph applied on the Refactoring Browser. It shows few stand-alone classes and a few deep hierarchies. The first thing that strikes the eye is the class *BrowserNavigator* (A) which has a huge number of methods (175) and lines of code (1495) compared to the other classes present in the system. At the same time it only has one instance variable (this is the reason for its very narrow look). It may be a case for refactoring. If we take a look at the inheritance tree on the right side we can spot the class *BRStatementNode* (B) which is completely empty. When I asked the developers of the Refactoring Browser about this case, they told me that they were aware of the problem and that this class had been created to duplicate a hierarchy of another program. The same case can be spotted on one of the stand-alone classes *RefactoringError* (D) which is also empty. The next point of interest is the class *BRScanner* (C) which has the most instance variables (14) while it implements comparatively few methods (52). Perhaps this massive stand-alone class could be split up into subclasses. Another thing we can see is, that in the inheritance hierarchy in the middle of the graph, the root class *Refactoring* (E) is implementing by far the most methods, while there are quite a few very small classes deeper down the inheritance chain.

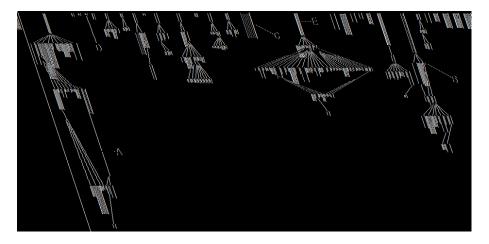


Figure 2.15: The system complexity graph applied on the Refactoring Browser using as size metrics NIV and NOM, and as color metric WLOC.

**Results with Duploc:** When we apply the SYSTEM COMPLEXITY graph on Duploc, we can spot the following in Figure 2.16: The system shows some very flat inheritance hierarchies, with many stand-alone classes which can have considerable sizes. This could mean that the system has not yet been refactored. There are three deep hierarchies, although in all three we can see that the main work is being done by the roots, which indicates top-heavy hierarchies. We also see that the main class *called DuplocApplication* (A) is very large and has only one very small subclass called it DuplocInformationMural<sup>7</sup>. Although *Duplo-cApplication* has the most methods and has the second most instance variables, the class with the most lines of code is *FastSparseCMatrix* (B). This class has only half the number of methods of *DuplocApplication* (74 vs. 130) but has nearly twice as much lines of code (1641 vs. 1060). Because of this we can already deduce that *FastSparseCMatrix* has some very long methods. The third point of interest are the classes on the left side (C): all of them are empty. These classes have become empty after being exported from the ENVY environment. The fourth eye-catch is the class *BinValueColoringModel* (D) on the right side. This class has the most instance variables (20), but only 52 methods. This may indicate that it is a service class which implements a lot of accessor methods. This supposition is being enforced by the light color value which is a sign for few lines of code (402), and is confirmed when we browse the source code of this class.

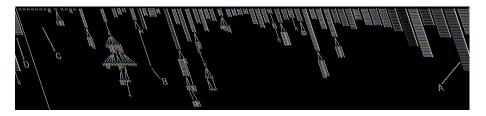


Figure 2.16: The system complexity graph applied on Duploc using as size metrics NIV and NOM, and as color metric WLOC.

**Possible Alternatives:** The color metric can be varied at will, especially class complexity metrics like NCV (number of class variables) prove to be useful.

**Evaluation:** This is certainly one of the first graphs that should be applied to a system, as it can return information on the structure and complexity of the subject system. However, it suffers one small drawback, which shows in very large systems: Sometimes the number of classes we want to display is so large that this graph takes several screens of place. It is difficult then to discern the outliers in the systems at one glance. The system hot spots graph discussed in Section 2.2.10.2 can counter this problem.

<sup>&</sup>lt;sup>7</sup>The InformationMural is a subapplication of Duploc included in a latter phase of development. Evidently the developer did not want to write an own main application class from scratch, but preferred to take the existing one, subclass it and override only some needed methods. This explains the small size of this class.

#### 2.2.10.2 System Hot spots

Graph	Checker, quadratic, sort according to width metric.	
Scope	Full system.	
Metrics		
Size	NOM (number of methods)	NIV (number of instance vari- ables)
Color	WLOC (lines of code)	
Position	-	-

**General Idea:** For very large systems it's hard to decide where to start looking for problems hot spots. One general rule is to look for very large or complex classes regarding their number of attributes and methods. The graph described here is a very simple display of all classes in the system sorted according to a certain metric. The nodes are placed next to each other to prevent overlapping. This graph detects outliers very easily because of the sorting. We distinguish the following:

- Large nodes at the bottom of the graph. These represent the biggest classes in the system.
- Small nodes at the top of the graph. These are the smallest classes which can sometimes even be empty.
- Very flat nodes. These nodes possess very few (if any) instance variables.
- Rather high nodes. This is seldom the case, as classes rarely have many attributes. Sometimes we can detect configuration classes like this.

**Results with the Refactoring Browser:** In Figure 2.17 we get a HOT SPOTS view on the Refactoring Browser. While in Figure 2.15 we had to search for the biggest and smallest nodes, this is made easy in this kind of graph because the nodes have been sorted: as before we can locate the class *BrowserNavigator* (A) and *BRScanner* (B). The sorting of the nodes makes it easy now to detect empty or very small classes, which find themselves at the top of the graph (D). Our attention is now also drawn to other classes like *BrowserApplicationModel* (C), which implements 38 methods while it defines no instance variable, which is visible by its flat shape. The view on the shape of the nodes is also facilitated, we can now detect classes like *MoveVariableDefinitionRefactoring* (E), which defines 6 instance variables while it implements only 7 methods (mainly accessors), giving it nearly a square shape.



Figure 2.17: The system hot spots graph applied on the Refactoring Browser using as size metrics NOM and NIV, and as color metric WLOC. The nodes have been sorted according to NOM.

**Results with Duploc:** The HOT POTS view on Duploc reveals also some information which could not be seen at first sight in Figure 2.16, as we see in Figure 2.18. We see Duploc has either very large classes (A)(B), or very small ones (D). We can also locate some classes with many instance variables (C). Two classes which could be interesting for further investigation because of are *DuplocCodeReader* (F) (32 methods, 17 instance variables) and *DuplocProgressMeter* (E) (15 methods, 9 instance variables): both classes have many instance variables and few methods, which could indicate service classes apt for refactoring.

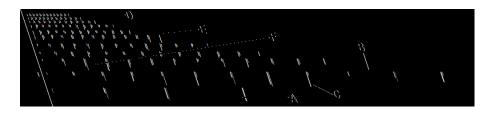


Figure 2.18: The system hot spots graph applied on Duploc using as size metrics NOM and NIV, and as color metric WLOC. Sort according to NOM.

**Possible Alternatives:** The color metric can be varied at will. A sort according to other metrics (especially WLOC and NCV) can also give interesting results which emphasise certain nodes.

**Evaluation:** The main drawback of the SYSTEM COMPLEXITY graph described in Section 2.2.10.1 is the fact that through the ordering of the nodes in tree structures we lose track of the size of the nodes all too easily. Only extreme cases strike our eyes. The SYSTEM HOT SPOTS graph described here makes this up through the sorting of the nodes and an ordering of them which reflects this sorting. However we lose the notion of inheritance in this case, since displaying the edges would mess up the view. A certain disadvantage of this graph is that the more nodes we display the more space is needed.

#### 2.2.10.3 Weight Distribution

Graph	Histogram, size addition layout, sort according to width metric.		
Scope	Full system.		
Metrics	etrics		
Size	NOM (number of methods)	-	
Color	HNL (hierarchy nesting level)		
Position	-	NOM	



Figure 2.19: The weight distribution graph applied on the Refactoring Browser. As width and vertical position metric we use NOM, as color metric we use HNL.

**General Idea:** With this graph we are able to make a general assessment on the system we are investigating. The width and the vertical position of the nodes is reflected by NOM, the color represents their HNL. This means that the deeper down (in the graph) the class nodes are, the more methods these classes implement. A dark node on the other hand means that the class it represents has a deep hierarchy nesting level. The possible assessments we can now make are:

- The system is *top-heavy*. This means that the classes that implement the most functionality are high up in the inheritance hierarchies. Such a graph has big nodes (on the bottom of the graph) which have very light color values (because their HNL is small). Top-heavy systems suffer when it comes to subclassing and reusing, because their root classes do too much themselves.
- The system is *bottom-heavy*. The most functionality is implemented in classes deep down the inheritance hierarchies. Such a case displays dark, big nodes on the bottom of the graph. Bottom-heavy systems are sometimes the results of overzealous abstracting mechanisms.
- The system is *even*. This display looks somehow chaotic, because the dark and light nodes distribute themselves over the whole graph. This case balances the two cases described above.

**Results with the Refactoring Browser:** The Refactoring Browser is an evenly distributed system, as we see in Figure 2.19: It's not possible to locate a majority of the dark or the light nodes on a certain area of the graph, although we can see there are three big classes marked as (A) high up the hierarchy.



Figure 2.20: The weight distribution graph applied on Duploc. As width and vertical position metric we use NOM, as color metric we use HNL.

**Results with Duploc:** Duploc is clearly a top-heavy system, as we see in Figure 2.20: The dark nodes are all very small (small NOM) and thus located on the top region of the graph. The big classes on the bottom of the graph are all very light (high up in the hierarchy). The system is thus to be classified as top-heavy, which is mainly due to its young age: Duploc has not yet undergone a reengineering or refactoring. It should be analysed on whether it's possible to introduce a supplemental abstraction level high up in the hierarchy.

**Possible Alternatives:** The width metric can be varied, especially NIV (number of instance variables) can give some supplemental information on the complexity of the classes. The color metric can also be changed, especially WLOC (lines of code) shows a good behaviour.

**Evaluation:** This graph can make a general assessment about the system. Such an assessment may not be very useful and will most probably not involve a specific problem, but upon such statements about the subject system we can vary our approach. In fact, the more we know about the system before we dive into its details, the more precisely we can deploy the other graphs.

## 2.2.10.4 Root Class Detection

Graph	Correlation.		
Scope	Full system or very large subsystems.		
Metrics			
Size	*	*	
Color	*		
Position	WNOC (total number of chil- dren)	NOC (number of children)	

**General Idea:** In very large systems with many inheritance hierarchies it may be difficult to identify at once the classes which have the most impact on their subclasses. The impact of a class on its descendants can be measured with the number of direct subclasses and the total number of subclasses of a class: the more there are, the more the functionality implemented in a root class is used. This graph shows the correlation between WNOC (total number of subclasses) and NOC (number of direct subclasses).

The further away from the origin such a class node is, the bigger is its impact. The type of inheritance used for a class can also be identified with this graph:

- If a node is positioned on the right side of the graph, while holding a vertical position near the top, this means that while the underlying class has a great number of descendants its direct subclasses are few. This is often the case when directly below a root class a supplemental abstraction level of classes has been introduced.
- If the node finds itself on the 45 degrees axis (it can't be further left because WNOC is always at least equal to NOC) and far away from the top of the graph, this means that the underlying class has a lot of direct subclasses. This is what we call a *flying saucer hierarchy* because the inheritance tree of this class resembles one.
- If a class node is positioned exactly along the 45 degrees axis this means that all its subclasses don't have subclasses themselves, and thus are leaf node classes in an inheritance tree.

**Results with the Refactoring Browser:** To make the effect of this graph more visible, in Figure 2.21 we see on the top left the root class detection graph while on the bottom right we see a display of two major inheritance trees. We see the class *Refactoring* (A) which has 43 descendants and 5 direct subclasses as root of major inheritance tree on the right side of the correlation graph. The other root class, *BrowserApplicationModel* (B) can also be identified on the right side of the graph. Two classes, *MethodRefactoring* (C) and *VariableRefactoring* (D), which are the heads of minor flying saucer hierarchies (14 and 13 direct subclasses) can be identified near the 45 degrees axis.

**Results with Duploc:** The results of this graph are somewhat deceiving in the case of Duploc, as its inheritance hierarchies are very flat. We can detect however two root classes, namely *PresentationModel-ControllerState* (A) and *PMCS* (B). In Figure 2.22 we see where the detected root classes are located in one of the inheritance hierarchies of Duploc.

**Possible Alternatives:** We do not make use of the color and size metrics, which could add information to this graph.

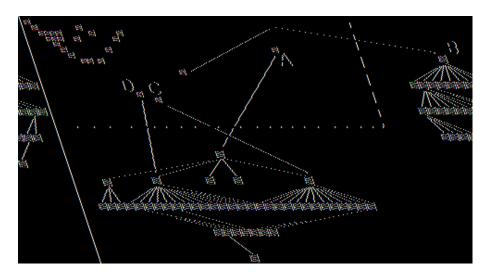


Figure 2.21: A root class detection graph applied on the Refactoring Browser. As position metrics we use WNOC and NOC.

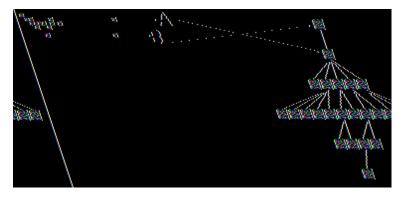


Figure 2.22: A root class detection graph applied on Duploc. As position metrics we use WNOC and NOC.

**Evaluation:** The detection of flying saucer hierarchies can of course be done through an inheritance tree display. The resulting tree graph has then to be searched for them. However, in some cases where the number of classes was very large, the resulting graph would become several screens big. In such cases it's not easy to detect flying saucers at once, and the graph described here comes into play. This graph can come in handy to see if there are some inheritance hierarchies upon which we want to apply inheritance specific graphs like intermediate abstract or inheritance impact.

### 2.2.10.5 Service Class Detection

Graph	Stapled, sort according to width metric.		
Scope	Subsystem or small full system.		
Metrics			
Size	NOM (number of methods)	WLOC (lines of code)	
Color	NOM		
Position	-	-	

**General Idea:** This graph has proven to be useful for the detection of so-called *service classes*. A service class is a class which mainly provides services to other classes. It often contains some tables and dictionaries which other classes can access for their purposes. Such classes often have an aberrant ratio between NOM and WLOC: they have very short methods which mainly access or return values. In this kind of graph we present a selection of some classes (a whole inheritance tree is often a good choice) as a stapled graph. The classes have been sorted according to their width, which represents NOM.

Because there tends to be a certain relation between NOM and WLOC, we should get a sort of staircase effect on the nodes the more we move to the right.

We can make out the following:

- If a class node breaks the staircase effect (by being too short) it is a candidate for a service class.
- This graph can also serve as detector for classes with overlong methods: If the class breaks the effect in the other direction (by being too tall) it's a candidate for method splitting, because this means that it has many lines of code (tall) and comparatively few methods (narrow, and because of the sorting pushed to the left side of the graph).

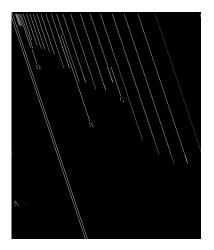


Figure 2.23: The service class detection graph applied on a subhierarchy of the Refactoring Browser. As width metric and sorting criterion we use NOM, the height metric is WLOC.

**Results with the Refactoring Browser:** In Figure 2.23 we selected a whole inheritance tree (26 classes) of the Refactoring Browser to be displayed in a SERVICE CLASS DETECTION graph. We see one huge class *BrowserNavigator* (A), which in fact is even bigger, but we cut it down because of space reasons. We see quite clearly that there is a certain tendency for a staircase which is severely broken in two places. The first

service class candidate is *CodeTool* (B), which has 22 methods and 49 lines of code. A closer inspection reveals that the methods are mainly get/set-methods (accessors). The second candidate is *CodeModel* (C) with 40 methods and 136 lines of code. The name itself already reveals the service function this class is intended to have. As method splitting candidate we detect the class *ClassCommentTool* (D) which has only 7 methods but 89 lines of code.



Figure 2.24: The service class detection graph applied on a subset of Duploc. As width metric and sorting criterion we use NOM, the height metric is WLOC.

**Results with Duploc:** We obtained the graph in Figure 2.24 by first applying the graph on the whole system and then by selecting a subset which looked interesting. We see there are some candidates for service classes: The class *CachedObservationData* (A) contains 20 methods for a total count of 50 lines of code. A closer inspection reveals it is truly a service class. Nearly the same ratio is visible in the classes *ComparisonMatrixBody* (B) (22/80), *PresentationModelControllerState* (C) (25/87) and *ObservationOnRawSubMatrix* (D) (30/122). Some classes tend to have overlong methods, namely *PMVSInformationMuralMode* (E) (22/396) and *DuplocCodeReader* (F) (32/530), and should be looked at for possible method splitting.

**Possible Alternatives:** Nearly the same results can be obtained if we use NIV (number of instance variables) instead of NOM: both NOM and NIV are closely related in service classes (because of the accessors). Sometimes abstract classes higher up the hierarchy tend to have the same properties as service classes, because their abstract nature makes them have several very short methods which are later overridden or extended by the subclasses. This can be alleviated if we use HNL (hierarchy nesting level) as color metric. Service class candidates which are true service classes tend then to have a darker color shade. Fake service classes like the abstract ones will have a lighter color shade because they are higher up the hierarchy.

**Evaluation:** As this graph addresses a special problem, it should be used in a second phase of reverse engineering. Experience has shown that it's advisable to apply it on subsystems, especially inheritance hierarchies.

## 2.2.10.6 Cohesion Overview

Graph	Checker, quadratic, sort according to width metric.		
Scope	Full system or subsystem.		
Metrics	S S		
Size	NOM (number of methods)	WNAA (number of direct ac- cesses on attributes)	
Color	NIV (number of instance variables)		
Position			

**General Idea:** In this graph the nodes differ greatly in shape and color. In the best case this graph can give us some clues on which classes we should inspect for a possible splitting. We distinguish the following:

- The flat nodes indicate that the methods of a class (the width indicates the number) do not access many times its instance variables. This is further emphasised by the small height (few instance variable accesses).
- The narrow and high nodes on the other hand, tend also to be very light colored. This case happens when the classes have many accesses but only few instance variables. This is mostly the case when the class defines an attribute which is then heavily accessed directly by its subclasses. This is not advisable because of the lacking encapsulation: a single access through an accessor which would then be invoked by other classes, instead of direct accesses on the attribute, would be much better.
- More or less rectangular nodes with darker color shades indicate a good cohesion inside those classes, although this is only provable after applying a class cohesion graph, which is described in Section 2.2.13.1.

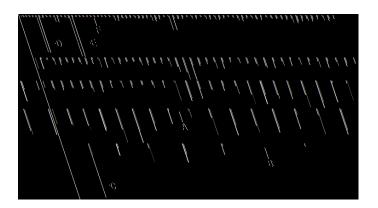


Figure 2.25: A cohesion overview graph applied on the Refactoring Browser. As size metrics we use NOM and WNAA. As color metric NIV is used.

**Results with the Refactoring Browser:** The resulting graph can be seen in Figure 2.25. The first thing we notice is that the nodes differ heavily in their shapes and colors. There are some white nodes that don't define instance variables (for example (A)) and because of this absence they can't have any instance variable access either. This is the reason for their flat shape. We also gather there are some empty or nearly empty ones (located around (F)). The class *BrowserNavigator* strikes once again the eye for its huge number of methods and its small number of instance variables (only one). The nodes (D) and (E) strike the

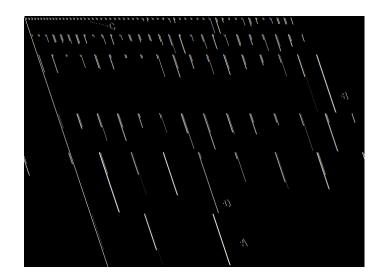


Figure 2.26: A cohesion overview graph applied on Duploc. As size metrics we use NOM and WNAA. As color metric NIV is used.

eye for their narrow shape and light color: Both have few methods and instance variables, (1,2) and (2,1) respectively, while at the same time they have a huge number of accesses. The reason for this is that their variables are directly accessed by their subclasses. The class *BRScanner* (C) shows a great complexity and heavy access.

**Results with Duploc:** The graph in Figure 2.26 shows a few characteristics of Duploc: Many empty or nearly-empty classes (C), quite a few heavy-access classes (B) and (D) and a few very large classes, for example *DuplocApplication* (A). We see there are quite a few classes that could be interesting for inspection with a class cohesion graph and do that for one special case, the class *DuplocApplication* in Section 2.2.13.1.

Possible Alternatives: None.

**Evaluation:** This graph can be seen as an *in-betweener*, because it comes after a graph for general overview and before a graph which treats class internals. The best result it can return is a collection of classes which we should further examine with a class cohesion graph, described in Section 2.2.13.1.

#### 2.2.10.7 Spinoff Hierarchy

Graph	Inheritance tree, centered, without sort.		
Scope	Subsystem, especially inheritance hierarchies.		
Metrics	3		
Size	WNOC (total number of chil- dren)	NOM (number of methods)	
Color	WNOC (total number of children)		
Position	-	-	

**General Idea:** We have noticed that in inheritance hierarchies the notion of inheritance is often carried on only by one or two classes on each level of the inheritance tree. This means that when a class has some subclasses often only one of them is really carrying on the weight of the inheritance, while its siblings tend to be *spinoff* classes implementing only few functionalities. Although this is not a bad thing per se, an easy detection of such spinoff hierarchies could make us focus on the inheritance carriers, while we could save time by ignoring (at least at the beginning) the less important spinoff classes. Spinoff classes often implement few methods and have few or no subclasses at all.

We distinguish the following:

- Small, light colored nodes. These are the *spinoff classes* with few or no children and few methods.
- Large, dark colored nodes. These are the *inheritance carriers*.

**Results with the Refactoring Browser:** In Figure 2.27 we see all inheritance hierarchies that make up the Refactoring Browser. We filtered out all stand-alone classes to get a clearer overview. We detect two cases of spinoff hierarchies:

- The one with the class *BrowserApplicationModel* (A) as root. We see two classes split up the second level of this tree, namely *CodeTool* (A21) and *Navigator* (A11). There are a few spinoff classes on this level, neither of them has subclasses. The same situation is present on the next level of this tree where the classes *BrowserTextTool* (A22) and *BrowserNavigator* (A12) carry on the weight of inheritance. A good example for spinoffs is visible between *CodeTool* (A21) and *BrowserTextTool* (A22): *CodeTool* has 7 subclasses but only one of them, *BrowserTextTool*, carries on the inheritance. Each one of its siblings is very small (keep in mind that the height reflects NOM) and is thus a spinoff.
- 2. The one with the class *Refactoring* (B) as root. Again two main inheritance threads are visible: The one consisting of *Refactoring* (B), *MethodRefactoring* (B11) and *ChangeMethodNameRefactoring* (B12). The other consists of *Refactoring* (B), *VariableRefactoring* (B21) and *RestoringVariableRefactoring* (B22).

The other inheritance trees in this display also show some property of a spinoff hierarchy, and could be a case of further investigation.

**Results with Duploc:** After removing the many stand-alone classes from Duploc, the remaining graph in Figure 2.28 can only show us the absence of spinoff hierarchies. Especially in the tree with the class *PresentationModelControllerState* (A) as root, we see that on the third level we have 5 siblings, 4 of which are all inheritance carriers, with only one tiny spinoff class with the meaningful name *PMCSDummyMode* (B).

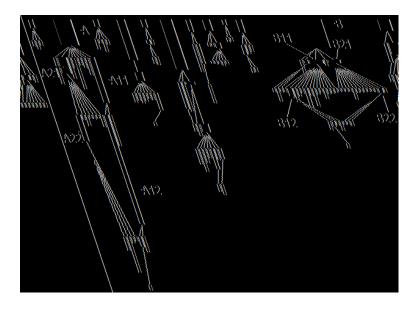


Figure 2.27: The spinoff hierarchy graph applied on the inheritance hierarchies of the Refactoring Browser. As size metrics we use WNOC and NOM, as color metric WNOC.



Figure 2.28: The spinoff hierarchy graph applied on Duploc. As size metrics we use WNOC and NOM, as color metric WNOC.

**Possible Alternatives:** We have to emphasise that a preprocessing consisting of filtering out all standalone nodes is advised for this graph, as they add unnecessary complexity to the displayed graph. This graph does not have real alternatives, as it addresses a special problem.

**Evaluation:** This graph should come into play in a later phase of the reverse engineering, as it addresses a special problem which may not be present at all in the system. The detection of an inheritance carrier could be important, as it is the place which should be checked out because subclasses depend on it. The spinoff classes on the other hand, can be examined for possible push-ups of functionality.

### 2.2.10.8 Inheritance Impact

Graph	Inheritance tree, without sort.		
Scope	Subsystem, especially inheritance hierarchies.		
Metrics	s i i i i i i i i i i i i i i i i i i i		
Size	NMO (number of methods over- ridden)NME (number of methods ex- tended)		
Color	NOM (number of methods)		
Position	-	-	

**General Idea:** This graph is able to tell us if there has been made an improper or suspect use of inheritance: it can tell us if a class that implements many methods does not make use of method overriding or method extension, or uses it only rarely. Overriding and extending methods is one of the powerful properties of object-oriented languages and should be used if possible.

Nodes that override or extend a lot are bigger, nodes that implement many methods are dark. We are looking for dark nodes (many methods) which are at the same time very small (no use or rare use of overriding and extension).

**Results with the Refactoring Browser:** One of the hierarchies of the Refactoring Browser seems to have one such class which should certainly be further investigated: In Figure 2.29 we can detect the class *BrowserNavigator* (A) which implements many methods (175), while it only overrides one and extends two methods.



Figure 2.29: The inheritance impact graph applied on an inheritance tree of the Refactoring Browser. As size metrics we use NMO and NME, as color metric NOM.

**Results with Duploc:** This graph returns no meaningful results if applied on Duploc.

**Possible Alternatives:** No real alternatives, as it addresses a specific problem. This graph is often obtained after filtering out all stand-alone classes and all inheritance hierarchies which show no sign we are looking for.

**Evaluation:** A graph which addresses a very special problem. It's not always useful, but if it can detect something, it can be an important discovery which can affect a whole inheritance hierarchy.

Graph	Inheritance tree, without sort.	
Scope	Subsystem, especially inheritance hierarchies.	
Metrics		
Size	NOM (number of methods)	NMA (number of methods added)
Color	NOC (number of children)	
Position	-	-

### 2.2.10.9 Intermediate Abstract Class

**General Idea:** This graph is useful for the detection of abstract classes or nearly-empty classes which are located somewhere in the middle of an inheritance chain. Often they tend to have a superclass which implements a lot of methods. The programmer then started to subclass this class. The number of direct subclasses would soon be too big so an attempt was made to logically group several subclasses under an abstract intermediate class.

Such an intermediate subclass would normally have many children, while at the same time its size is very small (because it is abstract or nearly empty). We thus have to look for small, dark nodes in the middle of inheritance hierarchies.

The dark color comes from the greater number of direct subclasses, while the small size from the small functionality implemented. We chose NMA as height metric to reflect the fact that often such intermediate abstract classes don't override superclass methods, which in turn means that is we use NOM as width metric, the node is square (no functionality implemented, or if there is a bit of implemented functionality, then it doesn't come from the superclass). Intermediate abstract classes are of some interest, because often we can try to push up some functionalities of its subclasses into it, thus concentrating them in one place, instead of spreading the functionality all over the subclasses, risking to obtain duplicated code.

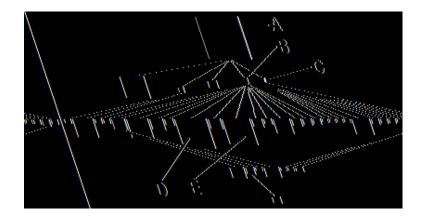


Figure 2.30: The intermediate abstract class graph applied on a subset of the Refactoring Browser. As size metrics we use NOM and NMA, as color metric NOC.

**Results with the Refactoring Browser:** The Refactoring Browser harbours in one of its inheritance hierarchies two intermediate abstract classes, as we see in Figure 2.30. The root class *Refactoring* (A) implements quite a few methods, while we can spot the two intermediate abstract classes as *MethodRefactoring* (B) and *VariableRefactoring* (C). These two classes implement themselves very few methods (2 and 1 respectively) and are the roots of smaller subhierarchies. In the case of *MethodRefactoring* we see that its

subclasses are implementing several methods, as we see in *InlineMethodRefactoring* (D) and *MoveMetho-dRefactoring* (E). Perhaps an attempt could be made to extract duplicated code and push it up into the intermediate abstract class.



Figure 2.31: The intermediate abstract class graph applied on an inheritance hierarchy of Duploc. As size metrics we use NOM and NMA, as color metric NOC.

**Results with Duploc:** One of Duploc's inheritance hierarchies also contains an intermediate abstract class, as we see in Figure 2.31: The subclass *PMCS* (B) of the root class *PresentationModelControllerState* (A) implements only 4 methods and is obviously an intermediate abstract class. The subclasses of *PMCS* should be searched for duplicated code which could be pushed up into *PMCS*.

Possible Alternatives: None.

**Evaluation:** The detection of abstract classes is very important: several object oriented languages either directly provide a declaration or support a standard idiom for identifying abstract classes. Abstract or nearly abstract classes can be seen as the hinges of the system, upon which several classes depend. It's where the common functionality is defined and where we should start to look at source code if we want to understand the logic of their subclasses.

# 2.2.11 Useful Graphs: Method Graphs

Method graphs can work at any level of granularity most of the time. However, the more method nodes we display, the harder it is to make out outliers. Methods are the entities which are responsible for the action in a system. This implies that every graph which uses method nodes is often followed by an examination of the actual underlying source code. This means that the graphs listed here have a very concrete context.

In this section we list the following graphs:

- METHOD EFFICIENCY CORRELATION, Section 2.2.11.1.
- CODING IMPACT HISTOGRAM, Section 2.2.11.2.
- METHOD SIZE NESTING LEVEL, Section 2.2.11.3.

Graph	Correlation.		
Scope	Full system, subsystem or single class.		
Metrics			
Size	NOP (number of parameters)	NOP	
Color	*		
Position	LOC (lines of code)	NOS (number of statements)	

#### 2.2.11.1 Method Efficiency Correlation

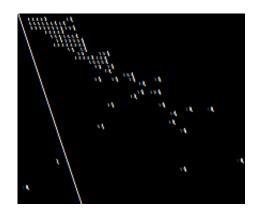


Figure 2.32: A method efficiency correlation graph.

**General Idea:** This graph is a good way to locate the *freaky entities* inside a group of methods, when it comes to their efficiency. By efficiency we mean how many statements are put on each line. By displaying the nodes in the correlation graph (as in Figure 2.32), we see that most of the nodes are near a certain correlation axis. However, there are a few which do not adhere to this rule.

The methods that are not near the correlation axis may have some problems, which may be

- 1. High LOC (lines of code) and low NOS (number of statements). This is for example the case with "forgotten methods", that at some point have been commented out and then been forgotten. This may also be the case for overzealous line indentation, when a single parenthesis is put on a line of its own or when many blank lines have been used.
- Low LOC and high NOS. This can be the case when the methods are written without indentation and several statements are on the same line, which is a bad thing too, since this decreases the readability, and it may also break the law of Demeter [LH89].
- 3. Long methods (high LOC and high NOS). Normally a case for redesign, since long methods should be split up in smaller, better understandable and reusable ones [Bec97].
- 4. Empty methods. These nodes position themselves on the top left of the graph. Although they can be viewed there by selecting and moving, the overlapping of the nodes which is characteristic for this graph makes it hard to see those empty methods at one glance. A better graph for the detection of empty methods is the Coding Impact Histogram described in Section 2.2.11.2.

Other hot spots can be detected by looking at the size of the nodes:

• Big nodes have many parameters. Although it's hard to define a threshold on the number of parameters, we think that methods taking more than 5 parameters should be looked at.

• Very small nodes on the outskirts of the graph should be looked at: these are very long methods which do not take any input parameter. Perhaps they could be split up easily.

The interesting property of this graph is its scalability. Since most of the nodes overlay each other, and those nodes are of no real interest to us, because they have average metric measurements, we can display several thousands of nodes at the same time. Our interest is drawn by the nodes which find themselves on the outer skirts of the graph, and which do not suffer overlaying, as their position is defined by their non-average metric measurements. The size of this graph is not affected by the number of displayed nodes, but on the maximum values for the position metrics.

**Results with the Refactoring Browser:** The method efficiency correlation graph shows some interesting results when applied to the Refactoring Browser. In Figure 2.33 we display all 2365 methods of the Refactoring Browser. We can spot several cases which should be looked into. The first nodes to meet the eye are those on the right edge of the graph (A). These three methods are very long (45, 51 and 65 lines of code) compared to the other methods in the system, which does not have a great distribution, thus signifying that the system is homogeneous related to the method lengths. The opposite case can be seen on the top left side of the graph (B). Upon closer inspection (by selecting and moving the nodes) we can see that the RefactoringBrowser contains 20 empty methods. The next point of interest is the method marked (C): this method takes 7 input parameters which is of course very much. The method *reInstallInterface* (D) on the top of the graph is also a case of closer study: While it has 16 lines of code it contains no statements. If we browse its source code, we see that the whole body of the method has been commented. The method *needsParenthesisFor:* (E) on the other hand contains 31 statements in only 19 lines of code and should perhaps be reformatted. The group of methods marked as (F) should also be looked into, since all of them contain comparatively few statements in long method bodies.



Figure 2.33: The method efficiency graph applied on the Refactoring Browser, using as position metrics LOC and NOS, as color metric HNL, and as size metric NOP.

**Results with Duploc:** When this graph is applied to Duploc, as we see in Figure 2.34, the first thing to strike the eye is the large distribution of the nodes. Duploc obviously does have some very long methods. The second thing that meets the eye is that the main correlation axis has a different angle compared to the Refactoring Browser in Figure 2.33. The method *putPerlCode:* (A) is 201 lines long but does have only 2 statements. Upon closer inspection we see that its purpose is to print out a very long string. We have some other very long methods, (B) with 135 lines, (C) with 95 lines, (D) with 109 lines. We have some method that are far away from the system correlation axis, like (A), (C), (E), (F) and (G). (E) for example has 64 lines of code with only 1 statement. A closer inspection reveals its body is mainly commented code for testing purposes, i.e. when the system is tested some parts of the method body are uncommented. (F) reveals the same situation, where the 18 lines long method body doesn't contain any statements at all. (G) has 32 statements packed in 14 lines of code. Reformatting makes it more readable. The empty methods

can of course be detected as (H), while we should also note the nodes around (I), which seem to be very short and at the same time badly formatted methods. The two methods (J) also draw attention due to their considerable size, which reflects the fact that they take 9 input parameters each.

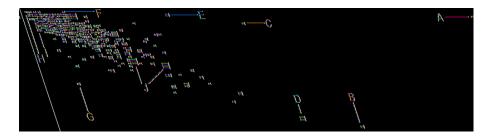


Figure 2.34: The method efficiency graph applied on Duploc, using as position metrics LOC and NOS, as color metric HNL, and as size metric NOP.

**Possible Alternatives:** We chose the size of the nodes to be represented by NOP (number of parameters). Since the distribution tends to get sparse the more we move to the right and to the bottom, we can see the methods which take many parameters more clearly, since it's normally the large methods that take more parameters. Generally in this graph the size metrics can be chosen freely, although it's advisable to use metrics which tend to have small measurements. Otherwise the nodes get very big and clutter up the view. The color metric can also be used freely. We chose HNL (hierarchy nesting level) in this case, but since the nodes in this graph tend to be very small, the color node metric doesn't really matter.

**Evaluation:** This is one of the few graphs which works very well at any level of granularity. As such it can be used anytime. We saw it can be useful to apply it on a subsystem before we dive into its details. At class level it can help to detect problem cases for a concrete reengineering.

Graph	Histogram, size addition layout, sort according to width metric.		
Scope	Single class or small subsystem.		
Metrics			
Size	LOC (lines of code) -		
Color	LOC		
Position	LOC -		

#### 2.2.11.2 Coding Impact Histogram

**General Idea:** This graph shows the coding impact of methods and where the most coding has happened. While the normal histogram can only tell us how methods are distributed regarding their lines of code, this graph (Figure 2.35) can reveal where the real programming effort has been made: Writing 20 methods each one line long is easier than writing one method 20 lines of code long. It shows if there are any aberrant methods that are too long or if the system is unbalanced because of too long and complex methods. As a nice side-effect we can also grasp at one glance if there are any empty methods (those at the very top of the graph). A good design should have a lot of tiny methods so this is where the biggest columns in the graph should be. Methods not following this rule should be analysed as possible "split candidates" which could be broken down into smaller pieces. While this graph is inefficient on whole systems because of the huge number of methods, it has proven to be very useful when applied to the methods of one single class. It should also be noted that the average length of a method implemented in typical industrial Smalltalk applications is around 6 lines [Bec97].



Figure 2.35: A coding impact histogram.

**Results with the Refactoring Browser:** Since this is one of the graphs which can hardly be applied on whole systems, but rather on specific small subsystems or singular classes, we do not compare the systems from our case studies with each other, but we rather show a few illustrative examples taken out randomly<sup>8</sup> from the Refactoring Browser. We selected only the two classes (*BrowserNavigator* (B) and *BRProgramNode* (A)) with the most methods for this graph. We see in Figure 2.36 that each class has its own coding impact topography. We see that *BrowserNavigator* (B) has many methods which tend to be overlong, and especially 6 very long ones which isolate themselves (B1) from the others. On the other hand *BRProgramNode* has an irregular topography with many accessors (A2) and one very long method (A1).

**Possible Alternatives:** This graph knows many useful mutations, especially those which keep LOC as vertical position metric, but use other size and color metrics and a different sort criterion. In these cases, especially NI (number of invocation) and NMAA (number of accesses on attributes) showed good behaviour.

<sup>&</sup>lt;sup>8</sup>This randomness should also express the interactive approach of such systems, which is guided by intuition rather than a systematic methodology, although experience has shown that at the beginning of a reverse engineering experiment we tend to apply a certain fixed set of graphs. This reflects the fact that the graphs address each a different level of abstraction.

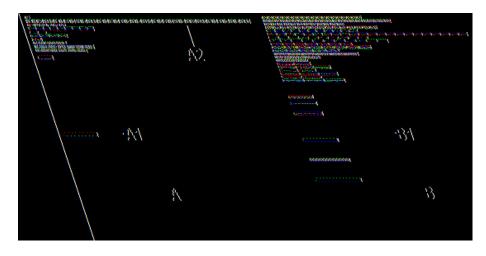


Figure 2.36: The coding impact graph applied on two classes of the Refactoring Browser. The width metric, as well as the color and vertical position metric is LOC.

**Evaluation:** This graph is very useful to *get a feeling* for certain classes or subsystems. It can show us what kind of implementation lies behind the subject entities and in certain cases what we should continue to explore.

Graph	Checker, quadratic, sort according to width metric.		
Scope	Subsystem, especially inheritance hierarchy. No stand-alone classes.		
Metrics			
Size	LOC (lines of code)	NOS (number of statements)	
Color	MHNL (hierarchy nesting level)		
Position	-	-	

#### 2.2.11.3 Method Size Nesting Level

**General Idea:** A general rule is that big methods should be split up [Bec97] into smaller chunks to increase their reusability and to make them easier to understand. This is especially true for methods that are implemented in classes deep down the inheritance hierarchy: perhaps parts of those big methods could be extracted and put up into a higher class to reuse them across several subclasses. The method size nesting level graph can help us to detect large methods deep down the inheritance hierarchy: It's a checker graph of methods with LOC and NOS as size metrics and MHNL as color metric. The nodes are sorted according to LOC, which puts the larger methods on the bottom area of the graph.

Since the color reflects the MHNL of the methods, we should be looking for big, dark nodes in the bottom area of the graph: these are possible split candidates. We call such methods split-and-push-up candidates.

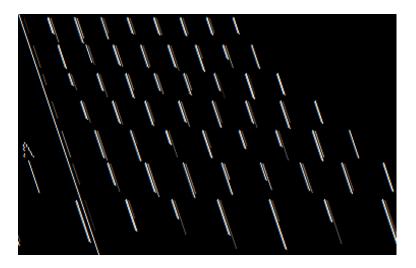


Figure 2.37: The method size nesting level graph applied on the largest Refactoring Browser methods. Size metrics: LOC, NOS. Color metric: MHNL.

**Results with the Refactoring Browser:** The Refactoring Browser shows in Figure 2.37 that is has been refactored itself a few times: there remain very few large methods, after filtering out all those with a LOC measurement smaller than 20. Yet, there are some large methods which also have medium MHNL values like those in the last row (A). Their lengths vary from 65 to 37 lines, which makes them also possible split-and-push-up candidates.

**Results with Duploc:** We display in Figure 2.38 only the methods that have more than 20 LOC and belong to non-stand-alone classes. The resulting graph shows us there are several very large methods, which on one hand don't have big MHNL values, but since they're not methods belonging to root classes either, are all the same split-and-push-up candidates. The biggest methods (A) have 201, 135 and 109

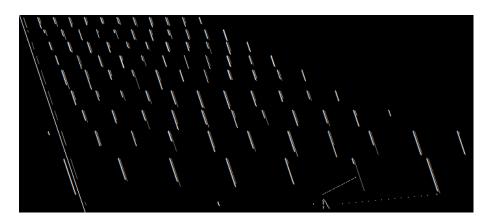


Figure 2.38: The method size nesting level graph applied on several Duploc methods. Size metrics: LOC, NOS. Color metric: MHNL.

LOC, which is way too much for Smalltalk methods. This excessive size is again due to the fact that most of them have never been refactored and written in one pull.

**Possible Alternatives:** The same graph using only LOC as size and color metric can be applied on whole systems (including stand-alone classes). In such a case the graph serves to easily detect very large methods which could be split up.

**Evaluation:** Since this graph is useful for classes belonging to inheritance hierarchies, it should primarily be used to get insights into such structures as to where the methods are which could be reengineering candidates.

## 2.2.12 Useful Graphs: Attribute Graphs

Attributes define the properties of classes. As such, it's mandatory that to understand the purpose of an attribute, we have to understand the class in which it is defined. This implies that very soon after applying one of the following graphs, we have to look at the source code of the class.

In this section we list the following graphs:

- DIRECT ATTRIBUTE ACCESS, Section 2.2.12.1.
- ATTRIBUTE PRIVACY, Section 2.2.12.2.

#### 2.2.12.1 Direct Attribute Access

Graph	Checker, quadratic, sort according to width metric.		
Scope	Full system or subsystem.		
Metrics			
Size	NAA (number of times accessed directly)	NAA	
Color	NAA		
Position	-	-	

**General Idea:** This is a graph of all attributes of a system or subsystem. As metrics we use NAA (number of times accessed) for the size and the color. We then also sort the nodes according to NAA. What we get is a clear display of which attributes are accessed the most in a system. These attribute nodes are positioned at the bottom of this graph. The largest nodes should be a case for closer inspection. The general rule should be that attributes which are accessed directly can break the system if the inner implementation of the attribute changes. This can be avoided by using an accessor method which returns the value(s) of the attribute. An accessor on such an attribute can provide a defensive wall of protection against such changes. There may also be some attributes which are never accessed and which may have been forgotten in the system and thus only add unnecessary complexity to it. They could be removed from the system. Such attribute nodes are positioned on top of the graph.

**Results with the Refactoring Browser:** In Figure 2.39 we notice at once that there is the attribute *class* (A) defined in the class *MethodRefactoring* which is directly accessed 86 times. We also see there are some never accessed attributes which should also be further investigated (B).



Figure 2.39: The direct attribute access graph applied on the Refactoring Browser. The size, color metric and sort criterion is NAA.

**Results with Duploc:** In Figure 2.40 we see that while in Duploc there are no attributes which are heavily accessed (the maximum is 31 direct accesses for the attribute *region* (A) defined in the class *AbstractRaw-SubMatrix*) there are many attributes which are never accessed (B) and which should be looked into for possible removal.

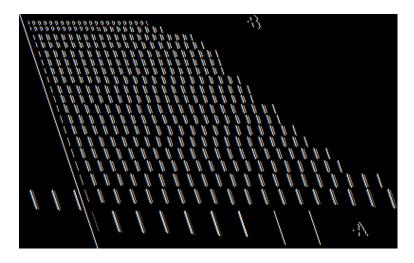


Figure 2.40: The direct attribute access graph applied on Duploc. The size, color metric and sort criterion is NAA.

**Possible Alternatives:** An interaction with interesting nodes is necessary to see if accessors have been implemented for them and if those accessor methods are used all the time.

**Evaluation:** A graph which works at every level of granularity. The next step which has to follow such a graph is to examine the classes in which the outlier attributes are defined. Note that this graph takes only the direct accesses into account. If an attribute is accessed very often through the use of an accessor method this will not show in this graph. Note that the quality of this graph depends heavily on the quality of the metamodel. Especially when building a model out of a CDIF file we have often seen that sometimes accesses are left out. This can lead us to wrong conclusion on never accessed attributes. Again, a check against the code has to be done to be sure.

## 2.2.12.2 Attribute Privacy

Graph	Checker, quadratic, sort according to width metric.		
Scope	Full system or subsystem. Better performance with C++ or Java.		
Metrics			
Size	NAA (number of times accessed directly)	NCM (number of classes which access this attribute)	
Color	*		
Position	-	-	

**General Idea:** Attributes may be directly accessed several times in a system. As we said in Section 2.2.12.1 such a situation is not ideal and can be detected with the graph described there. Apart from the number of times an attribute is accessed, another metric may prove to be useful for a similar graph: NCM, the number of classes which have methods that directly access a certain attribute. The attribute privacy graph is a checker graph which uses as size metrics NAA and NCM.

We are looking for wide, high nodes: such nodes are directly accessed a lot of times by many classes and should have an accessor at all costs, because the system easily breaks if such an attribute is tampered with.

Very wide but shallow nodes should also be looked at: although they are directly accessed a lot, it's by few or often only one class. If it's the case of only one accessing class, it should be checked it the attribute in question is private. If not, it can be made private without impact on the rest of the system.



Figure 2.41: The direct attribute access graph applied on the Refactoring Browser. The size metrics are NAA and NCM.

**Results with the Refactoring Browser:** In Figure 2.41 we can spot some heavily accessed attributes marked as (A) which are accessed by many classes. We also see there are some very flat but wide nodes which are attributes heavily accessed by only 1 or very few classes.

**Results with Duploc:** In Figure 2.42 we can see that as a difference to the Refactoring Browser, Duploc has attributes which are seldom accessed by more than one class. The maximum NCM value is 3. We deduce from that the implementor of Duploc keeps an eye on encapsulation<sup>9</sup>.

Possible Alternatives: None.

<sup>&</sup>lt;sup>9</sup>The implementor of Duploc used to implement a lot in C++, which could be a reason for the tight encapsulation.

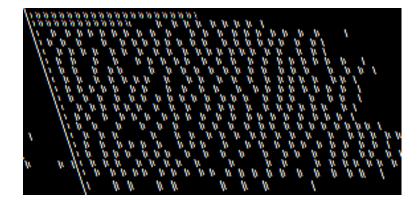


Figure 2.42: The direct attribute access graph applied on Duploc. The size metrics are NAA and NCM.

**Evaluation:** A graph whose purpose is to find attributes which have to be examined. Since such an examination takes place at textual level, it's a graph which can help find problems at once. The results are incomplete in this case: the last step after detecting wide and flat nodes would be to check if the attributes concerned are defined as private. If not they could be made private. However, this does not work in Smalltalk, so we had to leave that part out with our case studies.

# 2.2.13 Useful Graphs: Class Internal Graphs

A class internal graph treats the special case where the components of a class are displayed at the same time: methods and attributes.

In this case we find ourselves at a low level of abstraction, the source code is only one step away and it's necessary to look at it after applying a class internal graph.

In this section we list the following graph:

• CLASS COHESION, Section 2.2.13.1.

#### 2.2.13.1 Class Cohesion

Graph	Confrontation graph, nodes sorted according to their width metrics.			
Scope	Single class.			
Metrics (Method Nodes)				
Size	LOC (lines of code)	NOS (number of statements)		
Color	LOC			
Position	-	-		
Metrics (Attribute Nodes)				
Size	NAA (number of times accessed directly)	NAA		
Color	NAA			
Position	-	-		

**General Idea:** This graph is a confrontation graph where the edges represent instance variable accesses between methods and attributes. This graph can indicate us how strong the internal cohesion of a class is. If a class has many accesses and looks very chaotic, this means that the class is difficult to split. On the other hand, if we can make out two or more separate clusters in this display, this is an indication that the class is a good split candidate. If the root class of an inheritance hierarchy shows such characteristics it is a sign that the hierarchy tends to be top-heavy. If the class shows sparse attribute accesses it could be easier to subclass.

**Results with the Refactoring Browser:** In Figure 2.43 we displayed the methods and attributes of the class *BRScanner* which has been identified as (C) in Figure 2.25. We gather at once that this class is heavily coupled internally and that splitting such a class is next to impossible.

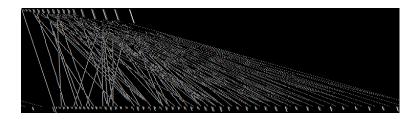


Figure 2.43: A class cohesion graph applied on the class BRScanner. The method nodes (in the lower row) use as size metric NOS and as color metric LOC. The attribute nodes (in the upper row) use as color and size metric NAA.

**Results with Duploc:** We obtained some impressive results when we applied this graph to some classes of Duploc. We show only one here: the class *DuplocApplication*. After filtering out all methods that never accessed attributes, we got the graph displayed in Figure  $2.44^{10}$ . We clearly see two distinct clusters of attribute and method nodes. This class is thus certainly a split candidate. This suspect was confirmed afterwards when I asked the implementor of Duploc about this class. He confirmed that this class was to be split up during the next redesign of the system.

<sup>&</sup>lt;sup>10</sup>Note that the graph resulted like this after direct manipulation of the graph (i.e. moving around nodes) and not because of a layout algorithm that can identify clusters. However, we included into CODECRAWLER the functionality to help us quickly identify such clusters.

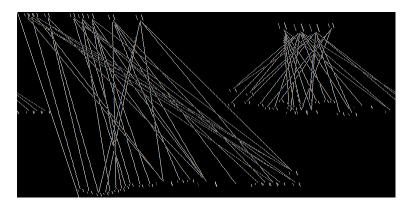


Figure 2.44: A class cohesion graph applied on the class DuplocApplication.

**Possible Alternatives:** We advise the user to remove all stand-alone nodes from the graph, as they are of no use in this case. The metrics, especially the color metric in the method nodes can be varied freely.

**Evaluation:** This graphs needs some interaction before it can express its full potential. However, its usefulness is indisputable: Up to this moment we haven't seen a technique which can detect split candidates with such an easy and quick method.

# 2.3 Grouping

**Author: Oliver Ciupke** 

## 2.3.1 The problem

Reengineering large software systems requires analysing and manipulating tasks at multiple levels of abstraction. The standard views of a system used during design and implementation, such as the class and method structure as well as the source code, are often not sufficiently abstract. These views provide a wealth of information, not all of which is relevant to the task of reengineering. Consequently, the essential information may be obscured.

Additionally, some problems arising during the design of software are themselves only understandable at a high level of abstraction. For example, the class structure may look well designed even though there are too many dependencies between the packages or subsystems. This can make the system unnecessarily complex thus compromising flexibility and increasing the cost to maintain the system.

Abstract views of a system are already supported in several design languages and methods. Examples include packages in UML [RJB99] or subjects in OOA [CY91]. Unfortunately, they are currently to vaguely defined making them insufficient for the tasks occurring in reverse- and reengineering. Here, a concept must be implementable by tools in a way that it is composable with other techniques (such as querying and problem-detection) in a clearly defined manner.

## 2.3.2 Grouping

To be able to analyse a system, we must have an exact way to come from elementary views to more abstract ones. They must be formally sound, so it is known with precision what the meaning of an abstraction and to ensure that it is possible to build tools able to deal with these abstractions.

What we need is a description of how to move from a detailed description of a system to a more abstract view. We call this process *grouping* (also referred to as *abstraction*, *compression* or *lifting*). Grouping means replacing a set of entities, often describing a common abstract concept, into one abstract entity called a *group* (or complex entity).

Figure 2.45 shows a model of a small C++ program including classes and methods together with their relationships. Figure 2.46 shows a more abstract view on the same program where methods have been grouped together with the classes they belong to. One can imagine, that these kinds of abstractions are often useful to get an appropriate overview and to understand overall dependencies. For larger systems, it is even necessary to go beyond this level, e.g., to view it on the level of subsystems or packages.

#### 2.3.2.1 Examples

In principle, any set of entities under consideration can be grouped together. Similarly, each entity can in principle be split up into its component parts in order to provide a more detailed view of the system. However, there are several groupings which are particularly useful in building up common abstractions. Examples for those are

**Classes to packages:** Grouping classes to packages (or subsystems or modules or clusters) is probably the application most commonly used. Figure 2.47 shows an example. The three upper left packages could be a framework, whereas the lower right package could be an instantiation of this framework towards a specific application. After grouping, dependencies between classes are propagated to the surrounding packages. The process leads to an overview over the overall architecture of the system.

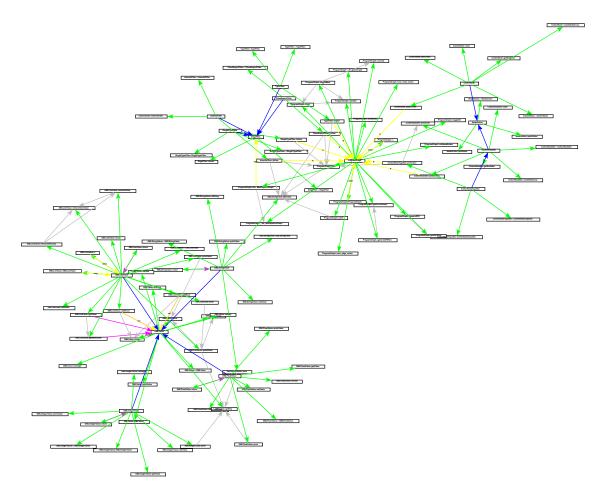


Figure 2.45: Program structure with classes and methods

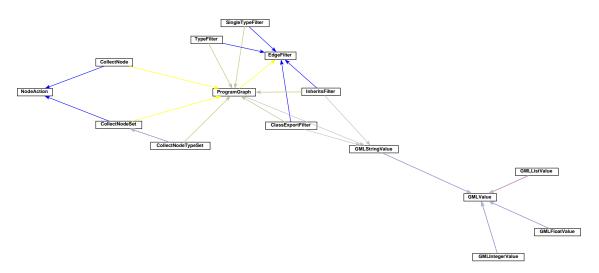


Figure 2.46: Program structure with collapsed methods

Interesting views in this example are as well those where not the classes of all packages (groups) are being replaced by a single node. We replace only those belonging to the framework itself, if we want to investigate how the application uses the framework. Or we could only group the instantiation if we want to see, which classes of the framework are actually needed.

- **Packages to packages:** In general packages may be grouped recursively giving a *grouping hierarchy*. In some programming environments, packages of different levels are given different names such as "subsystems", "program blocks", "service blocks" etc. Typically, this case occurs together with the former one (classes to packages).
- **Classes to super-classes:** Grouping classes with (possibly one of their) their super classes is an abstraction commonly applied when modelling a system or an algorithm. An algorithm working on viewable elements is described in terms of an abstract super-class only.

Since we must take into account all possible dependencies, it is not enough only to consider only the super class, again we have to propagate relationships of all original classes to the remaining representant.

- **Objects to classes (or types):** In a dynamic view, objects of a snapshot or a trace can be mapped to their classes. This is extremely useful due to the huge number of objects often being created during a program run. Therefore, similar functionality has so far often been implemented in tools such as visual debuggers (e.g., LOOK! <sup>11</sup> or profilers. Related work on visualising the behaviour of object-oriented systems are for example [PHKV93], [KLRZ94].
- **Dynamic method calls to pair of calling and called object:** A set of messages between objects, i.e., method invocations during run time, can be grouped together with respect to caller and callee. If messages are modelled as associations between objects and represented by edges in a graph, then this is an example for grouping edges instead of nodes. It is often useful to label the resulting edge with the number of messages it represents.
- **Dynamic method calls within a certain interval of time:** Messages can also be grouped with respect to given intervals of time. See for example [KM96].
- **Classes to files and files to directories:** In many programming environments, grouping classes to files and files to directories conforms to grouping classes to packages.
- **Objects to processes or physical machines:** Objects can be grouped to their containing processes (or threads) or the physical machines they reside on. This grouping is useful, e.g., to examine the run-time performance of an concurrent (for processes) or distributed (for machines) object-oriented system.

and many more.

The examples given above are applied in different contexts, mostly depending on the kinds of entities involved. Groupings containing *static* entities (classes, methods) are often used in model capture as well as forward and reverse engineering. *Dynamic* entities (objects, threads) are used when debugging or during performance or other dynamic analysis purposes. Groupings from or to *physical* entities (files, resources, databases, machines, processors) are needed for many development issues such as configuration management, checking consistency of compile dependencies, compile time analysis or documentation.

## 2.3.3 Definitions

We now want to provide an formal definition for our concept of grouping. First, we give a list of requirements which a concept should fulfill in order to be applicable in our tasks.

In the following, we will only consider grouping entities (nodes). Grouping associations (the elements of a relation, i.e., edges) can in principle be defined in an analogous way.

<sup>&</sup>lt;sup>11</sup>http://www.objectivesoft.com/ look.html

• The result of every legal grouping on every legal structure should be well defined.

A tool implementing grouping should be able to give a meaningful result for any legal input.

• Entities should not just disappear. Every entity of a detailed view should be represented in an abstract view, though several entities of a detailed view may be represented by the same entity in the abstract view.

A change in a method means in turn changing its containing class and package. For example, if an abstract view only represents packages, it must be possible to determine from the model which entities have been changed. Conversely, if we determine from the model that a package has to be examined then we must be able to find out which classes and methods are meant in detail.

• Dependencies should be maintained. If two entities in a detailed view obey a certain dependency, their representants should have a dependency of the same type.

For example, if we change a methods interface in one package then we want to see from an abstract view, showing only packages, whether this change may affect other packages.

• Our concept of grouping should be well defined also in combination with other concepts such as filtering information.

Put together, these requirements already set up a kind of an informal definition or concept of grouping.

We represent the structure of a system as a *typed graph*. The entities of a structure and their associations are represented by nodes and edges respectively. A set of associations between a set of entities defines a relation on these entities. Nodes and edges may have different *types*. *Node types* are for example "Class", "Method", "Package", "Object"; *edge types* for example "hasMethod", "contains", "creates".

A group in terms of a graph which represents a system structure is just a set of nodes. Grouping is then replacing one or more groups (i.e., the corresponding nodes) by a node representing these groups while maintaining the edges between the nodes (see Figure 2.48). Formally, grouping and related terms are defined as follows. We use graph theoretic and relational algebraic terms in their usual definitions. Please note: for simplicity, we use the terms "Graph" or "directed Graph" here for what is exactly spoken a *directed 1-Graph* (see for example [Har69] or [SS89]). This follows the terminology adopted throughout most of the related work.

- A group is a subset of all entities under consideration (including groups themselves).
- A grouping is a surjective graph homomorphism  $\Phi$  mapping a Graph G = (V, E) to a graph G' = (V', E'). The elements of a single group are being mapped to a single node, which represents this group.
- A certain grouping  $\Phi$  can be given
  - 1. by a function H mapping the nodes of G to the nodes of G'. Since  $\Phi$  is a graph homomorphism, it is completely determined by H. Due to the property of  $\Phi$  as a surjective homomorphism, H is a *total, unambiguous and surjective relation*, i.e., a surjective function from V to V'.
  - 2. by its *spanning relation* S mapping every entity to be grouped to its surrounding group. S may be any *unambiguous* relation, i.e., partial function. We can compute H from S by

$$H = (I \sqcap \overline{SL}) \sqcup S =: I \oplus S$$

This is probably the definition most frequently used as input for tools, since it does not need to store nodes mapped to themselves. The notion of the overwrite operator for relations " $\oplus$ " is borrowed from the specification language Z as described in [Spi92].

3. by an *equivalence relation* A defining a group in terms of entities that are all equivalent.

$$E = \{(x, y) | H(x) = H(y)\} = HH^T$$

In this case, the grouping H is defined as the *canonical mapping* of A producing the *equivalence classes* 

$$V/A = VH$$

- The *trivial grouping* is the identity, which maps every graph to itself.
- The *cardinality* of a group is the number of its elements (since a group is a set).

For the ease of writing, we often identify these representations H, S and A with the grouping  $\Phi$  itself in the remainder of this report.

#### 2.3.3.1 Filters and views

Filtering is a concept complementary to grouping. Both grouping and filtering are needed to define a concept of views. While a surjective homomorphism defines a grouping, an injective homomorphism defines a filter. A view on a system can be defined in a formal way by a combination of groupings and filters.

• A *filter* is an injective graph homomorphism. A filtered view is a *partial sub-graph* of an other view.

A filter may select only a subset of the nodes and edges or restrict the graph to those nodes and edges fulfilling certain properties<sup>12</sup>.

- A *view* on a model is the result of any combination (a sequence of concatenations) of groupings and filters applied to a model of a system. If a model is represented as a graph, then every view of this model is also a graph.
- The model itself is named the *complete view*.
- A view  $v_2$  is said to be *more abstract* than another one  $v_1$ , if a non trivial grouping H exists, which maps  $v_1$  to  $v_2$ , i.e.,  $v_2 = v_1 H$ .

This means that the process of abstraction is defined in a way that nothing just disappears, but details are (temporarily) hidden.

New views are produced from existing ones by sequences of abstractions and selections. In this way, abstraction and selection form operators which transform views. Since abstraction and selection can be used in different order, we will spend a short look on the properties regarding this matter.

Figure 2.49 shows an example where an abstraction and a selection are applied to a graph in different order. Unfortunately, the resulting graph differs depending on if abstraction or selection was first. In other words, abstraction and selection do not *commute* in every case.

The fact that the order of abstraction and selection matters is important for the interpretation of results of an analysis. Since abstraction only collapses information, while selection really omits it, abstraction should be performed first, if there is a choice.

 $<sup>1^{2}</sup>$  Selection and restriction are special kinds of filters on heterogeneous relations. The distinction is taken from relational algebra as used in database theory. To further confuse the situation in SQL a restriction is written within the "select" statement.

#### 2.3.3.2 Group types

Every entity has a type (e.g., class, object, method, file). These are types on the meta-level and should not be confused with the types declared in the program or in the specification of the programming language.

The type of a group is determined by the possible types of entities it can hold (which may be more than one). Additionally, a group may be represented by an entity of a further type. For example the group of objects of a given class may be represented by this class (though it is not equal to this class).

#### 2.3.3.3 Group operations

A model of a system may provide operations on its entities or relations. For example a user may want to click onto a node to see or edit its properties, e.g., the source code it represents. For groups, there are additionally special operations available:

**collapse:** replace the set of entities contained in the same group by a representation of this group (e.g., in a certain view)

expand: replace the group by its elements (e.g., in a certain view)

To perform an operation on each single element of a group or to filter a subset of elements from a group fulfilling certain properties (i.e., a predicate) there are well known operations from functional programming [BW88]:

**map:** takes an operation o and a set (or group) s and performs o on each element of  $s^{13}$ 

filter: takes a predicate p (a function returning true or false) and a set (or group) s and returns a set s' containing those elements x of s, for which p(x) is true

#### 2.3.3.4 Visualising groups

A major goal for the concept of grouping is to support high level visualisations. There are different possibilities, how to realise such visualisations.

Since grouping produces views which are simply graphs, the visualisation of the result of a grouping is straightforward. Groups that are collapsed are shown as a node replacing the elements contained in the group. Relations of those elements are propagated to the group as required by grouping as a graph homomorphism.

But we also want to show the original situation. There we want to see all detailed information, but we also want to see which are the groups and which entities belonging to the different groups. Groups can also be visualised when they are not collapsed to a single node, but when they are still (or again) expanded and all their elements are present. This is especially important when groups are defined interactively by the user. Possibilities for such visualisations are

- Drawing a shape (e.g., a box) around the elements to be grouped. This requires all elements of a group to be localised near each other. This is for example the solution defined in the UML for packages. We drew groups in this way in Figure 2.48.
- Drawing all the entities contained in a group in the same way, e.g., in the same colour or with the same shape.
- The group is shown as an additional node which is connected to its elements by its defining containment relation. This way, elements of the detailed view and the abstract view are shown at the same time. We used this visualisation in Figure 2.49.

<sup>&</sup>lt;sup>13</sup>Sometimes "map" is referred to as "foreach", but this term is often used with a different meaning in the area of concurrent programming.

## 2.3.4 Tool support

Our tool set GOOSE already supports several common use-cases requiring grouping functionality. Most of it is implemented in the program REVIEW for querying and manipulating graph data in GML format.

REVIEW reads a graph from input, a set of commands is being performed on the graph, and result is given as output. Several commands are available to deal with grouping:

**collapseList** Collapse a list of nodes to a new node. The command reads a file that contains the name of every node that is to be collapsed. Each line has to contain one node name. The first line is the name of the new node to create. The collapsed nodes are deleted from the resulting graph.

The new node replaces the deleted nodes in all edges, i.e., the new node is connected with another node of the graph, iff one of the nodes in the given list was and that node is not in the given list.

- **collapseMethods** Every node in the graph of type *Method* are collapsed with the class node that they belong to. The method nodes are deleted and the class node replaces the deleted nodes in all edges. So this command abstracts from methods but keeps the relations between classes through methods.
- **collapseSubsystem** All nodes that belong to a subsystem are collapsed with the subsystem node. (Similar to collapse nodes.)
- **collapseSubsystemList** Reads a list of subsystem node names from file. Each line has to contain one name. REVIEW invokes *collapseSubsystem* on each name in the file.
- **addHierarchy** Adds a hierarchy to the program graph. The hierarchy is created by subsystem nodes and *consistsOf* edges. It forms a tree where the leafs are class nodes.
- **makeSimple** If there is more than one edge between any node, then replace them by only one edge, even if the two edges are of different types. This command is useful for graph layout algorithms that require simple graphs, i.e., graphs with at most one edge between nodes.
- **makeSimpleTypes** If there is more than one edge of the same type between any node, then replace them by only one edge of the same type. Use this command, if you don't care about the individual relationship but about the fact that there is a relationship of a certain type.

All but the last two commands group nodes. The last two commands implement two frequently needed edge groupings.

To give an example, Figure 2.46 was derived from Figure 2.45 by REVIEW by using the commands:<sup>14</sup>

collapseMethods
deleteSingles
makeSimpleTypes

Another useful tool is a script analysing a directory structure and producing a subsystem hierarchy as required by the addHierarchy command. This is applicable in programming environments where the subsystem structure is denoted by the structure of files and directories [Rit98]. This is for example the concept used by Java in most cases.

A tool providing visualisation for groupings which conform to the UML notation has also been implemented within the context of this project, see [Kos98].

<sup>&</sup>lt;sup>14</sup>In fact during normal usage, all these commands are invoked automatically by a script.

## 2.3.5 Open issues

Some issues are still unresolved and need further research.

**Grouping edges** It is also possible to group edges. Grouping edges occurs far less frequently but may be useful in several situations. All edges to be grouped are mapped to the edge representing the group. In most cases edges between the same nodes (called parallel edges) are being grouped, e.g., function calls between the same classes. There is no clear definition of how to group non adjacent edges, yet. Adjacent nodes may be grouped too (may be ambiguous in non-directed graphs) or such edges may be mapped to hyper-edges. We will delay a clear definition until need arises from applications or tool development.

Our tool REVIEW currently supports grouping together all edges between the same nodes having the same type.

**Non-hierarchic grouping** A non-hierarchic grouping is a grouping defined by a non-hierarchical relation. A non-hierarchic grouping contains groups which share at least one element. E.g., two packages may share the same class. Non-hierarchic grouping introduces a "shared containment" relation. Groups are connected by this relation, if they both contain the same element. A non-hierarchic grouping is not any more defined by a homomorphism, since a non-hierarchical relation is not a function.

It is not fully clear yet, which are exactly the cases where non-hierarchical grouping is needed. E.g., a shared class could also be moved to a third package with relations to this package. One example is to split up a system with respect to functional responsibilities (communication, user interface), where entities may serve for several functionalities.

In practical applications, there are both examples explicitly forbidding or allowing non-hierarchic grouping:

- A package structure must be organised hierarchically in Java.
- In contrast, the object modelling method described by Coad and Yourdon [CY91] allows different packages (in this context named "Subjects") to share the same classes.

Another open question is whether shared elements must propagate all their relations to all groups containing them or if these relations may be split up between groups.

#### **2.3.6** Summary and future work

We presented the concept of grouping which can be applied for both static and dynamic structures as they are often used in the context of reengineering software systems. It provides a formally sound way to deal with views on different levels of abstraction. We provided formal definitions for grouping, filtering and views. Filtering information is a concept complementary to grouping. Consistent views on a system result from combinations of grouping and filtering. We gave examples for different ways to visualise abstract views, and gave a brief overview of our tools supporting different tasks related to grouping.

Some of the tasks within this topic still require further work. Among them are:

- Grouping is already being supported by our tool prototypes, but there is still a wish-list of improvements. For example, more flexible and reusable methods for the specification of groupings and support for complex entities within the visualisation.
- A further look has to be taken on some formal issues which are still open, as mentioned in Section 2.3.5. Some clarification to these questions will also have to come from further practical usage in case-studies.

• A lot of knowledge exists about what "good" design means on the levels of abstraction as they are usually used in forward-engineering (e.g., the class level). For large systems it is important to gain similar kind of knowledge on higher levels of abstraction, such as the subsystem or package levels. It has to be investigated, how this can be supported by our comprehension of grouping.

## 2.4 Reorganisation

#### Author: Benedikt Schulz

In this section we present a new technique for the reorganisation of object- oriented systems. The key idea of our approach is introduce design patterns into a system through the application of refactorings. The use of design patterns leads to a more flexible design, whereas the application of refactorings ensures the correctness of the transformations.

The section is organised as follows. In section 2.4.1 we describe the fundamentals of our approach involving design patterns and refactorings. These two techniques are combined into a powerful technique for reorganisation in section 2.4.2. Our approach is discussed and evaluated in section 2.4.3. We present related work in section 2.4.4 and conclude in section 2.4.5 with suggestions for further work.

## 2.4.1 Fundamentals

Changing a system by hand is an extremely time-consuming, difficult and error-prone task. Consider the case of simply renaming a method of a class. After the name has been changed, all places where this method is called, which can be spread over the whole system, have to be identified and changed accordingly. Because of these difficulties, a *technique for reorganisation should allow for tool-support*.

The aim of reorganisation is not to add new functionality to a system, but to make it more flexible in order to make it easier to add the desired new functionality. Therefore the reorganization should not change the behaviour of the system. Thus the *technique must define the notion of system behaviour making it possible to prove, that a performed reorganisation-step does not change this behaviour.* 

In this section, we present two approaches for the reorganisation of object-oriented systems. In 2.4.1.1 we describe the high-level approach first presented in [Zim97]. This approach describes design patterns as operators rather than as building blocks. The low-level approach of William Opdyke, who describes in [JO93] a set of transformations he calls *refactorings*, is presented in 2.4.1.2.

#### 2.4.1.1 Design Patterns as Operators

Design patterns have recently gained much attention in the field of software engineering. Many people consider them to be a promising approach for overcoming some fundamental problems in the design and reuse of object-oriented software. Design patterns were first introduced to the software community by Gamma et. al [GHJV95], [Gam91] in the early 90s. They are descriptions of solutions to a set of common, recurring design problems within a particular context. They consist of a pattern name, a problem description, the presentation of a well-proven solution and the consequences and trade-offs of applying the pattern. The essential advantages of the application of design patterns are:

- Better software design can be achieved via the reuse of well-proven solutions for general design problems.
- Design patterns form a standardised terminology for the modelling of software systems. They ease the documentation of design decisions and make them more understandable.
- Design patterns are language independent.

Over the years, some weaknesses in design patterns have been identified. One major problem is the absence of a more formal classification of design patterns and their mutual relationships [Zim95]. Due to the commonly very informal description of the purpose and the structure of a design pattern it is often hard to choose the right pattern variant.

Another major problem is the lack of a systematic way to integrate design patterns into existing systems. The descriptions of design patterns are so far mainly focused on the depiction of the target structure resulting from the application of the design pattern. The process of pattern application itself is often not taken into consideration. That means that up to now, design patterns have been seen as building blocks [BMR<sup>+</sup>96], rather than as operators. This allows for the easy construction of new systems, but makes it very hard to introduce design patterns into existing systems in order to reorganise them.

In order to overcome these shortcomings, a new approach is presented in [Zim97]. The author introduces a novel concept to support the systematic application of design patterns to existing software systems. In his approach he considers design patterns to be operators, whose application transforms an existing design into a target design, fulfilling the necessary requirements. In order to make a systematic approach possible, the process of design pattern application is split into five steps (see Figure 2.50). The example given in Figure 2.51 and 2.52 illustrates the algorithm of applying the pattern *Bridge* to an existing design:

- 1. **Identification of the problem structure**: The software engineer identifies the part of the existing design that should be reorganised. In the given example (see figure 2.51), the application of operator *Bridge* aims at de-coupling the abstraction (superclass Set) from its implementations. The problem structure comprises the class Set and its descendants.
- 2. Check of pre-conditions: All pre-conditions must hold in order to apply the design pattern operator. In our example, the implementations HashSet and ListSet must provide the same interface as the abstraction Set in order to be interchangeable.
- 3. **Parameterised transformation of the problem structure into the target structure**: In this step the design pattern application itself is put into place. The process is divided into generic aspects (depending on the design pattern) and application specific aspects (parameters). The user parameterises the process by choosing an appropriate variant of the design pattern. Additional parameters are for example the set of methods which will be delegated or the decision whether the class of the implementation object shall be interchangeable at run-time or not.

In our example, a new abstract class SetImpl (the new implementation superclass) is introduced. A new attribute impl is added to the class Set (**3a**). The methods of Set are divided into two groups; implementation-independent methods which remain in Set and implementation-dependent methods which are moved to SetImpl. This division is performed by the user. Set delegates all calls to implementation-dependent methods to SetImpl (**3b**). Now the user may decide whether he wants the implementation objects to be interchangeable at run-time or not (**3c**). In our example, the method changeImpl provides the corresponding functionality.

- 4. **Reorganisation of context**: The transformation of the problem structure often causes a change to its interface. The clients of the interface affected by the change have to be reorganised in order to provide the same functionality as before. In our example all calls of the form new(HashSet) are replaced by new(Set(HashSet)).
- 5. **Check of post-conditions**: Post-conditions must hold after the design pattern has been applied. It is often hard to formalise post-conditions as they usually depend on semantics. Therefore they are often described informally. In our example the condition "*the class of the implementation object is interchangeable at runtime*" must hold after the application of the operator *Bridge*.

The needed transformations are described in terms of meta model operations. For this purpose, a meta model comprising the entities *Class, Method, Param*, and *Attribute* and a description language are introduced. This language consists of three sets of constructs:

- transformation operations, e.g., applyDelegation(<from>, <to>, <Params>)
- parameterisations by the user, e.g., <InputVar>?:<VarType> [<InitList>]
- miscellaneous operations, e.g., conditions and labels

With the approach described above, the application of design patterns becomes an algorithmic process. In the "building block" approach, the user must know about and understand the informal description of the design pattern and all its relationships and variants in order to apply the pattern to the design. The explicit separation of generic and specific aspects (e.g., parameters) of the design pattern and the operational presentation simplify its application, since the user only has to cope with use-case specific aspects of the design pattern application. Beyond this, the transformational approach enables to use design patterns as *reorganisation operators* in order to reorganise existing systems by identifying the problem structure and transforming it into an improved target structure using the algorithm described above.

#### 2.4.1.2 Refactorings

In his PhD thesis [JO93], William Opdyke presented an approach for the support of the evolution of objectoriented systems. He defined a set of so-called *refactorings*, which are implementable in a tool and showed that the application of these refactorings does not change the system behaviour.

Opdyke calls the property of not changing the behaviour of a program *semantic equivalence* and defines it as follows:

Let the external interface to the program be via the function *main*. If the function *main* is called twice (once before and once after a refactoring) with the same set of inputs, the resulting set of output values must be the same.

Besides the semantic equivalence, Opdyke additionally defines several more invariants to ensure the syntactic and semantic correctness of the refactorings. He does not, for example, allow the use of *multiple inheritance* or the *redefinition of inherited member variables*. Every refactoring has to preserve these invariants.

The refactorings are divided into two subsets: One subset contains 26 *low-level refactorings* focusing on simplicity in order to make it easy to prove that system behaviour is not changed. The other subset consists of three *high-level refactorings*, which are defined by using the low-level refactorings. The proof of the behaviour preserving nature of the high-level refactorings is also furnished in terms of low-level refactorings.

Low-level refactorings allow the creation (e.g., create\_empty\_class), deletion (e.g., delete\_member\_function) or modification of program entities (e.g. add\_function\_argument). Furthermore, refactorings move member variables to (e.g., move\_member\_variable\_to\_superclass) and simple composite refactorings (e.g., convert\_code\_segment\_to\_function) belong to the low-level refactorings. The description of these refactorings is as follows:

- Name: Denotes the name of the refactoring.
- Description: Describes the refactoring and its variants.
- Arguments: Enumerates the set of arguments and their meaning.
- **Pre-conditions**: Lists the set of pre-conditions, which have to hold in order to assure that the semantics of the program remain unchanged when applying the refactoring.
- **Proof of correctness**: A semi-formal proof that ensures that the refactoring does not change either the invariants or the semantics of the program, assuming that all pre-conditions hold.

The three high-level refactorings described in [JO93] are:

• Creation of an abstract superclass: The common abstraction (methods, member variables) of two or more classes is factored out into an abstract superclass.

- **Subclassing and simplifying conditionals**: Conditional statements, dependent on the internal state of the object, which are used to select the different behaviour of an object are replaced by an implementation using various subclasses and polymorphism.
- **Converting an inheritance relation into an aggregation**: A relation modelled by inheritance is converted into aggregation. The functionality, which was inherited can be accessed by delegating it to the component.

The description of the high-level refactorings contains one additional section describing how it is decomposed into low-level refactorings, how the pre-conditions can be derived from them and how the proof of correctness is built upon the proof for the low-level refactorings.

Since Opdyke presented the pre-conditions as well as the necessary restructuring in a detailed way, it was possible to construct tools for the support of the refactoring task. Due to the complexity of the C++ language <sup>15</sup> Ralph Johnson and his group focused on Smalltalk and constructed a *Refactoring Browser for Smalltalk* [RBJ97a]. This tool supports a selection of the refactorings described in [JO93] and was successfully used for the development of *Hotdraw*.

Berthold Mohr showed in [Moh98] that it is possible to implement a tool supporting refactorings for a subset of C++. He implemented the refactoring for the conversion of an inheritance relationship into an aggregation including all necessary low-level refactorings.

In spite of the promising results in formalising [JO93] and implementing [RBJ97a][Moh98] refactorings, there are some restrictions to the approach:

- **Behaviour preservation**: The described approach of assembling high-level refactorings from low-level refactorings requires behaviour preservation for every application of a low-level refactoring. It may be desirable to ignore this constraint and to guarantee behaviour preservation for the entire high-level refactoring only.
- Language dependance: It is not possible to specify all the pre-conditions in a language-independent way, because the semantics of a program heavily depend on the semantics of the underlying programming language.
- **Pre-condition checking**: Checking whether a pre-condition holds or not, is a difficult task which requires global dataflow analysis techniques. This is, in general, an unsolvable problem.
- Level of abstraction: Even the high-level refactorings are not on a level of abstraction which is necessary for reengineering.

Especially the inappropriate level of abstraction is annoying when refactorings must be used to re-engineer large object-oriented systems. We present a new approach to overcome this shortcoming in the following section.

## 2.4.2 The FAMOOS approach

This section will begin with a definition of a set of requirements a reorganisation methodology must meet. Then we show that the approaches presented in section 2.4.1 are not adequate to fulfil these demands. In the last part, we introduce a new methodology for the reorganisation of large object-oriented systems.

We have identified the following requirements which must be fulfilled by a reorganisation methodology in order to be practically applicable for large object-oriented software systems:

<sup>&</sup>lt;sup>15</sup>Opdyke writes in [JO93]: "The C++ language is a semantically complicated language, supporting machine level operations such as pointer arithmetic; these complexities make it difficult to more precisely define what behaviour preservation means for C++ programs."

- Language independence: The methodology should not rely on a specific object-oriented programming language.
- Level of abstraction: The methodology should aim at the design level rather than the implementation level, since most of the problems of an object-oriented system which can be solved by reorganisation (e.g., lack of flexibility) concern the design of that system.
- **Behaviour preservation**: The goal of reorganisation is not to change the functionality of a given system, but to improve the structure of that system in order to make these sorts of functionality changes easier. Therefore a reorganisation operation should not change the systems behaviour. The methodology must be proven to preserve behaviour.
- **Tool support**: The methodology requires for tool support, since reorganising a large system by hand is error-prone and time-consuming.

In the preceding section, we presented two different approaches to the reorganisation of object-oriented systems. Although both of them are very promising, they suffer from some weaknesses which make it difficult to put them into practice. While design pattern operators fulfil the first two requirements, (i.e., they are language independent and deal with design aspects) they fail to fulfil the last. Thus, not only is the underlying model too abstract to be implemented in a tool, but the pattern operators also lack proof of behaviour preservation. Unlike low-level refactorings, the atomic operations which make up a design pattern operator do not necessarily lead to behaviour preserving states. In the example given in Figures 2.51 and 2.52, the application of step 3 changes the behaviour of that system as constructs of the form Set s = new(HashSet); are no longer valid due to the broken inheritance relationship. On the other hand, refactorings provide at least a semi-formal proof of behaviour preservation. Even more, it was shown that they can be implemented in tools [RBJ97a], [Moh98]. However, refactorings, even the high-level refactorings are not at a level of abstraction which is needed for reorganisation. Beyond this the specification of refactoring-specific pre-conditions heavily depends on the semantics of the underlying programming language. Thus it is not possible to use the refactoring-approach in a language independent way.

Our approach offers the best of both worlds. We combine the two approaches in order to overcome the shortcomings of both of them by *implementing the design pattern operators* [Zim97] *with Opdyke's refactorings* [JO93].

Whereas in [Zim97] the description of the operators is very informal, we are now able to formalise the operators, as well as their pre-conditions and post-conditions with the aid of the refactorings. This means, that we replace the meta-model and the operator-language defined in [Zim97] with the model and language defined in [JO93].

The different phases for applying a design-pattern operator in [Zim97] are refined as follows:

- 1. **Identification of the problem structure**: As in [Zim97] the software engineer has to identify the part of the system to be reorganised first. Depending on the selected problem structure and the design pattern operator to be applied, the sequence of refactorings can be determined.
- 2. Checking the pre-conditions: The preconditions are derived from the sequence of refactorings. They are checked before each application of a single refactoring. However it is possible to derive a set of preconditions which, when satisfied, ensures the success of the application of the whole sequence.
- 3. **Parameterised transformation of the problem structure into the target structure**: This step performs the transformation by executing the sequence of refactorings. As in [Zim97] this may require a parameterisation by the user.
- 4. **Reorganisation of context**: Since the reorganisation of the context is part of the application of a single refactoring, the context of the problem structure does not need to be reorganised after the execution of the sequence of refactorings.

5. **Checking the post-conditions**: The post-conditions hold after the application of the design pattern, because every refactoring preserves the given invariants, while at the same time not changing the behaviour of the system. An explicit checking of the post-conditions can thus be omitted.

Figures 2.53 and 2.54 depict the third step of the application of the design pattern operator *Bridge* as performed by our new approach. The sequence of necessary transformation steps is different from the sequence in [Zim97], since our approach requires behaviour preservation at every step.

In step 3a a new, unreferenced class AbstractSet is introduced. AbstractSet is assigned to the same superclass (not shown in the figure) as Set in step 3b. The introduction of a component relationship between Set and AbtractSet, the copying of the interface of Set to AbstractSet and the delegation of all methods of AbstractSet to its component Set is depicted in step 3c.

Step 3d as depicted in Figure 2.54 contains the most crucial refactoring. First the constructor method of AbstractSet is changed to include one parameter determing the actual implementation (e.g., HashSet or ListSet) of the set. Then every construction of a set (subclass of Set) is changed to use the new constructor of AbstractSet and every variable definition using Set or one of its subclasses as a type is changed to AbstractSet. Finally in step 3e a new method getImpl can be introduced in the class Set to determine the actual implementation class of the set.

Our approach fulfils all the already mentioned requirements. The method is **language-independent** on the design-level, because design patterns are language-independent. The design level is the **appropriate level** of **abstraction** for a reengineering task. The application of design pattern operators is correct in the sense of **behaviour preservation**. Since every single primitive transformation is fully specified together with pre-conditions and post-conditions, it is possible to implement **tools to support our approach**.

## 2.4.3 Discussion

In this section we present some results of the application of our approach to an existing system. For this purpose we implemented a tool prototype which provides several simple or complex refactorings for C++ (e.g., add\_class, change\_superclass, convert\_inheritance\_into\_aggregation), [Moh98] using the FAST C++- parser [Sem97]. The sequence and order of refactorings required to implement a particular design pattern is determined via shell scripts. We tested our tool with an existing software application for the visualisation of flow data from the field of hydraulic engineering. The aim of our experiment was to make some hot-spots in the system more flexible in order to ease the functional extension of the software. Some of our experiences are described below.

The software system mentioned was developed for a Windows NT platform using a proprietary graphic library providing several drawing classes (e.g., Shape, Circle) and implemented using the Windows Graphic Device Interface (GDI) as a basis. Due to performance aspects, with a new version of that software, it was decided to use the Windows DirectDraw interface. In order to be compatible with older versions of Windows NT, one requirement was that the platform-dependent graphic sub-system should be interchangeable.

In the original design the platform specific graphic output routines (e.g., drawPoint) were coded in Graphic using GDI-routines. All graphical shapes inherited the needed base functionality from Graphic (see Figure 2.55). Therefore it was hard to introduce new functionality to support more than one graphic sub-system, since the platform-dependent parts were strongly coupled to the domain-specific classes. On the other hand, the relationship between Shape and Graphics was equivalent to an *implementation inheritance* and Graphic was never used as a static type of a variable or in a polymorphic way throughout the entire system.

One way to solve that problem is to *objectify* the parts of the problem structure which should be interchangeable. We decided to apply the pattern operator *Bridge* in order to decouple the platform-specific functionality from the domain-specific parts. Further requirements were that the reorganisation must preserve the needed inheritance relationships of both Graphic and Shape, since some methods of those classes were implemented using inherited functionalities<sup>16</sup> and the process should provide the possibility for appropriate parameterisation by the user (e.g., the selection of delegated methods or inheritance relations).

In order to achieve the necessary target structure, we used a variant of pattern operator *Bridge* consisting of the refactoring convert\_inheritance\_into\_aggregation. This high-level refactoring converts an inheritance relationship into a component relationship [JO93] and consists of the following steps:

- 1. Create a new member variable impl of type Graphic.
- 2. For each member function inherited from Graphic create a member function in Shape. Calls to these functions are delegated to the corresponding methods in Graphic.

Due to the requirement, which states that the inheritance relationship should be preserved, the refactoring was slightly modified as follows. Instead of breaking up the inheritance relationship the refactoring change\_superclass is called for class Shape. Shape becomes a subclass of the superclass of Graphic. Now one can easily change the code for platform-dependent graphic support by creating the appropriate subclasses GraphicGDI and GraphicDDraw. The corresponding methods of Graphic are changed to become *abstract*. All these steps can be supported by tool-based refactorings as well.

The process of applying design pattern operators to existing systems resulted in some difficulties. First, pattern operators must obey existing inheritance relationships if necessary. This sometimes leads to slightly complicated refactoring sequences and often requires user interaction. Another problem was that sometimes our tool could not check all pre-conditions automatically. However, this does not occur in only this specific case but is generally a problem [JO93], since pre-condition checking often requires global data flow analysis. Beyond this, our tool only supported a subset of C++ syntax (not supported is pointer arithmetic and multiple inheritance to name a few). Nevertheless, our tool significantly eased the reorganisation of the software. In addition, our tool works directly on the source code and no pretty printing is required. Thus programmers recognise their own code after the reorganisation operation has been completed.

## 2.4.4 Related Work

Since reengineering and reorganisation has been recognised as an important field of research, many contributions to this topic can be found in the literature. We present in the following only work which is closely related to our work.

*Schema evolution of object-oriented databases* can be seen as the basis for Opdykes work. In [BCG<sup>+</sup>87] a set of primitive schema evolution operators for the object-oriented database-system *ORION* is presented and categorised. However the focus of this work is more on how to deal with different versions of persistent objects than on code transformations in general.

One aspect of improving the structure of object-oriented systems is discussed in [Cas92] and [Cas93]. The author describes both a global and an incremental approach for reorganisation of inheritance relationships. The approach focuses on the reduction of redundancies in the class definitions. The same author presented in [Cas95b] an excellent overview of existing techniques for the management of class evolution in object-oriented systems.

Several approaches cover parts of the reengineering life-cycle. In [Mei96][vW96][FMvW97] the authors present an approach for visualising *fragments* of object-oriented systems. These fragments can be classes, subsystems or even design patterns. Transformations of the fragments are implemented by using the refactoring browser for Smalltalk [RBJ97a]. However the authors do not describe any high-level transformations of fragments.

The *MeTHOOD*-project [GD97] covers almost the whole reengineering life-cycle. The description of the transformations is given in a way similar to the description of the design pattern operators. However no

<sup>&</sup>lt;sup>16</sup>Note: These inheritance relations were left out for simplicity in Figure 2.55 and 2.56

proofs of correctness are given and it seems that it is not possible to perform the transformations automatically.

Adaptive Programming [Lie95][HS96b] separates the structure of object-oriented systems from the algorithms operating on these structures. This approach aims at a change avoidance rather than at a systematic way of increasing the flexibility of object-oriented systems.

## 2.4.5 Summary

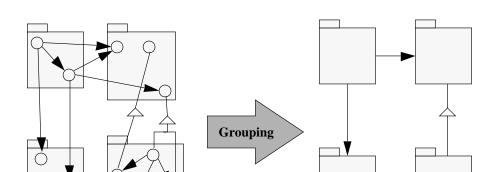
In this chapter we presented a new approach to the tool-based reorganisation of object-oriented software. Implementing design pattern operators with refactorings, we combined two existing approaches in order to overcome the weaknesses of both of them. Our methodology uses refactorings, which are proven to preserve behaviour and can be implemented in tools. Since every application of a refactoring leads to a behaviour preserving state, composing design pattern operators of refactorings ensures the behaviour preservingness of the entire pattern operator. Beyond this, design pattern operators are language independent, at least at design level. Pattern operators aim at the right level of abstraction, since most of the problems of an object-oriented system can only be solved by improving the systems design. We demonstrated the feasibility of our approach by applying design pattern operators composed of simple or complex refactorings to an existing system in order to reorganise the hot-spots of that system.

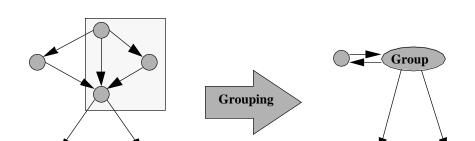
# 2.5 Reverse and Reengineering Patterns

We have shown how metrics, graphs and metrics combination and grouping software entities can help in understanding identifying important entities, overall applications organizations and detecting problems. Based on these results, refactorings can help fixing the identified problems.

Having techniques for helping in the reengineering tasks is of primary importance. Without techniques and tools supporting them reengineering applications would be nearly impossible. However, relying only on techniques and tools does not suffice. Indeed knowing how to apply a given technique is not enough and knowing when to apply or not is also a key point in the success of the reengineering. Moreover, knowing what are consequence of applying a given technique and in which order apply them is also a key information.

The following part is a first attempt to record all the implicit knowledge that reengineers have been developed over the years. This knowledge is described in pattern format to ease the readability and highlight the motivation and the forces behind every patterns. Please notice that this work is not into a final and definitive stage as the gathering of such an information is slow and difficult.





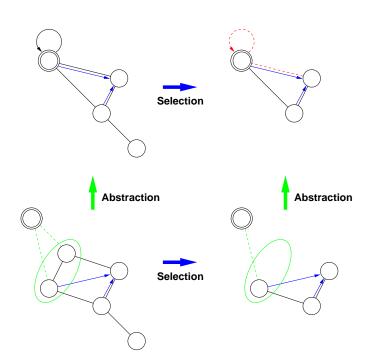


Figure 2.49: Abstraction and selection do sometimes not commute

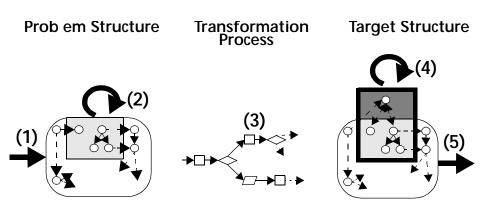


Figure 2.50: Systematic process of design pattern application

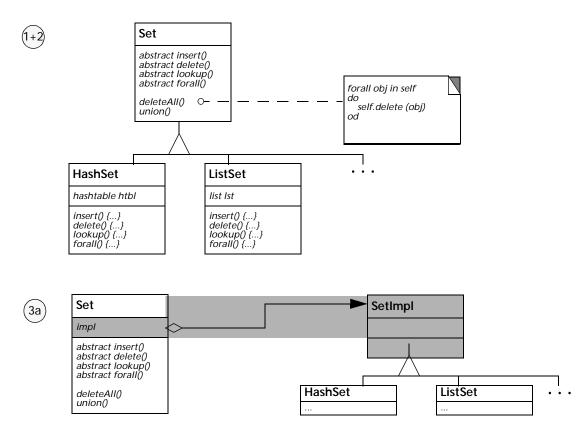


Figure 2.51: Application of Design Pattern Operator Bridge (I)

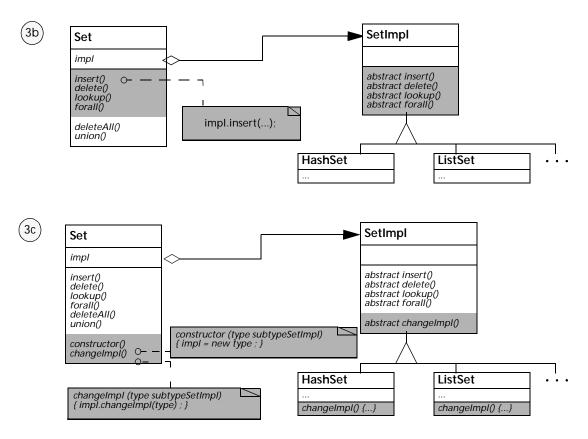


Figure 2.52: Application of Design Pattern Operator Bridge (II)

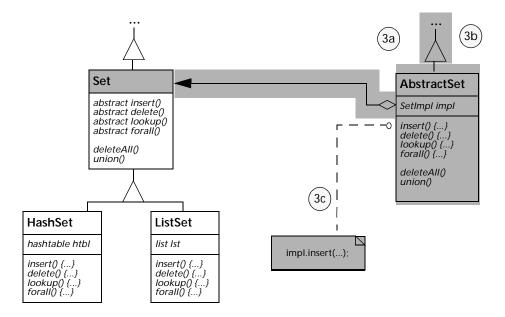


Figure 2.53: Application of a Bridge with refactorings (I)

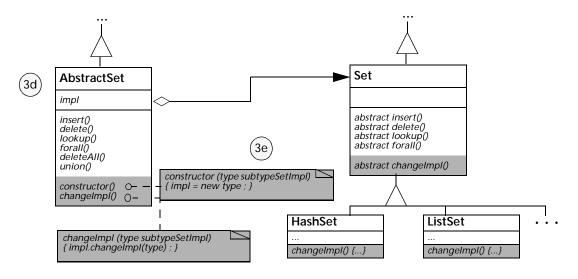


Figure 2.54: Application of a Bridge with refactorings (II)

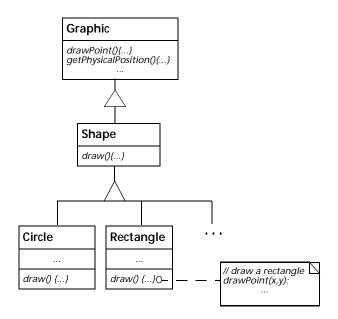


Figure 2.55: Problem structure before reorganisation

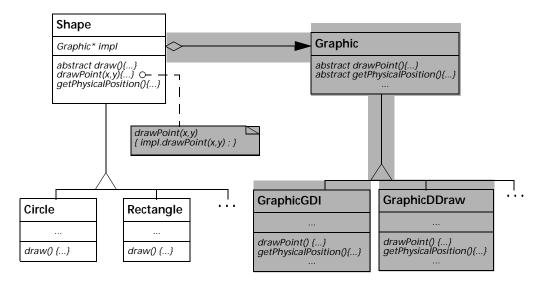


Figure 2.56: Improved target structure after the reorganisation and functional extension

# Part II

**Reverse Engineering** 

## **Chapter 3**

# **Reverse Engineering Patterns**

### **3.1** Patterns for Reverse Engineering

This pattern language describes how to reverse engineer an object-oriented software system. Reverse engineering might seem a bit strange in the context of object-oriented development, as this term is usually associated with "legacy" systems written in languages like COBOL and Fortran. Yet, reverse engineering is very relevant in the context of object-oriented development as well, because the only way to achieve a truly reusable object-oriented design is recognized to be iterative development (see [Boo94], [GR95], [JGJ97], [Ree96]). Iterative development involves refactoring existing designs and consequently, reverse engineering is an essential facet of any object-oriented development process.

The patterns have been developed and applied during the FAMOOS project [http: //www.iam.unibe.ch/  $\sim$ famoos/]; a project whose goal is to produce a set of re-engineering techniques and tools to support the development of object-oriented frameworks. Many if not all of the patterns have been applied on software systems provided by the industrial partners in the project (i.e., Nokia and Daimler-Chrysler). These systems ranged from 50.000 lines of C++ up until 2,5 million lines of Ada. Where appropriate, we refer to other known uses we were aware of while writing.

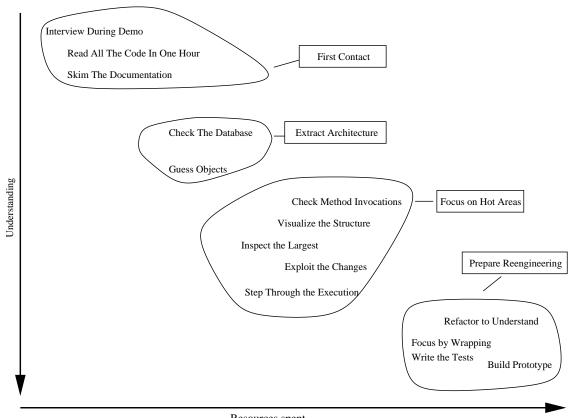
In its current state, the pattern language can still be improved and we welcome all kinds of feedback that would help us do that. We are especially interested in course grained comments —does the structure work? is the set of forces complete? is the naming OK ?— rather than detailed comments on punctuation, spelling and lay-out.

**Acknowledgments.** We would like to thank our EuroPLoP'99 shepherd Kyle Brown: his comments were so good we considered including him as a co-author. We also want to thank both Kent Beck and Charles Weir who shepherded a very rough draft of what you hold right now. Finally, we must thank all participants of the FAMOOS project for providing such fruitful working context.

## **3.2** Clusters of Patterns

The pattern language itself has been divided into *clusters* where each cluster groups a number of patterns addressing a similar reverse engineering situation. The clusters correspond roughly to the different phases one encounters when reverse engineering a large software system. Figure 3.1 provides a road map and below you will find a short description for each of the clusters.

**First Contact** (p. 113) This cluster groups patterns telling you what to do when you have your very first contact with the software system.



Resources spent

Figure 3.1: Overview of the pattern language using clusters.

- **Extract Architecture** (p. 125) Here, the patterns tell you how to get to the architecture out of a system. This knowledge will serve as a blueprint for the rest of the reverse engineering project.
- **Focus on Hot Areas** (p. 133) The patterns in this cluster describe how to get a detailed understanding of a particular component in your software system.
- **Prepare Reengineering** (p. 149) Since reverse engineering often goes together with reengineering, this cluster includes some patterns that help you prepare subsequent reengineering steps.

## 3.3 Overview of Forces

All the patterns in this pattern language tell you how to address a typical reverse engineering problem. To evaluate the situation before and after applying the pattern we introduce a number of *forces*. The forces are meant to help you assessing whether the pattern is appropriate for your particular situation.

- **Limited Resources.** Because your resources are limited you must select which parts of the system to reverse engineer first. However, if you select the wrong parts, you will have wasted some of your precious resources. In general the *less resources you need to apply, the better*.
- **Tools and Techniques.** For reverse engineering large scale systems, you need to apply techniques probably accompanied with tools. However, techniques and tools shape your thoughts and good reverse engineering, requires an unbiased opinion. Also, techniques and tools do require resources which you might not be willing to spend. In general, the *less techniques and tools required, the better*.

- **Reliable Info.** A reverse engineer is much like a detective that solves a mystery from the scarce clues that are available [WC96]. As with all good detective stories, the different clues and testimonies contradict each other, thus your challenge is to assess which information is reliable and solve the mystery by coming up with the most plausible scenario. In general, the *more reliable the information you get, the better*.
- Abstraction. The whole idea of understanding the inner complexities of a software system is to construct mental models of portions of it, thus a process of abstraction. Consequently, the reengineering taxonomy of Chikofsky and Cross [CCI90], defines reverse engineering as "the process of analyzing a subject system to [...] create representations of the system [...] at a higher level of abstraction". Of course, the target level of abstraction for your particular reverse engineering step depends very much on the subsequent demands and so you don't want to get too abstract. Still in general, the *more abstract the information obtained, the better*.
- **Sceptic Colleagues.** As a reverse engineer, you must deal with three kinds of colleagues. The first category are the faithful, the people who believe that reverse engineering is necessary and who thrust that you are able to do it. The second is the category of sceptic, who believe this reverse engineering of yours is just a waste of time and that its better to start the whole project from scratch. The third category is the category of fence sitters, who do not have a strong opinion on whether this reverse engineering will pay off, so they just wait and see what happens. To save your reverse engineering from ending up in the waste bag, you must keep convincing the faithful, gain credit with the fence sitters and be wary of the sceptic. In general, the *more credibility you gain, the better*.

## **3.4 Resolution of Forces**

Table 3.1 shows an overview of how the different patterns resolve the forces. This view is especially important because it emphasises the different trade-offs implied by the patterns. For instance, it shows that READ ALL THE CODE IN ONE HOUR and SKIM THE DOCUMENTATION take about the same amount of resources and also require about the same amount of techniques and tools (very little, hence the double plusses), yet score differently on the reliability and abstraction level of the resulting information. On the other hand, GUESS OBJECTS requires more resources, techniques and tools then the previous two (one minus), but achieves better results in terms of reliable and abstract information.

## **3.5** Format of a Reverse Engineering Pattern

The pattern presented hereafter have the following format.

- **Name.** Names the pattern after the solution it proposes. The pattern names are verb phrases to stress the action implied in applying them.
- **Intent.** Summarizes the purpose of the pattern, including a clarifying *Example* on when and how to apply the pattern.
- **Context.** Presents the context in which the pattern is supposed to be applied. You may read this section as the prerequisites that should be satisfied before applying the pattern.
- **Problem.** Describes the problem the pattern is supposed to solve. Note that the prerequisites defined in the 'Context' section are supposed to narrow the scope of the problem, so readers are encouraged to read both sections together.
- **Solution.** Proposes a solution to the problem that is applicable in the given context. This section may include a *Recipe* or a list of *Hints* and *Variations* to be taken in account when applying the solution.

	Limited Resources	Tools and Techniques	Reliable Info	Abstraction	Sceptic Colleagues		
First Contact							
READ ALL THE CODE IN ONE HOUR	++	++	+	-	+		
SKIM THE DOCUMENTATION	++	++	-	+	-		
INTERVIEW DURING DEMO	++	+	0	+	-		
Extract Architecture							
GUESS OBJECTS	-	-	++	++	+		
CHECK THE DATABASE	-	-	++	+	++		
Focus on Hot Areas							
INSPECT THE LARGEST	0	_	0	0	0		
VISUALIZE THE STRUCTURE	_		+	+	+		
CHECK METHOD INVOCATIONS	-	-	+	+	0		
EXPLOIT THE CHANGES			++	+	++		
STEP THROUGH THE EXECUTION	-	0	++	-	+		
PREPARE REENGINEERING							
WRITE THE TESTS		-	++	++	0		
REFACTOR TO UNDERSTAND		-	0	+	0		
BUILD A PROTOTYPE		-	+		++		
Focus by Wrapping		0	0	0	0		

Table 3.1: How each pattern resolves the forces. Very good: ++, Good: +, Neutral: 0, Rather Bad: -, Very bad: --. Limited Resources: The less resources you need to apply, the better. Tools and Techniques: The less techniques and tools required, the better. Reliable Info: The more reliable the information you get, the better. Abstraction: The more abstract the information obtained, the better. Sceptic Colleagues: The more credibility you gain, the better.

- Forces Resolved. Describes the situation after applying the pattern. This description is done in terms of the forces.
- Rationale. Explains the technical background of the pattern, i.e. why it works.
- **Known Uses.** Presents the know uses of this pattern. Note that all patterns in this pattern language have been developed and applied in the context of the FAMOOS project. Yet, this section presents other reported uses of the pattern we were aware of while writing the pattern.
- **Related Patterns.** Links the pattern in a web of other patterns, explaining how the patterns work together to achieve the global goal of reverse engineering. The section includes a *Resulting Context* which tells you how you may use the output of this pattern as input for another one.

## Chapter 4

# **Cluster: First Contact**

All the reverse engineering patterns in this cluster are applicable in the very early stage of a reverse engineering project when you are largely unfamiliar with the software system. Before tackling such a project, you need an initial assessment of the software system. However, accomplishing a good initial assessment is difficult because you need a quick and accurate result.

The patterns in this cluster tell you how to optimally exploit information resources like source code (READ ALL THE CODE IN ONE HOUR (p. 115)), documentation (SKIM THE DOCUMENTATION (p. 118)) and system experts (INTERVIEW DURING DEMO (p. 121)). The order in which you apply them depends mainly on your project and we refer you to the "*Related Patterns*" section in each pattern for a discussion on the trade-offs involved. Afterwards you will probably want to CONFER WITH COLLEAGUES (p. 156) and then proceed with EXTRACT ARCHITECTURE (p. 125).

### **Forces Revisited**

- **Limited Resources.** Wasting time early on in the project has severe consequences later on. *Consequently, time is the most precious resource in the beginning of a reverse engineering project.* This is especially relevant because in the beginning of a project you feel a bit uncertain and then it is tempting to start an activity that will keep you busy for a while, instead of something that confronts you immediately with the problems to address.
- **Tools and Techniques.** In the beginning of a reverse engineering project, you are in a bootstrapping situation: you must decide which techniques and tools to apply but you lack a profound basis to make such a decision. *Consequently, choose very simple techniques and very basic tools, deferring complex but time consuming activities until later.*
- **Reliable Info.** Because you are unfamiliar with the system, it is difficult to assess which information is reliable. *Consequently, base your opinion on certified information but complement it using supplementary but less reliable information sources.*
- **Abstraction.** At the beginning of the project you can not afford to be overwhelmed by too many details. *Consequently, favor techniques and tools that provide you with a general overview.*
- **Sceptic Colleagues.** This force is often reinforced in the beginning of a reverse engineering project, because as a reverse engineer –or worse, a consultant– there is a good chance that you are a newcomer in a project team. *Consequently, pay attention to the way you communicate with your colleagues as the first impression will have dire consequences later.*

	Limited Resources	Tools and Techniques	Reliable Info	Abstraction	Sceptic Colleagues
READ ALL THE CODE IN ONE HOUR	++	++	+	-	+
SKIM THE DOCUMENTATION	++	++	-	+	-
INTERVIEW DURING DEMO	++	+	0	+	-

Table 4.1: How each pattern of FIRST CONTACT resolves the forces. Very good: ++, Good: +, Neutral: 0, Rather Bad: -, Very bad: --

# READ ALL THE CODE IN ONE HOUR

Author(s): Serge Demeyer, Stéphane Ducasse and Sander Tichelaar

#### Intent

Make an initial evaluation of the condition of a software system by walking through its source code in a limited amount of time.

**Example.** You are facing a 500 K lines C++ program, implementing a software system to display multi-media information in real time. Your boss asks you to look how much of the source code can be resurrected for another project. Before actually identifying what may be reused, you will leaf through the source code to get a feeling for its condition.

### Context

You are starting a reverse engineering project of a large and unfamiliar software system. You have the source code at your disposal and you have reasonable expertise with the implementation language being used.

## Problem

You need an initial assessment of the internal state a software system to plan further reverse engineering efforts.

### Solution

Take the source code to a room where you can work undisturbed (no telephones, no noisy colleagues). Grant yourself a reasonably short amount of study time (i.e., approximately one hour) to walk through the source code. Take notes sparingly to maximize the contact with the code.

After this reading time, take about the same time to produce a report about your findings, including list of (i) the important entities (i.e., classes, packages, ...); (ii) the coding idioms applied (i.e., C++ [Cop92], [Mey98], [Mey96]; Smalltalk [Bec97]); and (iii) the suspicious coding styles discovered (i.e., "code smells" [Fow99]). Keep this report short, and name the entities like they are mentioned in the source code.

**Hints.** The fact that you are limited in time should force you to think how you can extract the most useful information. Below are some hints for things to look out for.

- Functional tests or units tests convey important information about the functionality of a software system.
- Abstract classes and methods reveal design intentions.
- Classes high in the hierarchy often define domain abstractions; their subclasses introduce variations on a theme.

- Occurrences of the Singleton pattern [GHJV95] may represent information that is constant for every complete execution of a system.
- Surprisingly large constructs often specify important chunks of functionality that should be executed sequentially.
- Some development teams apply coding styles and if they did, it is good to be aware of them. Especially naming conventions are crucial to scan code quickly.

## **Forces Resolved**

- **Limited Resources.** By applying this pattern, you spend 1/2 a day (plus the time to collect the source code) to end up with a short list that is a reasonable basis for planning further reengineering efforts.
- **Tools and Techniques.** Good source code browsers can speed you up and inheritance hierarchy browsers can give you a feel for the structure of a software system. However, be wary of fancy tools as they quickly overwhelm you with too much unnecessary information and may require a lot of time to configure correctly. Printing out the source code and reading a paper version may serve just as well.
- **Reliable Info.** The concentrated contact with the code –and code is the only testimony you are sure is correct<sup>1</sup>– provides you with a rather unbiased view to start with. Moreover, by applying this pattern –especially in combination with SKIM THE DOCUMENTATION (p. 118)– you may already have encountered some contradicting pieces of information, which is definitely worthwhile to explore in further depth.
- **Abstraction.** The information you get out is fairly close to the source code, consequently the abstraction level is quite low. However the fact that you work under time pressure forces you to skip details driving you towards an abstract view of the software system.
- Sceptic Colleagues. The mere effect of asking quite precise questions after only 1/2 a day of effort raises your credit tremendously, usually enough for being allowed to continue your attempts.

## Rationale

Reading the code in a short amount of time is very efficient as a starter. Indeed, by limiting the time and yet forcing yourself to look at all the code, you mainly use your brain and coding expertise to filter out what seems important. This is a lot more efficient than extracting human readable representations or organizing a meeting with all the programmers involved.

Moreover, by reading the code directly you get an unbiased view of the software system including a sense for the details and a glimpse on the kind of problems you are facing. Because the source code describes the functionality of the system –no more, no less– it is the only reliable source of information. Be careful though with comments in the code. Comment can help you in understanding what a piece of software is supposed to do. However, just like other kinds of documentation, comments can be outdated, obsolete or simply wrong.

Finally, acquiring the vocabulary used inside the software system is essential to understand it and communicate about it with other developers. This pattern helps to acquire such a vocabulary.

<sup>&</sup>lt;sup>1</sup>Remember the old Swiss saying: "If the map and the terrain disagree, trust the terrain"

## **Known Uses**

While writing this pattern, one of our team members applied it to reverse engineer the Refactoring Browser [RBJ97b]. The person was not familiar with Smalltalk, yet was able to identify code smells such as "Large Constructs" and "Duplicated Code". Even without Smalltalk experience it was possible to get a feel for the system structure by a mere inspection of class interfaces. Also, a special hierarchy browser did help to identify some of the main classes and the comments provided some useful hints to what parts of the code were supposed to do. Applying the pattern took a bit more than an hour, which seemed enough for a relatively small system and slow progress due to the unfamiliarity with Smalltalk.

The original pattern was suggested by Kent Beck, who stated that it is one of the techniques he always applies when starting consultancy on an existing project. Since then, other people have acknowledged that it is one of their common practices.

## **Related Patterns**

If possible, READ ALL THE CODE IN ONE HOUR (p. 115) in conjunction with SKIM THE DOCUMENTA-TION (p. 118) to maximize your chances of getting a coherent view of the system. To guide your reading, you may precede this pattern with INTERVIEW DURING DEMO (p. 121), but then you should be aware that this will bias your opinion.

**Resulting Context.** This pattern results in a list of (i) the important entities (i.e., classes, packages, ...); (ii) the presence of standard coding idioms and (iii) the suspicious coding styles discovered. This is enough to start GUESS OBJECTS (p. 127) and CHECK THE DATABASE (p. 130) to improve the list of important entities. Depending on whether you want to wait for the results of SKIM THE DOCUMENTATION (p. 118), you should consider to CONFER WITH COLLEAGUES (p. 156).

# SKIM THE DOCUMENTATION

Author(s): Serge Demeyer, Stéphane Ducasse and Sander Tichelaar

### Intent

Make an initial guess at the functionality of a software system by reading its documentation in a limited amount of time.

**Example.** You must develop a geographical information system. Your company has once been involved in a similar project, and your boss asks you to check if some of the design of this previous project can be reused. Before doing any design extraction on the source code, you will skim the documentation to see how close this other system is to what you are expected to deliver.

## Context

You are starting a reverse engineering project of a large and unfamiliar software system. You have the documentation at your disposal and you are able to interpret the diagrams and formal specifications contained within.

**Example.** If the documentation relies on use cases (see [JGJ97]) for recording scenarios or formal languages for describing protocols, you should be able to understand the implications of such specifications.

## Problem

You need an initial idea of the functionality provided by the software system in order to plan further reverse engineering efforts.

### Solution

Take the documentation to a room where you can work undisturbed (no telephones, no noisy colleagues). Grant yourself a reasonably short amount of study time (i.e., approximately one hour) to scan through the documentation. Take notes sparingly to maximize the contact with the documentation.

After this reading time, take about the same time to produce a report about your findings, including a list of (i) the important requirements; (ii) the important features (iii); the important constraints; (iv) references to relevant design information. Include your opinion on how reliable and useful each of these are. Keep this report as short as possible and avoid redundancy at all cost (among others, use references to sections and/or page numbers in the documentation).

Depending on the goal of the reverse engineering project and the kind of documentation you have at your disposal, you may steer the reading process to match your main interest. For instance, if you want insight into the original system requirements then you should look inside the analysis documentation, while knowledge about which features are actually implemented should be collected from the end-user manual or tutorial notes. If you have the luxury of choice, avoid spending too much time to understand the design documentation (i.e., class diagrams, database schema's, ...): rather record the presence and reliability of such documents as this will be of great help in later stages of the reverse engineering.

Be aware for documentation that is outdated with respect to the actual system. Always compare version dates with the date of delivery of the system and make note of those parts that you suspect unreliable.

Avoid to read the documentation electronically if you are not sure to gain significant browsing functionality (e.g., hypertext links in HTML or PDF). This way you will not spend times with versions, file format and platform issues that certain word processors and CASE tools do not succeed to address.

**Hints.** The fact that you are limited in time should force you to think how you can extract the most useful information. Below are some hints for things to look out for.

- A table of contents gives you a quick overview of the structure and the information presented.
- Version numbers and dates tell you how up to date the documentation is.
- References to other parts of the documentation convey chronological dependencies.
- Figures are a always a good means to communicate information. A list of figures, if present, may provide a quick path through the documentation.
- Screen-dumps, sample print-outs, sample reports, command descriptions, reveal a lot about the functionality provided in the system.
- Formal specifications, if present, usually correspond with crucial functionality.
- An index, if present contains the terms the author deems significant.

## **Forces Resolved**

- **Limited Resources.** By applying this pattern, you spend 1/2 a day (plus the time to collect the documentation) to end up with a short list that is a reasonable basis for planning further reengineering efforts.
- **Tools and Techniques.** As reading the documentation only requires the physical document, the tool interference is really low. Yet, when CASE tools have been applied, it may be necessary to consult the documentation on line. Note that CASE tools often enforce some documentation conventions so be sure to be aware of them.

No special techniques are necessary to apply this pattern, unless formal specification or special diagrams are used.

- **Reliable Info.** The success of this pattern depends heavily on the quality of the documentation. Applying this pattern (especially combined with READ ALL THE CODE IN ONE HOUR (p. 115)), you may have encountered some contradicting pieces of information, which is definitely worthwhile to explore in further depth.
- **Abstraction.** The abstraction level you get out depends largely on the abstraction level of the available documentation, but is usually quite high because documentation is supposed to be written at a certain abstraction level.
- **Sceptic Colleagues.** Unless good documentation is available, sceptics will almost certainly consider this activity a waste of time and you will probably loose some credibility with the faithful and fence sitters. This is a negative effect, so reduce its potential impact by limiting the time spend here.

## Rationale

Knowing what functionality is provided by the system is essential for reverse engineering. Documentation provides an excellent means to get an external description of this functionality.

However, documentation is either written before or after implementation, thus likely to be out of sync with respect to the actual software system. Therefore, it is necessary to record the reliability. Moreover, documentation comes in different kinds, i.e. requirement documents, technical documentation, end-user manuals, tutorial notes. Depending on the goal of your reengineering project, you will record the usability of each of these documents. Finally, documentation may contain large volumes of information thus reading is time consuming. By limiting the time you spend on it, you force yourself to classify the pieces of information into the essential and the less important.

## **Related Patterns**

You may or may not want to SKIM THE DOCUMENTATION (p. 118) before READ ALL THE CODE IN ONE HOUR (p. 115) depending on whether you want to keep your mind free or whether you want some subjective input before reading the code. INTERVIEW DURING DEMO (p. 121) can help you to collect a list of entities you want to read about in the documentation.

**Resulting Context.** This pattern results in a list of (i) the important requirements; (ii) the important features (iii); the important constraints; (iv) references to relevant design information plus an opinion on how reliable and useful each of these are. Together with the result of READ ALL THE CODE IN ONE HOUR (p. 115) and SKIM THE DOCUMENTATION (p. 118) this is a good basis to CONFER WITH COLLEAGUES (p. 156) and then proceed with GUESS OBJECTS (p. 127) and CHECK THE DATABASE (p. 130).

# INTERVIEW DURING DEMO

Author(s): Serge Demeyer, Stéphane Ducasse and Sander Tichelaar

## Intent

Obtain an initial feeling for the functionality of a software system by seeing a demo and interviewing the person giving the demo.

**Example.** You are asked to extend an existing database application so that it is now accessible via the world-wide web. To understand how the end-users interact with the application, you will ask one of the current users to show you the application and use that opportunity to chat about the systems user-interface. And to understand some of the technical constraints, you will also ask one of the system maintainers to give you a demo and discuss about the application architecture.

## Context

You are starting a reverse engineering project of a large and unfamiliar software system. You have found somebody to demonstrate the system and explain its usage.

## Problem

You need an idea of the typical usage scenario's plus the main features of a software system in order to plan further reverse engineering efforts.

## Solution

Observe the system in operation by seeing a demo and interviewing the person who is demonstrating. Note that the interviewing part is at least as enlightening as the demo.

After this demo, take about the same time to produce a report about your findings, including (i) some typical usage scenarios or use cases; (ii) the main features offered by the system and whether they are appreciated or not; (iii) the system components and their responsibilities; (iv) bizarre anecdotes that reveal the folklore around using the system.

**Hints.** The person who is giving the demo is crucial to the outcome of this pattern so care is demanded when selecting the person. Therefore, consider to apply this pattern several times with different kinds of persons giving the demo. This way you will see variances in what people find important and you will hear different opinions about the value of the software system. Always be wary of enthusiastic supporters or fervent opponents: although they will certainly provide relevant information, you must spend extra time to look for complementary opinions in order to avoid prejudices.

Below are some hints concerning people you should be looking for, what kind of information you may expect from them and what kind of questions you should ask them.

• An *end-user* should tell you how the system looks like from the outside and explain you some detailed usage scenarios based on the daily working practices. Ask about the situation in the company before the software system was introduced to assess the scope of the software system within the business processes. Probe for the relationship with the computer department to divulge bizarre anecdotes.

- A person from the *maintenance/development team* should clarify the main requirements and architecture of a system. Inquire how the system has evolved since delivery to reveal some of the knowledge that is passed on orally between the maintainers. Ask for samples of bug reports and change requests to assess the thoroughness of the maintenance process.
- A *manager* should inform you how the system fits within the rest of the business domain. Ask about the business processes around the system to check for unspoken motives concerning your reverse engineering project. This is important as reverse engineering is rarely a goal in itself, it is just a means to achieve another goal.

### **Forces Resolved**

- **Limited Resources.** By applying this pattern, you spend 1/2 a day (plus the time to set-up the demo) to end up with a short list that is a reasonable basis for planning further reengineering efforts.
- **Tools and Techniques.** Except for the equipment necessary to run the software system –which should be readily available– this pattern does not require anything special. The interviewing technique to apply requires a special listening ear though.
- **Reliable Info.** A demo is a reliable means to dig out what features are considered important, but you cannot trust on it to omit irrelevant features. Of course the reliability of the information obtained depends largely on the person who is giving the demo. Therefore, if possibly cross-check any information against other more reliable sources (requirements, progress and delivery reports, source code, log files, ...).
- **Abstraction.** The abstraction level achieved by seeing a demo is quite abstract, though it depends on the person who is giving a demo.
- **Sceptic Colleagues.** The users and maintainers of a software system are usually quite eager to show you the system and tell you what they like and dislike about it. If you have a good listening ear this is a good way to boost your credibility.

## Rationale

Interviewing people working with a software system is essential to get a handle on the important functionality and the typical usage scenario's. However, asking predefined questions does not work, because in the initial phases of reverse engineering you do not know what to ask. Merely asking what people like about a system will result in vague or meaningless answers. On top of that, you risk getting a very negative picture because people have a tendency to complain.

Therefore, hand over the initiative to the user by requesting for a demo. First of all, a demo allows users to tell the story in their own words, yet is comprehensible for you because the demo imposes some kind of tangible structure. Second, because users must start from a running system, they will adopt a more positive attitude explaining you what works. Finally, during the course of the demo, you can ask lots of precise questions, getting lots of precise answers, this way digging out the expert knowledge about the system's usage.

### **Known Uses**

One anecdote from the very beginning of the FAMOOS project provides a very good example for the potential of this pattern. For one of the case studies —a typical example of a 3-tiered application with a database layer, domain objects layer and user-interface layer— we were asked 'to get the business objects out'. Two separate individuals were set to that task, one took a source code browser and a CASE tool and extracted some class diagrams that represented those business objects. The other installed the system on his local PC and spend about an hour playing around with the user interface to came up with a list of ten questions about some strange observations he made. Afterwards, a meeting was organized with the chief analyst-designer of the system and the two individuals that tried to reverse engineer the system. When the analyst-designer was confronted with the class-diagrams he confirmed that these covered part of his design, but he couldn't tell us what was missing nor did he tell us anything about the rationale behind his design. It was only when we asked him the ten questions that he launched off into a very enthusiastic and very detailed explanation of the problems he was facing during the design — he even pointed to our class diagrams during his story! After having listened to the analyst-designer, the first reaction of the person that extracted the class diagrams from the source code was 'Gee, I never read that in the source code'.

## **Related Patterns**

For optimum results, you should perform several attempts of INTERVIEW DURING DEMO (p. 121) with different kinds of people. Depending on your taste, you may perform these attempts before, after or interwoven with READ ALL THE CODE IN ONE HOUR (p. 115) and SKIM THE DOCUMENTATION (p. 118).

**Resulting Context.** This pattern results in (i) some typical usage scenarios or use cases; (ii) the main features offered by the system and whether they are appreciated or not; (iii) the system components and their responsibilities; (iv) bizarre anecdotes that reveal the folklore around using the system. Together with the result of READ ALL THE CODE IN ONE HOUR (p. 115) and SKIM THE DOCUMENTATION (p. 118) this is a good basis to CONFER WITH COLLEAGUES (p. 156) and then move on to GUESS OBJECTS (p. 127) and CHECK THE DATABASE (p. 130).

## Chapter 5

# **Cluster: Extract Architecture**

The patterns in FIRST CONTACT (p. 113) should have helped you getting an initial feeling of the software system. Now is the right time to draw some blueprints of the complete system that will serve as a roadmap during the rest of the reverse engineering project. The main priority in this stage of reverse engineering is to get an accurate picture without spending too much time on the hairy details.

The patterns in this cluster tell you how to derive a system blueprint from source code (GUESS OBJECTS (p. 127)) and from a database schema (CHECK THE DATABASE (p. 130)). With these blueprints you will probably want to proceed with FOCUS ON HOT AREAS (p. 133).

### **Forces Revisited**

- **Reliable Info.** Since the blueprints resulting from these activities will influence the rest of your project, accuracy is the single most important aspect. *Consequently, take special precautions to make the extracted blueprints as reliable as possible.* In particular, plan for an incremental approach where you gradually improve the blueprints while you gain a better understanding of the system.
- Limited Resources. Results coming from this stage of reverse engineering are always worthwhile. *Consequently, consider* EXTRACT ARCHITECTURE *a very important activity and plan to spend a considerable amount of your resources here.* However, via an incremental approach you can stretch your resources in time.
- **Tools and Techniques.** While extracting an architecture, you can afford the time and money to apply some heavyweight techniques and purchase some expensive tools. *Yet —because accuracy is so important— do never rely on techniques and tools and always make a conscious assessment of their output.*
- **Abstraction.** Architectural blueprints are meant to strip away the details. Yet, computer science has this strange phenomenon that details are crucial to the overall system [Bro87]. *Consequently, favor different blueprints that emphasize one perspective and choose the most appropriate ones when necessary.* Adapt the notation to the kind of blueprint you are making ([Dav95] principle 21).
- **Sceptic Colleagues.** Good blueprints help a lot because they greatly improve the communication within a team. However, since they strip away details, you risk to offend those people who spend their time on these details. Also, certain notations and diagrams may be new to people, and then your diagrams will just be ignored. *Consequently, take care in choosing which blueprints to produce and which notations to use they should be helpful to all members of the team.*

	Limited Resources	Tools and Techniques	Reliable Info	Abstraction	Sceptic Colleagues
GUESS OBJECTS	-	-	++	++	+
CHECK THE DATABASE	-	_	++	+	++

Table 5.1: How each pattern of FIRST CONTACT resolves the forces. Very good: ++, Good: +, Neutral: 0, Rather Bad: -, Very bad: --

# GUESS OBJECTS

Author(s): Serge Demeyer, Stéphane Ducasse and Sander Tichelaar

## Intent

Progressively refine a model of a software system, by defining hypotheses about what should be in the code and checking these hypotheses against the source code

**Example.** You are facing a 500 K lines C++ program, implementing a software system to display multi-media information in real time. Your boss asks you to look how much of the source code can be resurrected from another project. After having READ ALL THE CODE IN ONE HOUR (p. 115), you noticed an interesting piece of code concerning the reading of the signals on the external video channel. You suspect that the original software designers have applied some form of observer pattern, and you want to learn more about the way the observer is notified of events. You will gradually refine your assumption that the class VIDEOCHANNEL is the subject being observed by reading its source code and tracing interesting paths.

## Context

You are in the early stages of reverse engineering a software system: you have an initial understanding of its functionality and you are somewhat familiar with the main structure of its source code. Due to this, you have identified a certain aspect of the system as especially important. You have on-line access to the source code of the software system and the necessary tools to manipulate it (i.e., from an elementary grep to a professional browser). You have reasonable expertise with the implementation language being used.

## Problem

You must gain an overall understanding of the internal structure of a software system and report this knowledge to your colleagues so that they will use it as a kind of roadmap for later activities.

## Solution

Take a notepad and/or sketchpad (not necessarily as an electronic tool). Based on your experience, and the little you already understand from the system, devise a model that serves as your initial hypotheses of what to expect in the source code. Check these hypotheses against the source code, using whatever tools you have available. Consciously keep track of which parts of the source code confirm and which parts contradict your hypotheses. Based on the latter, refine the initial model, recheck the hypotheses and rework the list of confirmations and contradictions. Do this until you obtain a more or less stable model.

Note that it is a good idea to sort the entities in your hypotheses models according to the probability of appearance in source-code. This is especially useful as names inside the source-code do not always match with the concepts they represent. This may be due to particular coding conventions or compiler restrictions (identifiers cannot exceed a certain length), or because of the native language of the original programmer.<sup>1</sup>

Afterwards, sit down to produce a boxes- and arrows diagram describing your findings. As a rule of the thumb, make sure your diagram fits on one page. It is better to have two distinct diagrams, where each

<sup>&</sup>lt;sup>1</sup>In one particular reverse engineering experience, we were facing source code that was a mixture of English and German. As you can imagine, grep is not a very good tool to check occurrences of English terms in German texts.

provides a clean perspective on the system than one messy diagram with too many details too read and memorize. People should be able to redraw the diagram from memory after they have seen it once; it is only then that your diagram will really serve as a roadmap.

**Variations.** The pattern itself is quite broad and thus widely applicable. Below are some suggestions of possible variants.

- *Guess Patterns*. While having READ ALL THE CODE IN ONE HOUR (p. 115), you might have seen some symptoms of patterns. You can use a variant of GUESS OBJECTS to refine this knowledge. (See the better known pattern catalogues [GHJV95], [BMR<sup>+</sup>96], [Fow97a] for patterns to watch out for. See also [Bro96] for a discussion on tool support for detecting patterns.)
- *Guess Object Responsibilities*. Based on the requirements resulting from SKIM THE DOCUMENTA-TION (p. 118), you can try to assign object responsibilities and check the resulting design against the source code. (To assign object responsibilities, use the noun phrases in the requirements as the initial objects and the verb phrases as the initial responsibilities. Derive a design by mapping objects on class hierarchies and responsibilities on operations. See [WBWW90] for an in depth treatment on responsibility-driven design.)
- *Guess Object Roles.* The usage scenarios that you get out of INTERVIEW DURING DEMO (p. 121) may serve to define some use cases that in turn help to find out which objects fulfill which roles. (See [JCJO92] for use cases and [Ree96] for role modeling.)
- *Guess Process Architecture*. The object-oriented paradigm is often applied in the context of distributed systems with multiple cooperating processes. A variant of GUESS OBJECTS may be applied to infer which processes exist, how they are launched, how they get terminated and how they interact. (See [Lea96] for some typical patterns and idioms that may be applied in concurrent programming.)

### **Forces Resolved**

- **Limited Resources.** The amount of resources you invest in this pattern depends mainly on the level of detail and accuracy that you want to achieve. Be wary of the hairy details though, as this pattern tends to have an exponential effort/gain curve. For detailed information, consider switching to STEP THROUGH THE EXECUTION (p. 146) instead.
- **Tools and Techniques.** Applying this pattern does not require a lot of tools: a a simple grep may be sufficient and otherwise a good code browser will do. Probably you will also need a tool for producing the final blueprint, as it is likely that someone will have to update the blueprint later on in the project. However, choose a a simple drawing tool rather then a special purpose CASE tool, as you will need a lot of freedom to express what you found.

In itself, the pattern does not require a lot of techniques. However, a large repertoire of knowledge about idioms, patterns, algorithms, techniques is necessary to recognize what you see. As such, the pattern should preferably be applied by experts, yet lots of this expertise may be acquired on the job.

- **Reliable Info.** The blueprints you extract by applying this pattern are quite reliable because of the gradual refinement of the hypotheses and confirmation against source code. Yet, be sure to keep the blueprint up to date while your reverse engineering project progresses and your understanding of the software system grows.
- **Abstraction.** If applied well, the different blueprints you achieve by means of GUESS OBJECTS provide the ideal abstraction level. That is, each blueprint provides a unique perspective on the software system that highlights the important facts and strips the unimportant details. Yet, navigating between the various blueprints provides you all the necessary perspectives to really understand the system.

**Sceptic Colleagues.** The results of GUESS OBJECTS pattern should drastically increase the confidence of your team in the success of the reverse engineering project. This is because the members of the team will normally experience an "*aha erlebness*", where the little pieces of knowledge they have fit the larger whole.

### Rationale

Clear and concise descriptions of a system are a necessary ingredient to plan team activities. However, being clear and concise is for humans to decide, thus creating them requires human efforts. On the other hand, they must accurately reflect what's inside the system, so somehow the source-code should be incorporated in the creation process as well. GUESS OBJECTS addresses this tension by using a mental model (i.e., the hypotheses) as the primary target, yet progressively refines that model by checking it against source code. Moreover, conciseness implies loss of detail, hence the reason to extract multiple blueprints offereing alternative perspectives.

### **Known Uses**

In [MN97], there is a report of an experiment where a software engineer at Microsoft applied this pattern (it is called 'the Reflexion Model' in the paper) to reverse engineer the C-code of Microsoft Excel. One of the nice sides of the story is that the software engineer was a newcomer to that part of the system and that his colleagues could not spend too much time to explain him about it. Yet, after a brief discussion he could come up with an initial hypotheses and then use the source code to gradually refine his understanding. Note that the paper also includes a description of a lightweight tool to help specifying the model, the mapping from the model to the source code and the checking of the code against the model.

## **Related Patterns**

All the patterns in the FIRST CONTACT (p. 113) cluster are meant to help you building the initial hypotheses to be refined via GUESS OBJECTS (p. 127). Next, some of the patterns in FOCUS ON HOT AREAS (p. 133) may help you to refine this hypothesis.

**Resulting Context.** After this pattern, you will have a series of blueprints where each contains one perspective on the whole system. These blueprints will help you during later reverse engineering steps, in particular the ones in FOCUS ON HOT AREAS (p. 133) and PREPARE REENGINEERING (p. 149). Consequently, consider applying CONFER WITH COLLEAGUES (p. 156) after applying GUESS OBJECTS (p. 127).

# CHECK THE DATABASE

Author(s): Serge Demeyer, Stéphane Ducasse and Sander Tichelaar

### Intent

Get a feeling for the data model inside a software system by checking the database schema.

**Example.** You are asked to extend an existing database application so that it is now accessible via the world-wide web. The initial software system manipulates the business objects (implemented in C++) stored inside a relational database. You will reconstruct the data model underlying your business objects by mapping the table definitions in the database on the corresponding C++ classes.

### Context

You are in the early stages of reverse engineering a software system, having an initial understanding of its functionality. The software system employs some form of a database to make its data persistent.

You have access to the database and the proper tools to inspect its schema. Or even better, you have samples of data inside that database and maybe you are even able to spy on the database queries during the execution of the system. Finally, you have some expertise with databases and knowledge of how data-structures from your implementation language are mapped onto the data-structures of the underlying database.

## Problem

You want to derive a data model for the persistent data in a software system in order to guide further reverse engineering efforts.

### Solution

Check the database schema to reconstruct at least the persistent part of the data model. Use your knowledge of how constructs in the implementation language are mapped onto database constructs to reverse engineer the real data model. Make samples of data inside the database to refine the data-model.

### **Forces Resolved**

- Limited Resources. Reconstructing the data model from the database schema takes considerable resources, although it depends largely on the underlying technology. Factors that affect this force in a positive way are the quality of the database schema (is it in normal form?), the correspondence between the database paradigm and the implementation language paradigm (inheritance hierarchies do not map directly to relational tables), the expressiveness of the database schema (does it include foreign keys ?). On the other hand, the reverse engineering of database schemas may include techniques like data sampling and run-time inspection, which takes even more resources.
- **Tools and Techniques.** This pattern can do without a lot of tool support: a dump of the database schema and some samples of data inside the tables is something all database systems can provide. However,

there are some tools available to support you in recovering object models (see [HEH<sup>+</sup>96], [PB94], [JSZ97]).

This pattern requires substantial technical expertise, because it requires knowledge of ways to manipulate data structures in both the implementation language and the database, plus ways to map one onto the other.

- **Reliable Info.** Because the pattern is based on analyzing persistent data, the reliability of the reconstructed data model is usually quite high. However, if the database system is manipulated by different software systems and if each of these software systems is build with different implementation technologies (CASE tools, 4GL, ...), the reliability of the data model tends to decrease because the database schema provides the most common denominator of all implementation technologies involved. Data sampling is a good way to cope with this problem though.
- **Abstraction.** The abstraction level of the reconstructed data model tends to be low, as it is closer to the underlying database schema than it is to the implementation language. However, this depends largely on the amount of resources spent. For instance, with data sampling and run-time inspection one can drastically improve the abstraction level.
- **Sceptic Colleagues.** If applied well, this pattern increases your credibility considerably, because a well defined data model is normally considered a collective source of knowledge which greatly improves the communication within a team. Moreover, almost all software engineers will have experience with data models and will appreciate their presence.

## Rationale

Having a well defined central data model is a common practice in larger software projects that deal with persistent data. Not only, it specifies common rules on how to access certain data structures, it is also a great aid in assigning development tasks. Therefore, it is a good idea to extract an accurate data model before proceeding with other reverse engineering activities.

## **Known Uses**

The reverse engineering and reengineering of database systems is a well-known problem, drawing certain attention in the literature (see [HEH+96], [PB94], [JSZ97]). Note the recurring remark that the database schema alone is too weak a basis and that data sampling and run-time inspection must be included for successful reconstruction of the data model.

## **Related Patterns**

CHECK THE DATABASE requires an initial understanding of the system functionality, like obtained by applying patterns in the cluster FIRST CONTACT (p. 113).

There are some patterns that describe various ways to map object-oriented data constructs on relational database counterparts. See among others [KC98], [CKR99].

**Resulting Context.** CHECK THE DATABASE results in a data model for the persistent data in your software system. Such a data model is quite rough, but it may serve as an ideal initial hypotheses to be further refined by applying GUESS OBJECTS (p. 127). The data model should also be used as a collective knowledge that comes in handy when doing further reverse engineering efforts, for instance like in the clusters FOCUS ON HOT AREAS (p. 133) and PREPARE REENGINEERING (p. 149). Consequently, consider to CONFER WITH COLLEAGUES (p. 156) after CHECK THE DATABASE.

## Chapter 6

# **Cluster: Focus on Hot Areas**

The patterns in FIRST CONTACT (p. 113) should have helped you getting an initial feeling of the software system, while the ones in EXTRACT ARCHITECTURE (p. 125) should have aided you deriving some blueprints of the overall system structure. The main priority now is to get detailed knowledge about a particular part of the system.

This cluster tell you *how*, and to some degree *where* you might obtain such detailed knowledge. The patterns involve quite a lot of tools and rely on substantial technical knowledge, hence are applicable in the later stages of a reverse engineering project only. Indeed, only then you can afford to spend the resources obtaining detailed information as only then you have the necessary expertise to know that your investment will pay off.

There are two patterns that explain you *where* to focus your attention: INSPECT THE LARGEST (p. 135) suggests to look at large object objects, while EXPLOIT THE CHANGES (p. 139) advises to look at the places where programmers have been changing the system. (Of course, no technique or tool will replace the human mind, hence to know where to focus your attention, be sure to CONFER WITH COLLEAGUES (p. 156) as well). Then, there are two patterns that inform you *where and how* to study program structures: VISUALIZE THE STRUCTURE (p. 142) tells about program visualisation techniques, while CHECK METHOD INVOCATIONS (p. 144) recommends to check invocations of both constructor and overridden methods. Finally, there is one pattern describing you *how* to investigate programs, namely STEP THROUGH THE EXECUTION (p. 146) which explains how to take advantage of your debugger.

Many reverse engineering projects prepare for a subsequent reengineering phase. If you're in such a situation, you might consider the patterns in PREPARE REENGINEERING (p. 149) as your next step. If you're not, then these patterns are the last ones we have to offer for helping you.

## **Forces Revisited**

- **Tools and Techniques.** To obtain the required details from a software system you must pay the price in terms of technical expertise and tools. This is the most important force during this stage of reverse engineering: *consequently, make sure your reverse engineering team does possess the necessary skills and tools.*
- Limited Resources. These patterns are applicable during the later stages of a reverse engineering project, thus resources are less scarce as you can be quite sure that your investment will pay off. On the other hand, the activities you apply require more resources. *Consequently, engage in detailed reverse engineering only when you are certain that you need to know the details about that part of a system.* The patterns in the previous clusters should have helped you obtaining that knowledge.

- **Abstraction.** All patterns in this cluster have in common that they extract detailed information, at an intermediate level of abstraction (i.e., between source code and design). Yet, detailed knowledge is necessary because in software engineering —and this is in contrast with many other engineering disciplines— details are very important [Bro87]. So, even during fine-grained reverse engineering, there are little details that seem so obvious, yet may obstruct the understanding of the system if you failed to state them.<sup>1</sup> *Consequently, when working on intermediate abstraction levels, make sure you provide enough context so that the relationship with both higher and lower levels is clear.*
- **Reliable Info.** As details are so important, you should be confident in the obtained results. *Consequently, favour extracting information from the trustworthy information sources.* Fortunately, because you're in the later stages of reverse engineering, you know which information sources are reliable and which ones are not.
- **Sceptic Colleagues.** You would not have arrived this far without the support of some colleagues, so at least you still have the support of the faithful. Moreover, you probably did satisfy the expectations, otherwise the sceptic would have succeeded to cancel your project. And if you did really well, you might even have won some fence sitters over into the camp of the faithful. At this stage, you will not achieve more support from your colleagues. *Consequently, keep on delivering the necessary results to avoid providing reasons for the sceptics to cancel your project.*

	Limited Resources	Tools and Techniques	Reliable Info	Abstraction	Sceptic Colleagues
INSPECT THE LARGEST	0	-	0	0	0
VISUALIZE THE STRUCTURE	-		+	+	+
CHECK OVERRIDDEN METHODS	-	-	+	+	0
EXPLOIT THE CHANGES			++	+	++
STEP THROUGH THE EXECUTION	-	0	++	-	+

Table 6.1: How each pattern of FIRST CONTACT resolves the forces. Very good: ++, Good: +, Neutral: 0, Rather Bad: -, Very bad: --

<sup>&</sup>lt;sup>1</sup>A typical example of such a harmful detail is the use of private/protected in a UML diagram. Depending on the favourite programming language of the author of the diagram, the interpretation is quite different, and readers of the diagram should be made aware of this. That is, with a C++ background the interpretation is class based, thus instances of the same class may access each other's private members. On the other hand, with a Smalltalk background, the interpretation is instance based, thus it is only the object itself that is allowed to access its members. Finally, in Java a protected member may also be accessed by classes in the same package.

# INSPECT THE LARGEST

Author(s): Serge Demeyer, Stéphane Ducasse and Sander Tichelaar

### Intent

Identify important functionality by looking at large constructs.

**Example.** You are facing an object-oriented system and you want to find out which classes do the bulk of the work. You will produce a list of all classes where the number of methods exceeds the average number of methods per class, sort the list and inspect the largest classes manually.

## Context

You are in a later stage of reverse engineering a software system. You have an overall understanding of its functionality and you know the main structure of its source code. You have a metrics tool at your disposal plus a code browser to inspect the source code. The metrics tool is configured in such a way that it provides you with a number of measurements of source code constructs you are feeding into it. Moreover, the metrics are defined in such a way that they have a high correlation with the amount of functionality implemented in the construct.

## Problem

You must identify those places in the source code that correspond with important chunks of functionality.

## Solution

Use the metrics tool to collect a limited set of measurements for all the constructs in the system. Sort the resulting list according to these measurements. Browse the source code for the largest among those constructs in order to understand how these constructs work together with other related constructs. Produce a list of all the constructs that appear important, including a description of how they should be used (i.e. external interface).

**Hints.** Identifying important pieces of functionality in a software system via measurements is a tricky business which requires expertise in both data collection and interpretation. Below are some hints that might help you getting the best out of your data.

• Which metrics to collect ? In general, it is better to stick to the simple metrics, as the more complex ones will not perform better for the identification of large constructs. This experience is backed up by empirical evidence, as it has been reported in the literature that size metrics have a high correlation (see among others [FP97]).

For identifying important functionality in object-oriented source code, look at methods and classes. For methods you may restrict yourself to counting the lines of code and if available the number of other methods invoked.<sup>2</sup> For classes, you should count the number of methods and the number of

<sup>&</sup>lt;sup>2</sup>Counting the lines of code can be done very efficiently without parsing, just by counting all occurrences of the <CR> character.

attributes defined on that class, plus the depth of the inheritance tree and probably also the lines of code (i.e., the sum of all the lines of code of all the methods of a class). (See the chapter on Metrics — p.22 for a more precise definition of each of the metrics and a list of other metrics you might collect).

- Which variants to use ? Usually, it does not make a lot of difference which variant is chosen, as long as the choice is clearly stated and applied consistently. Here as well, it is preferable to choose the most simple variant, unless you have a good reason to do otherwise. For instance, while counting the lines of code, you should decide whether to include or exclude comment lines, or whether you count the lines after the source code has been normalised via pretty printing. In such a case, do not exclude comment lines nor normalise the source code as the extra effort will not pay of. Another example of an alternative definition is the case of counting the number of methods, where one must decide how to deal with 'special' methods like class methods (i.e., the C++ static methods). In this case it is a good idea to count class methods separately as they represent a different kind of functionality.
- *Which thresholds to apply*? It is better not to apply thresholds to filter out those constructs which measurements fall into a given threshold interval. Indeed, 'large' is a relative notion and thresholds will distort your perspective of what constitutes large within the system.
- *How to interpret the results*? Do not only look for the largest construct while analysing the data. Before actually browsing the source code, check the distribution of measurements to see whether the 80/20 rule is satisfied. Also, gather several measurements in different columns one beside another and then look for unusual rows.

(Note that the 80/20 rule is a more formal expression of the rule of the thumb that most constructs in source code will be small, and only a few exceptional cases will be large. To be precise, the rule states that 80% of the constructs will be smaller than 20% of the size of the largest construct.)

## **Forces Resolved**

- Limited Resources. Once the metric tool is configured for the particular language of the software system, collecting the necessary data is not that resource consuming; in the worst case it can be done via batch jobs during the night. However, analysing the data and browsing the selected source code constructs requires a lot of resources depending on the desired level of detail. You can neutralise this effect to some extent by limiting the set of metrics.
- **Tools and Techniques.** To apply this pattern, you require a tool which should be able to collect the necessary measurements. However, since you can restrict yourself to simple counting metrics such a tool should be quite easy to obtain.

Analysing and interpreting the data however, requires a certain amount of knowledge. Some of this knowledge is summarised in our list of metric definitions (see the chapter on Metrics — p.22) and the rest you can learn on the job.

- **Reliable Info.** It is not because a software construct is large that it is important, neither is it true that a small construct is always irrelevant. Therefore, the results contain quite a lot of noise, hence are somewhat unreliable. Still, given the amount of resources required, this pattern usually provides a good return on investment, especially since the large constructs will often point you to other more important but smaller constructs.
- **Abstraction.** The abstraction level of this pattern results mainly from browsing the source code, not so much from measuring the constructs in the system. Therefore, the abstraction level of the results should be considered quite low.
- Sceptic Colleagues. Metrics are often associated with process and quality control, therefore some programmers may believe that you will use the metrics to examine their productivity. Be careful if

you have such programmers among your faithful as it may be a way to turn them into sceptics. In particular, do not blindly deduce that large constructs are bad and should be rewritten.

### Rationale

The main reason why size metrics are often applied during reverse engineering is because they provide a good focus (between 10 to 20% of the software constructs) for a low investment, even though the results are somewhat unreliable. With such a good focus, you can afford some erroneous results which you will compensate anyway via code browsing.

The results are a bit unreliable because 'large' is not necessarily the same as 'important'. Quite often large constructs are irrelevant as they would have been refactored into smaller pieces if they were important. Conversely, small constructs may be far more important than the large ones, because good designers tend to distribute important functionality over a number of highly reusable and thus smaller components. Still, different larger constructs may share the same smaller construct, so via the larger constructs you may be able to identify some important smaller constructs too.

The main disadvantage of the pattern is that it forces you to look at the largest constructs first. Large constructs are usually the most complicated ones, therefore understanding the corresponding source code may prove to be difficult. Another disadvantage is that the analysis of the metrics data results in a list of raw software constructs. For program understanding, it is usually more important to know how these constructs work together with other ones, something which must be revealed by code browsing.

Note that by restricting yourself to a limited set of simple metrics you already avoid one of the most common pitfalls. Indeed, metrics tools usually offer you a wide range of metrics and since collecting data is so easy, you may be tempted to apply all metrics that are available. However, the more data you collect, the more data you must analyse and the amount of numbers will quickly overwhelm you. Moreover, some metrics require substantial parsing effort, which in turn requires the configuration of the parser to your software system, which can be painstaking and time-consuming. By limiting the amount a metrics and keeping the metrics simple, you circumvent these problems.

## **Known Uses**

In several places in the literature it is mentioned that looking for large object constructs helps in program understanding (see among others, [MLM96a], [Kon97], [FNP98a], [FNP98b], [Mar98], [LS98], [Nes88]). Unfortunately, none of these incorporated an experiment to count how much important functionality remains undiscovered. As such it is impossible to assess the reliability of size metrics for reverse engineering.

Note that some metric tools visualise information via typical algorithms for statistical data, such as histograms and Kiviat diagrams. Visualisation may help to analyse the collected data. Datrix [MLM96a], TAC++ [FNP98a], [FNP98b], and Crocodile [LS98] are tools that exhibit such visualisation features.

### **Related Patterns**

Looking at large constructs requires little preparation but the results are a bit unreliable. By investing more in the preparation you may improve the reliability of the results. For instance, if you VISUALIZE THE STRUCTURE (p. 142) you invest in program visualisation techniques to study more aspects of the system in parallel, thereby increasing the quality of the outcome. Also, you can EXPLOIT THE CHANGES (p. 139) to focus on those parts of the system that change, thereby increasing the likelihood of identifying interesting constructs and focussing on the way constructs work together.

**Resulting Context.** By applying this pattern, you will have identified some constructs representing important functionality. Some other patterns may help you to further analyse these constructs. For instance, if you VISUALIZE THE STRUCTURE (p. 142) you will obtain other perspectives and probably other insights as well. Also, if you STEP THROUGH THE EXECUTION (p. 146) you will get a better perception of the runtime behaviour. Finally, in the case of a object-oriented code, you can CHECK METHOD INVOCATIONS (p. 144) to find out how a class is related to other classes.

Even if the results have to be analysed with care, some of the larger constructs can be candidates for further reengineering: large methods may be split into smaller ones (see [Fow99]), just like big classes may be cases of a GOD CLASS (see [BMMM98]).

# EXPLOIT THE CHANGES

Author(s): Serge Demeyer, Stéphane Ducasse and Sander Tichelaar

## Intent

Recover design issues by asking where, how and why the developers have been changing the implementation.

**Example.** You must understand an old but evolving software system, where the evolution is controlled through a configuration management system. You will filter out those modules that have been changed most often and find out what these changes where about and why they were necessary.

**Example.** You must understand an object-oriented framework that has been adapted several times as the developers gained insight into the problem domain. You will filter out all classes where the number of methods and attributes has decreased significantly and find out where that functionality has been moved to. With that knowledge, you will make a guess at the design rationale underlying this redistribution of functionality.

## Context

You are in a later stage of reverse engineering an evolving software system. You have an overall understanding of its functionality and you know the main structure of its source code. You have several releases of the source code at your disposal plus a way to detect the differences between the releases, i.e. a configuration management system and/or a metrics tool.

## Problem

You must identify those parts in the design that played a key role during the system's evolution.

## Solution

Use whatever means at your disposal to compile a list of targets of important/frequent changes. For each target, put yourself in the role of the original developer and ask yourself what the change is about and why it was necessary. With this insight, produce a list of crucial system parts, including a description of the design issues that makes them important.

**Variations.** The pattern comes in two variants corresponding to the way the targets of changes are identified.

• *The configuration database variant* requires that all changes to the system were done via a configuration management system which logs all changes in the configuration database. In that case you can take advantage of the query facilities provided by the configuration database to produce a list of components that have been changed. Sort the list according to the frequency of changes and inspect the corresponding source code plus the comments in the configuration database to find out how and why this component has changed.

• *The change metrics variant* identifies changes by comparing subsequent releases and measuring differences in size. With the change metrics variant, the first step is to measure the size of named constructs in two subsequent releases. Afterwards, you compile a list with three columns: the name of the construct and both measurements. Sort the list according to the largest decrease in size. For each decrease in size, ask yourself where this functionality has been moved to and then deduce how and why this construct has changed.

**Hints.** If you consider applying the change metrics variant on object-oriented source code, we can recommend three heuristics that help identifying the following changes.

- *Split into superclass / merge with superclass.* Look for the creation or removal of a superclass (change in hierarchy nesting level HNL), together with a number of pull-ups or push-downs of methods and attributes (changes in number of methods NOM and number of attributes NOA).
- Split into subclass / merge with subclass. Look for the creation or removal of a subclass (change in number of immediate subclasses- HNL), together with a number of pull-ups or push-downs of methods and attributes (changes in number of methods NOM and number of attributes NOA).
- *Move functionality to superclass, subclass or sibling class.* Look for removal of methods and attributes (decreases in number of methods NOM and number of attributes NOA) and use code browsing to identify where this functionality is moved to.
- Split method / factor out common functionality. Look for decreases in method size (via lines of code LOC, or number of message sends MSG, or number of statements NOS) and try to identify where that code has been moved to.

## Rationale

A configuration management tool maintains and controls the different versions of the components that constitute the entire software system. If such a tool has been used for the software system you are reverse engineering, its database contains a wealth of information about where, how and why the software system has evolved. As a reverse engineer, you should exploit the presence of this database.

But even without a configuration management system, it is feasible to identify where, how and to some degree why a system has evolved by comparing subsequent releases and measuring changes. With change metrics, the results are less accurate than it is the case with the configuration database variant, mainly because the rationale for the change is not recorded thus must be deduced. On the other hand, because you focus on constructs that decrease in size, you are likely to identify places where functionality has been moved to other locations. Such moving of functionality is always relevant for reverse engineering, as it reveals design intentions from the original developers.

Satisfying the prerequisite of having different releases of the source code plus the necessary tools to assess the differences, the main advantages of looking at changes are the following. (i) It concentrates on relevant parts, because the changes point you to those places where the design is expanding or consolidating. (ii) It provides an unbiased view of the system, because you do not have to formulate assumptions of what to expect in the software (this is in contrast to GUESS OBJECTS (p. 127) and VISUALIZE THE STRUCTURE (p. 142)) (iii) It gives an insight in the way components interact, because the changes reveal how functionality is redistributed among constructs (this is in contrast to INSPECT THE LARGEST (p. 135)).

## **Known Uses**

There is a company called MediaGeniX, which incorporates a scaled down version of the configuration database variant into their development process and tools. It is based on the so-called tagging tool, which

automatically updates one comment line in a method body each time this method is modified. The comment line records information like the date of the change, the name of the programmer and a reference into their configuration management system. The reference reveals the nature of the change (i.e., bug fix or a new feature) and via consultation of the actual configuration management system even what this change was about. Afterwards, they run queries to identify which features are localised to a few modules and which features cross-cut a large number of modules to identify where they may improve the design of the framework. Also, they have identified methods that are modified a lot when bug fixing, and used this information as input for their code reviewing. They have even identified cycles in the bug fixing, in the sense that the modification of one method fixed a bug but immediately introduced another bug and then the repair of the newly introduced bug again introduced the older bug. More information about the usage of the tagging tool in the context of reverse engineering can be found in [Hon98].

Besides the tagging tool, we are aware of two other projects where people have been exploiting the version control system for reverse engineering purposes. First, there is the SeeSoft tool, developed at Bell Labs, which visualises source code changes and has been used successfully for reverse engineering purposes [BE96]. Second, there is the ARES project (see http://www.infosys.tuwien.ac.at/Projects/ARES/) which also experimented with visualisation of changes using the 3DSoftVis tool [JGR99].

Finally, concerning the change metrics variant, we ran an experiment on three medium sized systems implemented in Smalltalk. As reported in [DDN99], these case studies suggest that the heuristics support the reverse engineering process by focussing attention on the relevant parts of a software system.

## **Related Patterns**

Inspecting changes is a costly but very accurate way of identifying areas of interest in a system. If you VISUALIZE THE STRUCTURE (p. 142) or INSPECT THE LARGEST (p. 135) you will get less accurate results for a lower amount of resources.

**Resulting Context.** By applying this pattern, you will have identified some parts in the design that played a key role during the system's evolution. Some other patterns may help you to further analyse these constructs. For instance, if you VISUALIZE THE STRUCTURE (p. 142) you will obtain other perspectives and probably other insights as well. Also, if you STEP THROUGH THE EXECUTION (p. 146) you will get a better perception of the run-time behaviour. Finally, in the case of a class, you can CHECK METHOD INVOCATIONS (p. 144) to find out how this class is related to other classes.

# VISUALIZE THE STRUCTURE

Author(s): Serge Demeyer, Stéphane Ducasse and Sander Tichelaar

#### Intent

Obtain insight in the software system's structure —including potential design anomalies— by means of well-known visualisations.

**Example.** You want to understand an object-oriented class structure in order to improve it. In particular, you would like to redistribute responsibilities, by splitting large superclasses and hooking the subclasses underneath the appropriate ancestor. To analyse the situation, you will display the inheritance hierarchies, paying special attention to large classes high up in the hierarchy. Afterwards, for classes identified that way, you will display a graph showing which method accesses which attributes to analyse the class' cohesion and find out whether a split is feasible.

### Context

You are in a later stage of reverse engineering a software system. You have an overall understanding of the system's functionality and based on that understanding, you have selected part of the software system for further inspection. You have a program visualisation tool at your disposal plus a code browser to inspect the source code.

## Problem

You want to obtain insight in the structure of a selected part of a software system, including knowledge about potential design anomalies.

### Solution

Instruct the program visualisation tool to show you a series of graphical layouts of the program structure. Based on these graphical layouts, formulate yourself some assumptions and use the code browser to check whether your assumptions are correct. Afterwards, produce a list of correct assumptions, classifying the items in one of two categories: (i) helps program understanding, or (ii) potential design anomaly.

**Hints.** Obtaining insight in the structure of a software system via visualisation tools is difficult, especially when searching for potential design anomalies. We have included our expertise with program visualisation in a separate chapter and we refer the interested reader to the chapter on Visualisation — p.31 for further details.

### Rationale

Program visualisation is often applied in reverse engineering because good visual displays allow the human brain to study multiple aspects of a problem in parallel. This is often phrased as "one picture conveys a

thousand words", but then of course the problem is which words they convey, thus which program visualisations to apply and how to interpret them. For the program visualisations listed in the the chapter on Visualisation — p.31 we describe both when to apply them and how to interpret the results. For other visualisations you will have to experiment to find out when and how to use them.

## **Related Patterns**

If your program visualisation tool scales enough to accommodate the system your facing, then you can start to VISUALIZE THE STRUCTURE right away. However, since program visualisations rarely scale well, it is preferable to first filter out which parts of the source code are relevant for further analysis. Therefore, consider to INSPECT THE LARGEST (p. 135) or to EXPLOIT THE CHANGES (p. 139) before you VISUALIZE THE STRUCTURE.

**Resulting Context.** By applying this pattern, you will have obtained an overview of the structure of a selected part of a software system, including potential design anomalies. Some other patterns may help you to further analyse these constructs. For instance, if you STEP THROUGH THE EXECUTION (p. 146) you will get a better perception of the run-time behaviour and if you CHECK METHOD INVOCATIONS (p. 144) you can find out how a class is related to other classes. If you have identified design anomalies, you should consider to refactor them (see [Fow99]). Some typical design anomalies including the way to refactor them can be found in the part on Reengineering — p.167.

# CHECK METHOD INVOCATIONS

Author(s): Serge Demeyer, Stéphane Ducasse and Sander Tichelaar

### Intent

Find out how a class is related to other classes by checking the invocations of key methods in the interface of that class. Two examples of key methods that are easy to recognise are constructors and overridden methods.

**Example.** You have identified a number of classes that represent part of the domain model. You want to learn about the aggregation relationships between these classes and therefore, you will inspect for all constructor methods which methods are invoking them.

**Example.** You have identified a part of a class hierarchy where the designers relied on template methods to customise the design. To learn how the subclasses interact with their superclasses, you will retrieve all methods overriding another one, and inspect who is invoking these methods.

## Context

You are in a later stage of reverse engineering a software system implemented in an object-oriented language. You have an overall understanding of the system's functionality and based on that understanding, you have selected a part of the class hierarchy for further inspection. You have a code browser at your disposal that allows you to jump from a method invocation to the places where the corresponding method is defined.<sup>3</sup>

## Problem

You want to find out how a class is related to the other classes in the system.

## Solution

Select key methods in the interface of the class and inspect who is invoking these methods.

Variations. The pattern has two variants depending on the selected methods in the public interface.

- *The constructor method variant* suggests you to look at invocations of constructor methods to reveal aggregation relationships between classes.
- *The template method variant* recommends you to select methods that are overridden in a subclass plus the methods invoking them to infer template methods [GHJV95].

 $<sup>^{3}</sup>$ Note that your code browser should take polymorphism into account. Polymorphism implies that one invocation has several candidates for being the defining method. Because the actual target can only be resolved at run-time, your browser must show all candidates.

**Hints.** If you consider applying the above variants, following suggestions may help you getting the best out of your efforts.

- For *the constructor method variant*, you must trace the chain of invocations until the result of the constructor is stored into an attribute. The class defining this attribute is the aggregation. Also, look out for invocations of constructor methods where the invoking object is passing itself as an argument and where this argument is stored into an attribute of the constructor class. In that case, the constructor class is the aggregation.
- *The template method variant*, explicitly states that you should look a methods that are overridden and not methods that are declared abstractly. The reason is that not all template methods distinguish the hook method via an abstract method, but that often a concrete method is used to specify the default behaviour. By looking for overridden methods, you are certain that you will cover the latter case as well.

## Rationale

If the object-oriented paradigm is applied well, state should be encapsulated behind the interface of a class (see [Mey97] and [Bec97] among others). Therefore, to understand how a class is related to other classes, method invocations are more reliable than attribute declarations. Yet, because the amount of method invocations is large you must choose which invocations to analyse. This pattern helps you in this choice by suggesting two specific kinds of methods that are easy to identify and result in well-known class relationships.

## **Known Uses**

In [DG98] we report on a case study where we applied the template method variant.

## **Related Patterns**

Checking method invocations of classes is quite tedious, thus it is best to start with a small amount of classes. Therefore, consider to INSPECT THE LARGEST (p. 135) or EXPLOIT THE CHANGES (p. 139) or VISUALIZE THE STRUCTURE (p. 142) to limit the amount of classes to inspect.

**Resulting Context.** By applying this pattern, you will know how a class is related to the other classes in the system. If you STEP THROUGH THE EXECUTION (p. 146) you will get a better perception of the run-time behaviour of these relationships.

## **STEP THROUGH THE EXECUTION**

### Intent

Obtain a detailed understanding of the run-time behaviour of a piece of code by stepping through its execution.

**Example.** You have a piece of code that implements a graph layout algorithm and you must understand it in order to rewrite it. You will feed a graph into the program and use the debugger to follow how the algorithm behaves.

## Context

You are in a later stage of reverse engineering a software system. You have an overall understanding of the system's functionality and based on that understanding, you have selected part of the software system for further inspection. You have a debugger at your disposal that allows you to inspect data structures and to interactively follow the step by step execution of a piece of code. You know a set of representative input data to feed into that piece of code to launch a normal operation sequence.

## Problem

You want to obtain a detailed understanding of the run-time behaviour of a piece of code.

### Solution

Feed representative input data in the piece of code to launch a normal operation sequence. Use the debugger to follow the step by step execution and to inspect the internal state of the piece of code.

#### Hints.

- Test programs usually provide representative input data in their initialisation code.
- Usage scenarios, like the ones resulting from INTERVIEW DURING DEMO (p. 121), may give clues on what is a normal operation sequence.

## **Forces Resolved**

- **Limited Resources** Once you know what you really want to understand, this pattern works well in a limited resource context. However, stepping through the code can be highly inefficient whitout a clear focus.
- Tools and Techniques The success of this pattern relies on the ability to use a good interactive debugger.
- **Reliable Info** By following the step by step execution of a program, you get a very reliable view of a piece of code. However, beware that the input data is indeed representative for a normal operation sequence.

Abstraction The abstraction level is quite low, unless you can tie the step by step execution to a typical usage scenario.

Sceptic Colleagues Neutral.

## Rationale

In [Mey97], object-oriented programming is defined as "designing a system around the functionality it offers rather then the data structures it operates upon". Hence, understanding the run-time behaviour is crucial to understand an object-oriented program. And the best way to get a view on the run-time behaviour is to see the events as they actually occur in a real execution, a view which is provided by interactive debugging tools.

## **Known Uses**

In [RE93] you can find some interesting debugging techniques applicable in the context of Smalltalk. Many of them will generalise to other programming environments as well.

## **Related Patterns**

Stepping through program executions is quite tedious, thus it is best to focus on a small piece of code. Consider to INSPECT THE LARGEST (p. 135) or to EXPLOIT THE CHANGES (p. 139) or to VISUALIZE THE STRUCTURE (p. 142) to obtain such a focus. Also, you need some typical usage scenarios which may be provided by INTERVIEW DURING DEMO (p. 121).

**Resulting Context.** By applying this pattern, you will have a detailed understanding of the run-time behaviour of a piece of code. This may be necessary to apply patterns in PREPARE REENGINEERING (p. 149).

## **Chapter 7**

# **Cluster: Prepare Reengineering**

The reverse engineering patterns in this cluster are only applicable when your reverse engineering activities are part of a larger reengineering project. That is, your goal is not only understanding what's inside the source code of a software system, but also rewriting parts of it. Therefore, the patterns in this cluster will take advantage of the fact that you will change the source code anyway.

	Limited Resources	Tools and Techniques	Reliable Info	Abstraction	Sceptic Colleagues
WRITE THE TESTS		-	++	++	0
REFACTOR TO UNDERSTAND		-	0	+	0
BUILD A PROTOTYPE		-	+		++
FOCUS BY WRAPPING		0	0	0	0

Table 7.1: How each pattern of PREPARE REENGINEERING resolves the forces. Very good: ++, Good: +, Neutral: 0, Rather Bad: -, Very bad: --

# WRITE THE TESTS

Author(s): Serge Demeyer, Stéphane Ducasse and Sander Tichelaar

## Intent

Record your knowledge about how a component reacts to a given input in a number of black box tests, this way preparing future changes to the system.

**Example.** You are asked to extend a parser for a command language so that it is able to parse two additional commands. Before actually changing the of parser, you will write a number of test programs that check whether the parser accepts all valid command sequences and rejects some typical erronous ones.

## Context

You are at the final stages of reverse engineering a software system, just before you will start to reengineer a part of that system. You have sufficient knowledge about that part to predict its output for a given input.

## Problem

Before starting to reengineer the system, you want to make sure that all what used to work keeps on working.

## Solution

Write a number of black box tests that record your knowledge about the input/output behaviour.

# **REFACTOR TO UNDERSTAND**

Author(s): Serge Demeyer, Stéphane Ducasse and Sander Tichelaar

## Intent

Obtain better readable — thus more understandable — and better organised — thus more extensible — code via renaming and refactoring.

**Example.** You are asked to extend a parser for a command language so that it is able to parse two additional commands. Before actually extending the parser, you will improve the readability of the source code. Among others, you will rename key methods and classes to reflect your understanding of a parser and you will split long and complex methods into smaller ones. As an example of the former, you will rename the class StreamIntf into Scanner and the method rdnxt into nextToken. An example of the latter would be to split the nextToken method, so that it becomes a large case statement, where each branch immediately invokes another method.)

## Context

You are at the final stages of reverse engineering a software system, just before you will start to add new functionality to that system. You have a good programming environment that allows you to rename things easily and that operates on top of a version control system.

## Problem

The shape of the code is such that it is difficult to read —hence to understand— and difficult to add the new functionality.

## Solution

Reorganise the code so that its structure reflects better what the system is supposed to do.

# BUILD A PROTOTYPE

Author(s): Serge Demeyer, Stéphane Ducasse and Sander Tichelaar

## Intent

Extract the design of a critical but cryptic component via the construction of a prototype which later may provide the basis for a replacement.

**Example.** You have a piece of code that implements a graph layout algorithm. You have an idea on how the algorithm works, but the code is too cryptic to map your knowledge of the algorithm onto the code. You will write a prototype that implements your understanding of the algorithm and map pieces of your code onto the existing code.

Context

Problem

**Solution** 

# FOCUS BY WRAPPING

Author(s): Serge Demeyer, Stéphane Ducasse and Sander Tichelaar

## Intent

Wrap the parts you consider unnecessary for the future reengineering in a black box component.

**Example.** You have to migrate a graph manipulation program from a Unix to Macintosh userinterface platform. The original program is well designed and has separated out most of the platform specific operations into a separate layer. You will clean up this layering by moving **all** platform specific behaviour into a separate layer, this way wrapping the obsolete part into a separate component.

Context

Problem

**Solution** 

Chapter 8

**Cluster: misc** 

## CONFER WITH COLLEAGUES

## Intent

Share the information obtained during each reverse engineering activity to boost the collective understanding about the software system.

**Example.** Your team has to reverse engineer a workflow system containing lots of complex rules on how tasks get transferred. Each team member investigates a part of the system and as such the knowledge about the workflow rules is distributed over the team. To increase the overall understanding, you will devote 15 minutes of the weekly team meeting to discuss reverse engineering results made during the last week.

## Context

You are a member of a software development team and part of the job assigned to your team is the reverse engineering of a software system. Different members of the team perform different reverse engineering activities and consequently the knowledge about the system is scattered throughout the team.

### Problem

How do you ensure that every team member contributes to the overall understanding of the software system.

## Solution

Use whatever means at your disposal (meetings, e-mail, intra-nets, ...) to ensure that whenever any team member finishes a reverse engineering step, the obtained information is shared with the rest of the team.

### Hints.

• To avoid information overload, choose the communication channels in such a way that sharing the information fits well with the culture within your team. For instance, do not organise a special team meeting devoted to reverse engineering results; rather use an existing meeting as a vehicle for applying this pattern.

## Rationale

Reverse engineering is sometimes compared with solving a puzzle [Will96b]. If team members keep some pieces of the puzzle for themselves it will never be possible to finish the puzzle. Consequently, it is imperative that a reverse engineering team is organised in such a way that information may be shared among the various team members.

## **Chapter 9**

# **Pattern Overview**

The followin tables summarize the patterns for reference purposes.

The first series of tables lists the patterns together with their problem and their solution, this way aiding reverse engineers to identify which patterns may be applicable to their problem.

The second series of tables show how all the patterns work together to tackle an overall reverse engineering project. For each pattern, the tables list the context and prerequisites plus the pattern results and how these results may serve as input for other patterns.

	FIRST CONTACT	
Pattern	Problem	Solution
READ ALL THE CODE IN ONE HOUR (p. 115)	You need an initial assessment of the internal state a software system to plan further reverse engineer- ing efforts.	Grant yourself a reasonably short amount of study time to walk through the source code. Afterwards produce a report including a list of (i) the important entities; (ii) the coding idioms applied ; (iii) the suspicious coding styles discovered
SKIM THE DOCUMENTATION (p. 118)	You need an initial idea of the functionality pro- vided by the software system in order to plan fur- ther reverse engineering efforts.	Grant yourself a reasonably short amount of study time to scan through the documentation. Afterwards produce a report includ- ing a list of (i) the important requirements; (ii) the important features (iii); the important constraints; (iv) references to rele- vant design information.
INTERVIEW DURING DEMO (p. 121)	You need an idea of the typical usage scenario's plus the main features of a software system in or- der to plan further reverse engineering efforts.	Observe the system in operation by seeing a demo and inter- viewing the person who is demonstrating. Afterwards produce a report including a list of (i) some typical usage scenarios or use cases; (ii) the main features offered by the system and whether they are appreciated or not; (iii) the system components and their responsibilities; (iv) bizarre anecdotes that reveal the folklore around using the system.
	EXTRACT ARCHITECTURE	E
Pattern	Problem	Solution
GUESS OBJECTS (p. 127)	You must gain an overall understanding of the in- ternal structure of a software system and report this knowledge to your collegues so that they will use it as a kind of roadmap for later activities.	Based on your experience, and the little you already understand from the system, devise a model that serves as your initial hy- potheses of what to expect in the source code. Check these hy- potheses against the source code, refine the initial model and recheck the hypotheses. Afterwards, produce a boxes- and ar- rows diagram describing your findings.
CHECK THE DATABASE (p. 130)	You want to derive a data model for the persistent data in a software system in order to guide further reverse engineering efforts.	Check the database schema to reconstruct at least the persistent part of the data model. Use your knowledge of how constructs in the implementation language are mapped onto database con- structs to reverse engineer the real data model. Use the samples of data inside the database to refine the data-model.

	FOCUS ON HOT AREAS	
Pattern	Problem	Solution
INSPECT THE LARGEST (p. 135)	You must identify those places in the source code that correspond with important chunks of functionality.	Use a metrics tool to collect a limited set of measurements for all the constructs in the system. Sort the resulting list according to these measurements. Browse the source code for the largest among those constructs in order to understand how these con- structs work together with other related constructs. Produce a list of all the constructs that appear important, including a de- scription of how they should be used (i.e. external interface).
Exploit The Changes (p. 139)	You must identify those parts in the design that played a key role during the system's evolution.	Use whatever means at your disposal to compile a list of tar- gets of important/frequent changes. For each target, put yourself in the role of the original developer and ask yourself what the change is about and why it was necessary. With this insight, pro- duce a list of crucial system parts, including a description of the design issues that makes them important.
VISUALIZE THE STRUCTURE (p. 142)	You want to obtain insight in the structure of a se- lected part of a software system, including knowl- edge about potential design anomalies.	Instruct the program visualisation tool to show you a series of graphical layouts of the program structure. Based on these graphical layouts, formulate yourself some assumptions and use the code browser to check whether your assumptions are correct. Afterwards, produce a list of correct assumptions, classifying the items in one of two categories: (i) helps program understanding, or (ii) potential design anomaly.
CHECK METHOD INVOCATIONS (p. 144)	You want to find out how a class is related to the other classes in the system.	Select key methods in the interface of the class and inspect who is invoking these methods. Two examples of key methods that are easy to recognise are constructors and overridden methods.
STEP THROUGH THE EXECUTION (p. 146)	You want to obtain a detailed understanding of the run-time behaviour of a piece of code.	Feed representative input data in the piece of code to launch a normal operation sequence. Use the debugger to follow the step by step execution and to inspect the internal state of the piece of code.

Context: You are starting a rev	Context: You are starting a reverse engineering project of a large and unfamiliar software system.	unfamiliar software system.	
Pattern	Prerequisites	Result	What next?
READ ALL THE CODE IN ONE HOUR (p. 115)	<ul> <li>source code</li> <li>expertise with the implementation language</li> </ul>	<ul> <li>the important entities (i.e., classes, packages,)</li> <li>ages,)</li> <li>the coding idioms applied</li> <li>the suspicious coding styles discovered</li> </ul>	<ul> <li>SKIM THE DOCUMENTATION (p. 118) and INTERVIEW DURING DEMO (p. 121) to get alternative views</li> <li>CONFER WITH COLLEAGUES (p. 156) to report findings</li> <li>GUESS OBJECTS (p. 127) and CHECK THE DATABASE (p. 130) to refine the list of important entities</li> </ul>
SKIM THE DOCUMENTA- TION (p. 118)	<ul> <li>documentation</li> <li>you are able to interpret the diagrams and formal specifications contained within</li> </ul>	<ul> <li>important requirements</li> <li>important features</li> <li>important constraints</li> <li>references to relevant design information.</li> <li>an assessment of the reliability and usefulness for each of the above.</li> </ul>	<ul> <li>READ ALL THE CODE IN ONE HOUR (p. 115) and INTERVIEW DURING DEMO (p. 121) to get alternative views</li> <li>CONFER WITH COLLEAGUES (p. 156) to report findings</li> <li>GUESS OBJECTS (p. 127) and CHECK THE DATABASE (p. 130) to map the in- formation on an overall system blueprint.</li> </ul>
INTERVIEW DURING DEMO (p. 121)	<ul> <li>running system</li> <li>somebody who can demonstrate how to use the system</li> </ul>	<ul> <li>typical usage scenarios or use cases</li> <li>the main features offered by the system and whether they are appreciated or not the system components and their respon- sibilities</li> <li>bizarre anecdotes that reveal the folklore around using the system</li> </ul>	<ul> <li>READ ALL THE CODE IN ONE HOUR (p. 115) and SKIM THE DOCUMENTA- TION (p. 118) to get alternative views</li> <li>CONFER WITH COLLEAGUES (p. 156) to report findings</li> <li>GUESS OBJECTS (p. 127) and CHECK THE DATABASE (p. 130) to map the in- formation on an overall system blueprint.</li> </ul>

EXTRACT ARCHITECTURE	Context: You are in the early stages of reverse engineering a software system. You have an initial understanding of its functionality and you are somewhat familiar with the main structure of its source code. (This initial understanding might have been obtained by the patterns in FIRST CONTACT (p. 113)).	Result What next ?	mportant aspects of • a series of blueprints, each one containing • CHECK THE DATABASE (p. 130) if you emperation a perspective on the whole system are interested in the data model	to the source code • all patterns in FOCUS ON HOT AREAS (p. 133) if you want to refine the blueprints	pertise with the im- anguage being used	employs some • a data model of the persistent part of your • system	<ul> <li>ss to the database,</li> <li>proper tools to in-</li> <li>all patterns in FOCUS ON HOT AREAS (p. 133) if you want to refine the datamodel</li> </ul>	how data-structures nplementation lan- napped onto the of the underlying
	stages of reverse engir e of its source code. (]	Prerequisites	<ul> <li>knowledge of important aspects of a software system</li> </ul>	<ul> <li>on-line access to th plus the necessary to ulate it</li> </ul>	• reasonable expertise with the implementation language being used	<ul> <li>software system e form of a database</li> </ul>	• you have access to the database, including the proper tools to in- spect its schema and samples of the data	<ul> <li>knowledge of how data-structures from your implementation lan- guage are mapped onto the data-structures of the underlying database</li> </ul>
	Context: You are in the early stages of reverse eng familiar with the main structure of its source code.	Pattern	GUESS OBJECTS (p. 127) (variants: guess patterns, guess object responsabil-			CHECK THE DATABASE (p. 130)		

		FOCUS ON HOT AREAS	
Context: You are in a later stage of revers with the main structure of its source code.	age of reverse engineering a softwa source code.	ure system. You have an overall und	Context: You are in a later stage of reverse engineering a software system. You have an overall understanding of its functionality and you are fairly familiar with the main structure of its source code.
Pattern	Prerequisites	Result	What next ?
INSPECT THE LARGEST (p. 135)	<ul> <li>a code browser</li> <li>The metrics tool is configured with a number of size metrics</li> </ul>	• a list of constructs represent- ing important functionality	<ul> <li>VISUALIZE THE STRUCTURE (p. 142) to obtain other perspectives on those constructs.</li> <li>STEP THROUGH THE EXECUTION (p. 146) to get a better perception of the run-time behaviour.</li> <li>(in the case of object-oriented source code) CHECK METHOD INVOCATIONS (p. 144) to find out how classes are related to eachother</li> <li>refactoring if you want to split some of these larger contructs into smaller ones</li> </ul>
EXPLOIT THE CHANGES (p. 139) (variants:configuration database, change metrics)	<ul> <li>several releases of the source code</li> <li>a configuration management system and/or a metrics tool</li> </ul>	• a list of design parts that played a key role during the system's evolution	<ul> <li>VISUALIZE THE STRUCTURE (p. 142) to obtain other perspectives on those constructs.</li> <li>STEP THROUGH THE EXECUTION (p. 146) to get a better perception of the run-time behaviour.</li> <li>(in the case of object-oriented source code) CHECK METHOD INVOCATIONS (p. 144) to find out how classes are related to eachother</li> </ul>
VISUALIZE THE STRUC- TURE (p. 142)	<ul> <li>a part of the software system</li> <li>a program visualisation tool</li> <li>a code browser</li> </ul>	<ul> <li>insight in the selected part</li> <li>list of potential design anomalies</li> </ul>	<ul> <li>STEP THROUGH THE EXECUTION (p. 146) to get a better perception of the run-time behaviour.</li> <li>(in the case of object-oriented source code) CHECK METHOD INVOCATIONS (p. 144) to find out how classes are related to eachother</li> <li>refactoring if you want to split some of these larger contructs into smaller ones</li> </ul>

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re system • a list of classes and the rela- • STEP THROUGH THE EXECUTION (p. 146) to get a tion tool tionships between them better perception of the run-time behaviour. allows you there the the the the the hod is de-	system • insight into the run-time be- • PREPARE REENGINEERING (p. 149) if you need to haviour of a piece of code reengineer that piece of code ive input
<ul> <li>a part of the software system</li> <li>a program visualisation tool</li> <li>a code browser that allows you to jump from a method invocation to the places where the corresponding method is defined</li> </ul>	<ul><li>a part of the software system</li><li>an interactive debugger</li><li>a set of representative input data</li></ul>
<ul> <li>CHECK METHOD INVOCA-</li> <li>a part of the software system TIONS (p. 144)</li> <li>a program visualisation tool ods, overridden methods)</li> <li>a code browser that allows y to jump from a method investigation to the places where t corresponding method is cfined</li> </ul>	STEP THROUGH THE EXE- cution (p. 146)

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# Part III

Reengineering

## Chapter 10

# **Reengineering Patterns**

A *reengineering pattern* describes how to go from an existing *legacy* solution to a *refactored* solution that better suits the current requirements. In this chapter we explain why we choose the pattern form to communicate reengineering expertise and present the reengineering pattern form. We stress the differences between the Design Patterns and the Reengineering Patterns, and also between the Reengineering patterns themselves. It should be noted that the reengineering patterns are not linked together in a pattern language due to a lack of time.

## **10.1 Reengineering Patterns: a Need**

Reengineering projects, despite their diversity, often encounter some typical problems again and again. These can be problems at different levels and due to different practices [FY97]. But it is unlikely that one methodology or process will be appropriate for all projects and organisations [SP98], just as not one tool or technique can be expected to solve all the technical problems encountered in a reengineering project. To allow reengineering projects to benefit from the experience gained in previous efforts, an appropriate form is required for transferring expertise. This form should be small enough to be easily consulted and navigated, and stable enough as to be useful for many reengineering projects.

In the object-oriented software engineering community *Design Patterns* [GHJV95] have been adopted as an effective way of communicating expertise about software design. A design pattern describes a solution for a recurring design problem in a form which facilitates the reuse of a proven design solution. In addition to the technical description of the solution, an important element of a design pattern is its discussion of the advantages and disadvantages of applying the pattern.

We propose the use of the pattern form as a means of communicating expertise in the area of reengineering. *Reengineering patterns* codify and record knowledge about modifying legacy software: they help to diagnose problems and identify weaknesses hindering further development of the system as well as aiding the search for solutions that better satisfy the new requirements. We see reengineering patterns as stable units of expertise which can be consulted in any reengineering effort: they describe a process without proposing a complete methodology, and they suggest appropriate tools without 'selling' a specific one. A more thorough discussion of the advantages of the pattern form as a vehicle for reengineering expertise can be found in [SP98], which discusses patterns closely related to ours.

All the reengineering patterns presented hereafter address problematic legacy solutions typically found in object-oriented code, and describe how to move from the *legacy* solution to a new *refactored* solution. The patterns presented are all of a technical nature, dealing directly with source code problems. There exist however higher-level reengineering patterns which describe overall strategies for dealing with legacy systems. See for instance the *Systems Reengineering Patterns* [SP98] which address broader methodological

issues. The 'Deprecation' pattern [SP98], for example, describes how to iteratively change interfaces of a system in a friendly way for the client of the system under change.

## **10.2 Reengineering Patterns and Related Work**

The reengineering patterns presented here and the *Systems Reengineering Patterns* of [SP98] are closely. The principle differences is that here the patterns are source-level rather high-level and they are focused on bject-oriented legacy systems<sup>1</sup>. Note that our patterns cannot be used to evaluate whether or not an application should be reengineered in the first place; this difficult task has been tackled by [STS97] and [RSW98]. In [BS95] a methodology is proposed to help in the migration of legacy systems (principally legacy database systems) to new platforms.

Reengineering patterns differ from *Design Patterns* [GHJV95] in their emphasis on the *process* of moving from an existing *legacy* solution that is no longer appropriate to a new *refactored* solution. The mark of a good reengineering pattern is (a) the clarity with which it exposes the advantages, the cost and the consequences of the target solution with respect to the existing solution, and not how elegant the target solution is, and (b) the description of the change process: how to get from the legacy version of the system to the refactored version.

We also contrast reengineering patterns with *AntiPatterns* [BMMM98]. Antipatterns, as exposed by Brown et al., are presented as "bad" solutions to design and management issues in software projects. Many of the problems discussed are managerial concerns that are outside the direct control of developers. Moreover, the emphasis in antipatterns is on prevention: how to avoid making the mistakes which lead to the "bad" solution. Consequently, antipatterns may be of interest when starting a project or during development but are no longer helpful when we are confronted with a legacy system. In reengineering, though, we prefer to withhold the judgement inherent in the notion of "bad solution" and use the term "legacy solution" or "legacy pattern" for a solution which at the time, and under the constraints given, seemed appropriate. In reengineering it is too late for prevention, and reengineering patterns therefore concentrate on the cure: how to detect problems and move to more appropriate solutions.

Finally, our reengineering patterns are different from code *refactorings* [JO93, JO93]. A reengineering pattern describes a process which starts with the detection of the symptoms and ends with the refactoring of the code to arrive at the new solution. A refactoring is only the last stage of this process, and addresses only the technical issue of automatically or semi-automatically modifying the code to implement the new solution. Reengineering patterns also include other elements which are not part of refactorings: they emphasise the context of the symptoms by taking into account the constraints being faced and include a discussion of the impact of the changes introduced by the refactored solution.

A reengineering pattern may describe a solution that would not be ideal if one is designing a system from scratch, but is a good solution under the current constraints of the legacy system. For example, if the constraint is that changes must be kept local some solutions are clearly not applicable even if they seem at first hand to be the best solutions.

## **10.3** Form of a reengineering pattern

The primary goal of a reengineering pattern is to help developers in solving reengineering problems. The idea being that a developer must diagnose a problem, identify the available options and choose a particular course of action. Furthermore, the relevant weaknesses must be identified, where relevant is defined in terms of the desired flexibility or some other quality, and the system transformed so that it possesses the desired quality.

<sup>&</sup>lt;sup>1</sup>We do not address problems of reengineering procedural applications to object-oriented ones.

The pattern form has been defined principally for *reengineering* patterns: that is, patterns which describe a transformation of a existing design to a more appropriate, target design. In general the pattern form has been defined with the following requirements in mind, although some patterns may - due to specific needs - add or omit sections described in the presented pattern format:

**Focus on Reengineering Process.** A reengineering pattern is different from a design pattern. It should go beyond discussing good and bad designs. A reengineering pattern should also discuss the reengineering *process*! For example if we know a design to contain problems according to present requirements, then how can these problems be discovered; or, what are the pitfalls in transforming a system? Both are issues that have to do with the reengineering process itself. A typical reengineering pattern will describe a process that transforms a system with a design that is no longer adequate into a system with an improved design. The reengineering pattern should clearly identify these two states of the system and their relationship.

**Easy Navigation.** The idea is that a handbook user (i.e. reengineer) should be able to determine if a pattern is applicable within the *first* page of description.

**Separate out Tool and Language Dependent Issues.** To make the patterns as generally applicable as possible, tool and language dependent issues should be separated out as much as possible. The main part of the patterns describes stable reengineering knowledge, whereas tools are more subject to evolution and in some cases language dependent issues can be interesting but not of influence on the core idea of the pattern.

**Standard Terminology and Notation.** A language neutral terminology and notation is mandatory if the patterns are to be kept as language independent as is reasonably possible. The rule for terminology is: as far as it is defined the UML terminology [SMHP+97] is used for object oriented concepts and if a term is not defined by UML then the terminology of the FAMOOS Information Exchange model (see section ??) is used. All other terms are to be defined in a glossary that is part of this handbook. For the homogeneity of the patterns, a strong requirement is that all the drawings should be done using UML notations.

The specific aspects of reengineering patterns lead us to the definition of an adapted form for the reengineering patterns. This form is structured as follows:

- **Pattern Name.** The name is based on the reengineering operation that is performed as this is the most natural way of discussing the pattern in the context of reengineering. It must form the basis for a terminology for reengineers to talk about reengineering a system. As a temporary solution, patterns which miss a good word will be named by a short sentence with a verb that emphasises the kind of reengineering transformation.
- Intent. A description of the process, together with the result and why it is desirable.
- **Applicability.** When is the pattern applicable? When is it not applicable? This section includes a list of symptoms, a list of reengineering goals and a list of related patterns. Symptoms are those experienced when reusing, maintaining or changing the system. For example, correlations between editing different parts of a system for making a certain change can indicate the need for a particular kind of reengineering. Reengineering goals present the qualities improved through the application of this pattern.
- **Motivation.** This section presents an example: it must acquaint the reader with a concrete example so they can better understand the more abstract presentation of the problem which follows in the structure and process sections. The example must clearly describe the structure of the legacy system, the structure of the reengineered system, and the relation between the two. The state of the system before and after the application of the pattern are described.

- **Structure.** This section describes the structure of the system before and after reengineering. Each structure section is similar to the structure section in the Gang of Four pattern book. The participants and their collaborations are identified. Consequences discuss the advantages and disadvantages of the target structure in comparison to the initial structure.
- **Process.** The process section is subdivided into three sections: the detection, the recipe and the difficulties. The detection section describes methods and tools to detect that the code is indeed suffering from the suspected problem and that the process given below can help to alleviate this problem. The recipe states how to perform the reengineering operation and the possible variants. The optional difficulties subsection discusses situations where the reengineering operation is not feasible or is compromised by other problems.
- **Discussion.** In this section, forces of the legacy solution are discussed first. Indeed, often legacy solutions fulfilled the requirements at the time they were implemented. When requirements change however, the solution must evolve to accommodate the new requirements. The forces of the refactored solution are stressed in terms of cost and benefit tradeoffs of applying the pattern. What is the cost of detecting this problem? What is the magnitude of the problem? What is the benefit gained by applying the pattern? This discussion should aid an engineer in deciding (once he knows the pattern is applicable to the code) whether or not it is, in this specific case, worth applying the pattern. Moreover, relationships with Design Patterns or AntiPatterns should be documented.

The sections above form the core of the pattern. The sections described below deal with more concrete issues and are essential to the reengineering handbook, where engineers need information about language specific issues and existing tool support and know applications.

- Language Specific Issues. This section lists what must be specifically resolved for each language. What makes it more difficult? More easy?
- **Tool Support.** Lists and describes tools to support the detection of symptoms, detection of participants and collaborations, and to aid in the transformation of the system.
- **Known Uses.** Gives known cases where the pattern has been applied, successfully or not. In our context, references are made to our industrial case studies.

## **10.4** Pattern Navigation

The patterns that we have collected for this handbook solve a quite random set of reengineering problems. To provide the reader with some guidance, we present ways to select the patterns that might be of interest in solving a specific problem of yours. The usefulness of the navigational support is inherently limited due to the rather small set of available patterns, but it provides an overview of what is already covered by the patterns.

We currently provide navigation based based on forces that typically play a role in reengineering (see section 10.4.1). For each type of navigation, a table is shown that gives a broad overview of which pattern covers what. Following the table, the patterns are listed – according to the same structure as the table – with a short explanation of their appearance in the table. This allows for a quick scan of the available patterns.

### **10.4.1** Forces

In this section we categorise the patterns according to how they affect the forces that are at work in a reengineering project. In table 10.1 we show the reengineering patterns and their influence on the different forces. The patterns are shown vertically and the forces horizontally. A '+' means that the pattern increases

the effect of, or has a high impact on, the particular force. A '-' means that the pattern reduces the impact of, or has a low impact on, the force. No sign means that the pattern either has no or an unpredictable influence on the force. To give an example, the Code Duplication Detection pattern requires only minor effort when applying it, scales up well and does not require much parsing.

	Forces						
	Flexibility	Understandability	Reusability	Effort	Scalability	Parsing effort	Global Impact
Type Check Elimination in Clients	+	+	+				-
Type Check Elimination within a Provider Hierarchy	+	+					
Detection of Duplicated Code				-	+	-	
Repairing a Broken Architecture	+	+					
Transforming Inheritance into Composition	+		+				
Distribute Responsibilities	+	+	+				

Table 10.1: How the individual patterns affect the forces of a reengineering effort.

### Flexibility

TYPE CHECK ELIMINATION IN CLIENTS. By reducing the coupling between the clients and the provider class hierarchy by refactoring the interface of the provider classes and the client code that depends on these interfaces thus making the clients much more robust. This greatly facilitates extending the functionality of the provider hierarchy without breaking the client code.

TYPE CHECK ELIMINATION WITHIN A PROVIDER HIERARCHY. Transforms a single provider class being used to implement what are conceptually a set of related types into a hierarchy of classes. Decision structures, such as case statements or if-then-elses, over type information are replaced by polymorphism. This results in increased modularity and facilitates the extension of functionality through the addition of new subclasses.

REPAIRING A BROKEN ARCHITECTURE. Detects and removes dependencies between packages of a system that aren't allowed according to the designated system architecture. These architecture breaking dependencies may prohibit the exploitation of the architecture's advantages and cause unexpected effects at maintenance work.

TRANSFORMING INHERITANCE INTO COMPOSITION. This pattern describes, how to transform an inheritance relationship into a component relationship. This increases flexibility, because a component relationship can be changed dynamically whereas an inheritance relationship can only be changed statically.

DISTRIBUTE RESPONSIBILITIES. Distributes the responsibilities equally among the classes of an object-oriented system to prevent large, hardly maintainable and reusable classes.

### Understandability

TYPE CHECK ELIMINATION IN CLIENTS. The reduced coupling between client and provider classes as well as the refactored interface of the provider classes present a more modular view of what is essentially the same functionality thus facilitating understandability.

TYPE CHECK ELIMINATION WITHIN A PROVIDER HIERARCHY. Breaking a single complex class into a hiearchy of simpler but more specialised classes facilitates partial understanding rather than requiring complete understanding. This simplifies understanding how the hierarchy can be extended by separating information that is relevant to the entire hierarchy from that which is specific to just a few classes.

REPAIRING A BROKEN ARCHITECTURE. Detects and removes dependencies between packages of a system that aren't allowed according to the designated system architecture. These architecture breaking dependencies may prohibit the exploitation of the architecture's advantages and cause unexpected effects at maintenance work.

DISTRIBUTE RESPONSIBILITIES. Distributes the responsibilities equally among the classes of an object-oriented system to prevent large, hardly maintainable and reusable classes.

### Reusability

TYPE CHECK ELIMINATION IN CLIENTS. The refactored interface of the classes from the provider hierarchy more accurately reflect the needs of any client classes. This increases the likelyhood that classes from the provider hierarchy can be reused while simplifying their reuse.

TRANSFORMING INHERITANCE INTO COMPOSITION. This pattern describes, how to transform an inheritance relationship into a component relationship. This increases flexibility, because a component relationship can be changed dynamically whereas an inheritance relationship can only be changed statically.

DISTRIBUTE RESPONSIBILITIES. Distributes the responsibilities equally among the classes of an object-oriented system to prevent large, hardly maintainable and reusable classes.

#### **Relatively minor effort**

DETECTION OF DUPLICATED CODE. Collecting duplication data is fully automatised. Filtering the data to find candidates for refactoring automatically is possible only in some special cases. Normally, human intervention and expertise is required to assess the duplication and decide on the possible refactoring operations.

#### Patterns that easily scale up

DETECTION OF DUPLICATED CODE. Code Duplication is done with tool support. All of the tools are built to scale up well.

#### Patterns that need a low parsing effort

DETECTION OF DUPLICATED CODE. Depending on the technique used to detect duplication, more or less parsing is required. The DUPLOC tool (see Chapter ??) developed in FAMOOS requires only minimal parsing.

### **Global impact**

TYPE CHECK ELIMINATION IN CLIENTS. The pattern involves refactoring the interface of a class hierarchy in order to better support the clients of these classes. Consequently, most clients will be affected which may potentially require that changes be made throughout the system.

### **Higher number of Classes**

DISTRIBUTE RESPONSIBILITIES. Distributes the responsibilities equally among the classes of an object-oriented system to prevent large classes that are difficult to maintain and reuse.

TYPE CHECK ELIMINATION WITHIN A PROVIDER HIERARCHY. Separates a single complex class into a hiearchy of simpler more specialized classes representing a cleaner separation of concerns. The increased number of classes allow a greater precision in expressing dependencies.

Chapter 11

**Cluster: Type Check Elimination** 

# TYPE CHECK ELIMINATION IN CLIENTS

Author(s): Stéphane Ducasse, Robb Nebbe and Tamar Richner

### Intent

Transform *client* classes that depend on type tests (usually in conjunction with case statements) into *clients* that rely on polymorphism. The process involves factoring out the functionality distributed across the clients and placing it in the provider hierarchy. This results in lower coupling between the *clients* and the *providers* (class hierarchy).

## Applicability

### Symptoms.

- Large decision structures in the *client* over the type of (or equivalent information about) an instance of the *provider*, either passed as an argument to the client, an instance variable of the client, or a global variable.
- Adding a new subclass of the *provider* superclass requires modifications to *clients* of the *provider* hierarchy because functionality is distributed over these clients.

#### **Reengineering Goals.**

- Localise functionality distributed across *clients* in the *provider* hierarchy.
- Improve usability of *provider* hierarchy.
- Lower coupling between *clients* and the *provider* hierarchy.

**Related Reengineering Patterns.** A closely related reengineering pattern is TYPE CHECK ELIMINA-TION WITHIN A PROVIDER HIERARCHY, where the case statements over types are in the *provider* code as opposed to the *client* code. The essential distinction is if the decision structure is over the type or an attribute functioning as a type of: (a) an instance of *another* class (this pattern) or (b) an instance of the class to which the method belongs (see TYPE CHECK ELIMINATION WITHIN A PROVIDER HIERARCHY in section 11).

## Motivation

The fact that the clients depend on provider type tests is a well known symptom for a lack of polymorphism. This leads to unnecessary dependencies between the classes and it makes it harder to understand the program because the interfaces are not uniform. Furthermore, adding a new subclass requires all clients be adapted. **Initial Situation.** The following code illustrates poor use of object-oriented concepts as shown by Fig. 11.1. The function makeCalls takes a vector of Telephone's (which can be of different types) as a parameter and makes a call for each of the telephones. The case statement switches on an explicit type-flag returned by phoneType(). In each branch of the case, the programmer calls the phoneType specific methods identified by the type-tag to make a call.

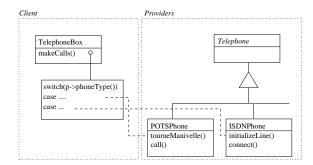


Figure 11.1: Initial relation and structure of clients and providers.

```
void makeCalls(Telephone * phoneArray[])
{
 for (Telephone *p = phoneArray; p; p++) {
   switch(p->phoneType()) {
   case TELEPHONE::POTS: {
    POTSPhone * potsp = (POTSPhone *) p;
    potsp->tourneManivelle();
    potsp->call(); break;}
   case TELEPHONE::ISDN: {
    ISDNPhone * isdnp = (ISDNPhone *) p;
    isdnp->initializeLine();
    isdnp->connect(); break;}
   case TELEPHONE::OPERATORS: {
    OperatorPhone * opp = (OperatorPhone *) p;
    opp->operatormode(on);
    opp->call(); break;}
   case TELEPHONE::OTHERS:
   default:
      error(....);
```

**Final Situation.** After applying the pattern the corresponding ringPhones() will look as follows and the structure as shown by the Fig. 11.2.

```
void makeCalls(Telephones *phoneArray[])
{
  for(Telephone *p = phoneArray; p; p++) p->makeCall();
}
```

Note that the client code, which represents distributed functionality, has been greatly simplified. Furthermore, this functionality has been localised within the **Telephone** class hierarchy, thus making it more complete and uniform with respect to the clients needs.

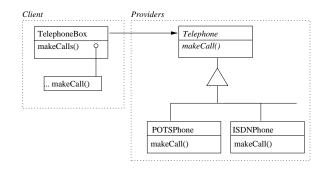


Figure 11.2: Final relation and structure of clients and providers.

### Structure

### Participants.

- provider classes (Telephone and its subclasses)
  - organised into a hierarchy.
- the clients (TelephoneBox) of the provider class hierarchy.

**Collaborations.** The collaborations will change between all clients and the providers as well as the collaboration within the provider hierarchy.

Initially, the clients collaborate directly with the provider superclass and its subclasses by virtue of type tests or a case statement over the types of the subclasses. After reengineering the only direct collaboration between the clients and the providers is through the superclass. Interaction specific to a subclass is handled indirectly through polymorphism.

Within the provider hierarchy the superclass interface must be extended to accurately reflect the needs of the clients. This will involve the addition of new methods and the possible refactorisation of the existing methods in the superclass. Furthermore, the collaborations between the provider superclass and its subclasses may also evolve, i.e. it must be determined whether the new/refactored methods are abstract or concrete.

**Consequences.** Relying on polymorphism localises the protocol for interacting with the provider classes within the superclass. The collaborations are easier to understand since the interface actually required by the clients is now documented explicitly in the provider superclass. It also simplifies the addition of subclasses since their responsibilities are defined in a single place and not distributed across the clients of the hierarchy.

## Process

**Detection.** The technique described in the pattern TYPE CHECK ELIMINATION WITHIN A PROVIDER HIERARCHY to detect case statements is applicable for this pattern. Whereas in the pattern TYPE CHECK ELIMINATION WITHIN A PROVIDER HIERARCHY, the switches are located in the same class, hence in one file for a language like  $C^{++}$ , in this pattern the case statements occur in several classes which can be spread over different files.

**Recipe.** The process consists of two major steps. The first is to encapsulate all the responsibilities that are specific to the provider classes within the provider hierarchy. The second is to make sure that these responsibilities are correctly distributed within the hierarchy.

- 1. Determine the set of clients to which the pattern will be applied.
- 2. Define a new abstract method in the provider superclass and concrete methods implementing this method in each of the subclasses based on the source code contained within each branch of the case statement.
- 3. Refactor the interface of the provider superclass to accurately reflect the protocol used by the clients. This involves not only adding and possibly changing the methods included but determining how they work together with the subclasses to provide the required behaviour. This includes determining whether methods are abstract or concrete in the provider superclass.
- 4. For each client, rewrite the method containing the case statement so that it uses only the interface of the provider superclass.

### Difficulties.

- 1. The set of clients may all employ the same protocol; in this case the pattern needs to be applied only once. However, if the clients use substantially different protocols then they can be divided into different kinds and the pattern must be applied once for each kind of client.
- 2. If the case statement does not cover all the subclasses of the provider superclass a new abstract class may need to be added and the client rewritten to depend on this new class. For example, if it is an error to invoke the client method with some subclasses as opposed to just doing nothing then the type system should be used to exclude such cases. This reduces the provider hierarchy to the one starting at the new abstract class.
- 3. Refactoring the interface will affect all clients of the provider classes and must not be undertaken without examining the full consequences of such an action.
- 4. Nested case statements indicate that multiple patterns must be applied. This pattern may need to be applied recursively in which case it is easiest to apply the pattern to the outermost case statement first. The provider classes then become the client classes for the next application of the pattern. Another possibility is when the inner case statement is also within the provider class but some of the state of the provider classes should be factored out into a separate hierarchy.

## Discussion

During the detection phase one can find other uses of case statements. For example, case statements are also used to implement objects with states [Bec94], [ABW98]. In such a case the dispatch is not done on object type but on a certain state as illustrated in the State pattern [GHJV95], [ABW98]. Moreover, the Strategy pattern [GHJV95], [ABW98] is also based on the elimination of case statement over object state.

 $C^{++}$  In  $C^{++}$  virtual methods can only be used for classes that are related by an inheritance relationship. The polymorphic method has to be declared in the superclass with the keyword virtual to indicate that calls to this methods are dispatched at runtime. These methods must be redefined in the subclasses.

Since  $C^+$  does not offer runtime type information, type information is encoded mostly using some enum type. A data member of a class having such an enum type and a method to retrieve these tags are usually a hint that polymorphism could be used (although there are cases in which polymorphic mechanism cannot substitute the manual type discrimination).

**ADA** Detecting type tests falls into two cases. If the hierarchy is implemented as a single discriminated record then you will find case statements over the discriminant. If the hierarchy is implemented with tagged types then you cannot write a case statement over the types (they are not discrete); instead an if-then-else structure will be used.

If a discriminated record has been used to implement the hierarchy it must first be transformed by applying the TYPE CHECK ELIMINATION WITHIN A PROVIDER HIERARCHY pattern.

**SMALLTALK** In SMALLTALK the detection of the case statements over types is hard because few type manipulations are provided. Basically, methods isMemberOf: and isKindOf: are available. anObject isMemberOf: aClass returns true if anObject is an instance of the class aClass, anObject isKindOf: aClass returns true if aClass or a superclass of anObject. Detecting these method calls is not sufficient, however, since class membership can also be tested with self class = anotherClass, or with property tests throughout the hierarchy using methods like isSymbol, isString, isSequenceable, isInteger.

## Tools

Glimpse and agrep can be found at ftp://ftp.cs.arizona.edu/glimpse.

## **Known Uses**

In the FAMOOS mail sorting case study, we identified 28 matches (a match is not equivalent to a file because a same file may contain several switches) for the expression agrep 'switch;type', 185 matches for the sole expression agrep 'switch'. In the same time agrep 'if' gave us 10976 matches whereas using the perl script shown above we reduce the matches to 497.

In this case study, we identify three obvious lacks of polymorphism but they were not corresponding with the presented pattern but its companion pattern TYPE CHECK ELIMINATION WITHIN A PROVIDER HIER-ARCHY. We found also cases that implement state object [GHJV95].

This pattern has been also applied in one of the FAMOOS case studies written in Ada. This considerably decreased the size of the application and improved the flexibility of the software. In one of the FAMOOS C++ case studies, manual type check also occurs implemented statically via # ifdefs.

# TYPE CHECK ELIMINATION WITHIN A PRO-VIDER HIERARCHY

### Author(s): Stéphane Ducasse, Robb Nebbe and Tamar Richner

# Intent

Transform a single *provider* class being used to implement what are conceptually a set of related types into a hierarchy of classes. Decision structures, such as case statements or if-then-elses, over type information are replaced by polymorphism. This results in increased modularity and facilitates the extension of functionality through the addition of new subclasses.

# Applicability

### Symptoms.

- Methods contain large decision structures over an instance variable of the *provider* class to which they belong.
- Extending the functionality of the *provider* class requires modifying many methods.
- Many *clients* depend on a single *provider* class.

### **Reengineering Goals.**

- Improve modularity.
- Simplify extension of *provider* functionality.

**Related Reengineering Patterns.** A closely related pattern is TYPE CHECK ELIMINATION IN CLIENTS where the case statements over types are in the client code as opposed to the provider code. The essential distinction is if the decision structure is over an instance variable of the class (this pattern) or another class (see TYPE CHECK ELIMINATION IN CLIENTS in section 11).

# Motivation

Case statements are sometimes used to simulate polymorphic dispatch. This often seems to be the result of the absence of polymorphism in an earlier version of the language (Ada'83  $\rightarrow$  Ada'95 or C  $\rightarrow$  C<sup>++</sup>). Another possibility is that programmer don't fully master the use of polymorphism and as a result do not always recognise when it is applicable. In any language that supports polymorphism it is preferable to exploit the language support rather than simulate it.

In the presence of polymorphism the process of dispatching is part of the language. In contrast, with case statements or other large decision structures the simulated dispatch must be hand coded and hand maintained. Accordingly, changing or extending the functionality are more difficult because they often

affect many places in the source code. It also results in long methods with obscured logic that are hard to understand.

Programmers often fall back to the language they are most familiar with – in the Variable State pattern Kent Beck shows an example of such a situation related to Lisp programmers [Bec97]. Thus, they may continue to implement solutions which do not exploit polymorphism even when polymorphism is available. This could occur especially when programmers extend an existing design by programming around its flaws, rather than reengineering it.

**Initial Situation.** Our example, taken in a simplified form from one of the case studies, consists of a message class that wraps two different kinds of messages (TEXT and ACTION) that must be serialised to be sent across a network connection as shown in the code and the figure 11.3.

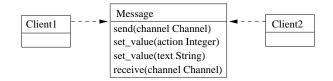


Figure 11.3: Initial relation and structure of clients and providers.

A single provider class implements what is conceptually a set of related types. One attribute of the class functions as surrogate *type* information and is used in a decision structure to handle different variations of functionality required.

```
class Message {
public:
 Message();
 set_value(char* text);
 set_value(int action);
 void send(Channel c);
 void receive(Channel c);
...
private:
 void* data;
 int type_;
ł
// from Message::send
 const int TEXT = 1;
 const int ACTION = 2;
 switch (type_) {
 case TEXT: ...
 case ACTION: ... };
```

### **Final Situation.**

The case statements have been replaced by polymorphism and the original class has been transformed into a hierarchy comprised of an abstract superclass and concrete subclasses. Clients must then be adapted to create the appropriate concrete subclass.

Initially there may be a large number of dependencies on this class, making modification expensive in terms of compilation time, and increasing the effort required to test the class. The target structure improves all of these problems with the only cost being the effort required to refactor the provider class and to adapt the clients to the new hierarchy.

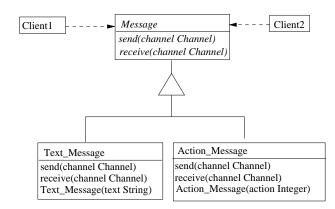


Figure 11.4: Final relation and structure of clients and providers.

```
class Message {
public:
 virtual void send(Channel c) = 0;
 virtual void receive(Channel c) = 0;
...
};
class Text_Message: public Message {
public:
 Text_Message(char* text);
 void send(Channel c);
 void receive(Channel c);
private:
 char* text;
...
};
class Action_Message: public Message {
public:
 Action_Message(int action);
 void send(Channel c);
 void receive(Channel c);
private:
 int action;
...
};
```

# Structure

### Participants.

- A single **provider** (Message) class that is transformed into a hierarchy of classes (Message, Text\_Message and Action\_Message)
- A set of client classes

**Collaborations.** The single provider class will be transformed into a hierarchy, thereby increasing modularity and facilitating extension of functionality.

Initially, the clients are all dependent on a single provider class. This class encompasses several variants of functionality and thus encapsulates all the collaboration that would normally be handled by polymorphism. This results in long methods typically containing case statements or other large decision structures.

The situation is improved by refactoring the single provider class into a hierarchy of classes: an abstract superclass and a concrete subclass for each variant. Each of the new subclasses is simpler than the initial class and these are relatively independent of each other.

**Consequences.** The functionality of the hierarchy can be extended adding a new subclass without modifying the superclass. The increased modularity also impacts the clients who are now likely to be dependent on separate subclasses in the provider hierarchy.

### Process

#### **Detection.**

A class having many long methods is a good candidate for further analysis. A line of code per method metric may help to narrow the search. If these methods contain case statements or complex decision structures all based on the same attribute then the attribute is probably serving as surrogate type information. In C<sup>++</sup>, where it is a good practice to define a class per file, the frequency of case statements in the same file can be also used as a first hint to narrow the search for this pattern.

**Example: detection of case statements in C**<sup>++</sup> Knowing if the pattern should be applied requires the detection of case statements. Regular-expression based tools like emacs, grep, agrep help in the localisation of case statements based on explicit construct like C<sup>++</sup>'s switch or Ada case. For example, grep 'switch' 'find . -name "\*.cxx" -print' enumerates all the files with extension .cxx contained in a directory tree that contains switch. The grep facilities for grep are extended in agrep so it is possible to ask for finer queries. For example, the expression agrep 'switch;type' -e 'find . -name "\*.cxx" -print' extracts all the files containing lines having switch and type.

However, even for a language like  $C^{++}$  that provides an explicit case statement construct, detecting case statements based on explicit if thenelse structures is necessary. The tools above are not well suited for such a task, since their detection capabilities are restricted to one line at a time. One possible solution is to use perl scripts - a perl script which searches the methods in  $C^{++}$  files and lists the occurrences of case statements can be found in the appendix.

### Recipe.

- 1. Determine the number of conceptual types currently implemented by the class by inspecting the case statements. An enumeration type or set of constants will probably document this as well.
- 2. Implement the new provider hierarchy. You will need an abstract superclass and at least one derived concrete class for every variant.
- 3. Determine if all of the methods need to be declared in the superclass or if some belong only in a subclass.
- 4. Update the clients of the original class to depend on either the abstract superclass or on one of its concrete subclasses.

### Difficulties.

- If the case statements are not all over the same set of functionality variants this is a sign that it might be necessary to have a more complex hierarchy including several intermediate abstract classes, or that some of the state of the provider should be factored out into a separate hierarchy.
- If a client depends on both the superclass and some of the subclasses then you may need to refactor the client class or apply the TYPE CHECK ELIMINATION IN CLIENTS pattern because this is an indication that the provider does not support the correct interface.

### Discussion

During the detection phase one can find other uses of case statements. For example, case statements are also used to implement objects with states [Bec94], [ABW98]. In such a case the dispatch is not done on object type but on a certain state as illustrated in the State pattern [GHJV95], [ABW98]. Moreover, the Strategy pattern [GHJV95], [ABW98] is also based on the elimination of case statement over object state.

In his thesis Opdyke [JO93] discusses the automatisation of code refactoring. His "Refactoring To Specialise", in which he proposed to use class invariant as a criteria to simplify conditionals, is similar to this pattern.

### Language Specific Issues.

 $C^{++}$  Detection: in C polymorphism can be emulated either by using function pointers or through union types and enum's.  $C^{++}$  programmers are likely to use a single class with a void pointer and then cast this pointer to the appropriate type inside a switch statement. This allows them to uses classes which are nominally object-oriented as opposed to unions which they have probably been told to avoid. The use of constants is typically favoured over the use of enum's.

Difficulties: If void pointers have been used in conjunction with type casts then you should check to see if the classes mentioned in the type casts should be integrated into the new provider hierarchy.

**ADA** Detection: because Ada83 did not support polymorphism (or subprogram access types) discriminated record types are the preferred solution. Typically an enumeration type provides the set of variants and the conversion to polymorphism is straightforward in Ada95.

**SMALLTALK** In SMALLTALK the detection of the case statements over types is hard because few type manipulations are provided. Basically, methods isMemberOf: and isKindOf: are available. anObject isMemberOf: aClass returns true if anObject is an instance of the class aClass, anObject isKindOf: aClass returns true if aClass or a superclass of anObject. Detecting these method calls is not sufficient, however, since class membership can also be tested with self class = anotherClass, or with property tests throughout the hierarchy using methods like isSymbol, isString, isSequenceable, isInteger.

### Tools

Glimpse and agrep can be found at ftp://ftp.cs.arizona.edu/glimpse.

# **Known Uses**

In one FAMOOS case study several instances of this problem were found. In the example studied in depth (DialogElement) it appears in conjunction with a class that groups together user interface and core model functionality. There is a data member called \_type that is used in the various switch statements. Furthermore a void pointer is frequently cast to an appropriate type based on the value of \_type.

Chapter 12

**Cluster: Duplicated Code** 

# DETECTION OF DUPLICATED CODE

#### Author(s): Matthias Rieger and Stéphane Ducasse

## Intent

Detect code duplication in a system, without prior knowledge of the code. Identifying the duplicated code is a first important step towards application refactoring.

# Applicability

The only prerequisite is the availability of the source code.

#### Symptoms.

- You already saw the same source somewhere else in the application.
- You already fixed the same error in another piece of code.
- You make a conceptual change and in adapting the software to the new concept have to edit similar pieces of code over and over again.
- You know you employed copy and paste programming during development, but do not remember exactly were it was.

**Reengineering Goals.** Some of the following reengineering goals are not only linked to the identification of duplicated code but also to its removal by refactoring:

- **Identifying unknown duplicated code.** This pattern is well-suited to identify **unknown** and middle size (4 to 100 lines) of duplicated code. If you are looking for occurrences of a particular line of code, use **sed** or **grep**-like tools or emacs (regexp and etag) facilities. If you are sure that the developers had to use copy and paste coding (e.g. your software contains about 4 millions lines of code and was developed by 2 people during one year) but want to know what has been copied and pasted, apply this pattern.
- **Identifying duplicated code in large scale system.** Following the previous point, if you are looking for a way to identify duplicated code in a big (100'000 lines) to huge system apply this pattern.
- **Improving maintenance.** Detection helps the maintainer of a system to make sure that some code fragment, where an error has been fixed, is not copied a number of times with the error still in it, or, complicating matters further, is fixed differently at each location by maintainers who have no knowledge of each other's activities.
- **Reducing maintenance cost.** By detecting clones of a piece of code to be maintained and merging the code into one instance, the multiplied effort otherwise necessary to maintain all the clone instances is removed.

- **Improving the code readability.** By identifying duplicated code and refactoring it, the size of code is reduced. The level of abstraction is elevated when similar code pieces are refactored in a new method, ultimately leading to the SMALLTALK ideal of 6 lines of code per method. In one of the FAMOOS case studies, we found a method of 6000 lines of C++ code, which is a nightmare in complexity by any standards.
- Improving compilation time. The less lines of code you have, the faster your system is compiled.
- **Reducing the footprint of the application.** The less lines of code you have, the smaller the executable of your application gets.

#### **Related Reengineering Patterns.**

- The CUT AND PASTE anti-pattern [BMMM98] explains what practices lead to code duplication. The pattern discussed here focuses on the *detection* of the duplicated code.
- Patterns describing the factoring and reorganisation of code within the class hierarchy or by creating new classes. Such patterns detail how the detected clones can be merged into a single instance.

### **Motivation**

The duplication of code occurs frequently during the development phase when programmers reuse tried and tested code in a new context, but are reluctant or, due to severe time pressure unable, to invest the time necessary to generalise the existing code to be used in the old and the new context. Since duplication is an *ad hoc*/copy&paste activity more than something that is planned for, occurrences of duplication are not documented and have to be detected.

# Process

In order to detect code duplication in an unknown system, one cannot search for specific patterns. Rather, the self similarity of the system has to be discovered. Each copy is equal or similar to its clones and this similarity is revealed by comparing the entire system to itself. This comparison is on the one hand computing intensive and on the other hand produces a remarkable amount of data of possibly copied code pieces. It is therefore necessary to automatically narrow down the candidates that have to be examined in detail by a human.

Recipe. The applicability of the recipe is based on the availability of a tool for duplication detection.

- 1. Start with an automatic search for clones. The tool should create a database of all locations where code duplication possibly occurred.
- 2. Deciding on the level or size of duplication that is interesting, filters are defined that remove the uninteresting candidates.
- 3. For each clone family (i.e.  $n \ge 2$  copies of the same piece of code) that is left after the filtering step, a list of source code locations, possibly already with citations of the offending code pieces, is presented to the maintainer so s/he can decide on how to remove the duplication.

Note that the recipe in this pattern does not concern itself with the actual problems of refactoring the code.

**Difficulties.** The approaches used to compare actual pieces of code work on syntactical representations. Therefore, one cannot detect duplicated *functionality* that does not bear any syntactical resemblance.

#### Language Specific Issues.

Language dependency can stem from the parser that transforms the source code in the format that is used for comparisons by the tool. Depending on this format, the parser can be of variable complexity. For example, comparing the source code as text with only minimal transformations, e.g. removing comments and superfluous white space, only needs a very simple lexer, which keeps language dependency at a low level. Comparing abstract syntax tree of the source code, however, requires a full blown parser. The complexity of the first transformation step thus correlates directly with language dependency.

### **Tools**

Tool support is vital for applying this pattern.

- We have implemented a SMALLTALK tool called DUPLOC (see Chapter ??), which is specifically aimed at supporting a visual approach of code duplication detection. At the moment, the tool uses textual comparisons only. It allows the user to compare source code file by file, enabling him to examine the source code by clicking on the dots. Noise filtering can be done by removing uninteresting lines.
- DUP [Bak92] is a tool that detects parameterised matches and generates reports on the found matches. It can also generate scatter-plots of found matches.
- DOTPLOT [Hel95] is a tool for displaying large scatter-plots. It has been used to compare source code, but also filenames and literary texts.
- DATRIX [MLM96b] is a tool that finds similar functions by comparing vectors of source metrics.

# Discussion

This pattern is valuable to apply if your system has the symptoms identified above or if your reengineering goals belong to the set of the mentioned reengineering goals. It is also advisable, though, to apply it as a precautionary measure in the maintenance process as a *code investment* [BMMM98]. If you plan to revamp an old system, duplication detection can help to plan parts of the effort.

Moreover, if your system should be migrated from one paradigm to another one—e.g. from COBOL to an object oriented language like SMALLTALK —and you suspect duplicated code, this pattern is valuable to identify which parts of the old system have been duplicated. Assessing the similarities and differences of the parts will also improve your understanding of the systems functionalities.

The approach that has been taken in the development of DUPLOC (Chapter ??) has the following advantages:

- It is lightweight: it does not use complicated algorithms like elaborate parsing techniques.
- It is visual: the human eye is built to detect configurations and this can be fully exploited with a matrix visual representation.
- It is language independent: Since we use textual comparison, the tool is language independent to a high degree and can be used for a number of languages without a change.

**Technical.** The algorithm that is used to compare the source lines determines what level of fuzziness is allowed to recognise a match. The simplest algorithm—which compares the source lines character per character—finds only exact matches. More complicated algorithms (see for example [Bak95]) can find *parameterised matches*. Parameterised matches point out the possibility to refactor code into a parametrisable function, where exact matches emphasise more the repetitive structures in the source code.

# **Known Uses**

The pattern has been applied in biology research to detect DNA sequences [PK82]. In the context of software reengineering, the pattern has been applied to detect duplicated code in FAMOOS case-studies containing up to 1'000'000 lines of C++. It also has been applied to detect duplicated code in a COBOL system of 4 millions of line of code. The DUP tool [Bak92] has been used to investigated the source code of the X-Window system, and DATRIX has investigated multiple versions of a large telecommunications system, wading through 89 million lines of code all in all [LPM<sup>+</sup>97]. DOTPLOT [Hel95] has been used to detect similarities in man-files, literary texts and names from file systems.

Chapter 13

**Cluster: Improving Flexibility** 

# REPAIRING A BROKEN ARCHITECTURE

Author(s): Holger Bär and Oliver Ciupke

# Intent

Detect and remove dependencies between packages of a system that aren't allowed according to the designated system architecture. These architecture breaking dependencies may prohibit the exploitation of the architecture's advantages and cause unexpected effects at maintenance work.

# Applicability

This pattern is only applicable if the system to re-engineer should suit a certain architecture like *Model View Controller* (MVC), should be layered or should obey other documented restrictions concerning the dependencies between its packages.

To clarify the further discussion we note that a dependency between two classes located in different packages implies a dependency between the corresponding packages with the same direction.

#### Symptoms.

- If you are to carry out a change on the system which is supported by the system's documented architecture, e.g. replacing the top level package in a layered architecture or adding a view to the model in a MVC architecture, the effort is higher than expected. This is due to extra dependencies breaking the architecture and resulting in a cascade of changes to the rest of the system.
- Analyzing the system one encounters forbidden dependencies between packages, e. g. model classes depending on their visual representation in an MVC architecture.

#### **Reengineering Goals.**

- If the conformance to a certain architecture is proven or recovered the benefits of the architecture can be exploited, e.g. it's easy to add a new view to a model within the MVC architecture.
- Understandability: the dependency constraints of the architecture reduce the number of dependencies.

# **Motivation**

The following example describes a typical three tier architecture for business applications with a user interface, application logic layer and database layer. The architectural restriction on the dependencies between these three packages is that the user interface may depend on the application logic which may depend on the database, but nothing more.

**Initial Situation.** In our example the application logic layer implements financial transaction management and offers a service named reportTransactions (from, to) to report the transactions for a certain period of time. Figure 13.1 shows the three packages and their dependencies. Evidently there's a dependency breaking the architecture from the application logic layer to the user interface: the call new ListOutput (reportList) to create a new output window for lists offered by the user interface layer. The reason for introducing this call instead of returning the resulting list to the user interface layer might be the fear of performance penalties.

In general reasons leading to a broken architecture are:

- Altering the system without having understood the architecture.
- A system architecture which seems to have performance penalties that can be overcome by breaking the architecture.
- Favoring "quick-and-dirty" instead of "nice-and-clean".

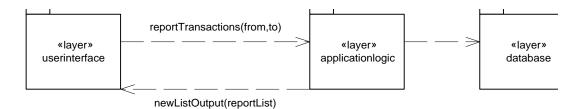


Figure 13.1: A broken architecture

**Final Situation.** The solution for the problem described above is quite straightforward. Just let the service computing the report return the transaction list to the user interface instead of itself displaying the list. But first it's not always that easy and second it's only easy after the dependency breaking the architecture has been found within the whole system.

### Structure

There is no common problem structure for breaking an architecture because architectures themselves do not have a common structure. So a target structure is missing also.

### Process

### Detection.

- 1. Analyze the actual high level dependency structure, i.e. the dependencies between packages.
- 2. Search for dependencies which are not allowed by the documented architecture. You can do this
  - **Manually:** Visualize the dependency graph with a graph layout tool and search manually for dependencies breaking the architecture.
  - **Automatically:** The process can be automated with a tool that is able to analyze and manipulate graphs or relational queries on the given data.

- (a) Set up a second graph containing the packages and the allowed dependencies between them according to the documented architecture.
- (b) Compute the set of actual dependencies minus the set of allowed dependencies. The result are the dependencies breaking the architecture.
- 3. To find architecture violations in a system that should be layered search for cycles in the dependency graph.

### Recipe.

A violating dependency exists either between two packages where no dependency is allowed or the dependency is just in the wrong direction like the one in our motivating example. In the first case there is no general solution, but in the second case the dependency can be reversed in a generic way:

- 1. Create a new abstract class with the same interface as the target of the dependency.
- 2. Replace the dependency on the target class by one on the new abstract class.
- 3. Let the target of the dependency inherit from the new abstract class. Now both the original source and target of the dependency are dependent on the abstract class.
- 4. Move the abstract class to a package where both the source and the target package may depend on it. In the case of reversing the dependency, this is the package containing the source of the dependency.

Figure 13.2 shows a broken architecture. Package P2 may depend on package P1, but P1 is not supposed to depend on P2. Actually class B depends both on classes C and D which makes P1 dependent on P2.

Figure 13.3 shows the solution to the problem. Instead of C and D, B calls now the abstract classes C\_abstr and D\_abstr. C and D inherit from their abstract counterpart and implement the methods called by B.

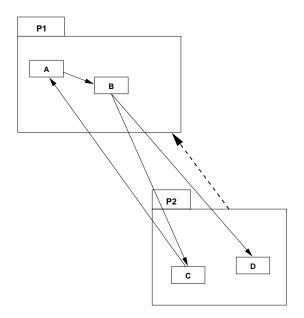


Figure 13.2: A broken architecture

There are also special solutions for similar problems like a model that has to update its various views without knowing how many views there are and of which type they are. This problem is solved by the OBSERVER pattern [GHJV95].

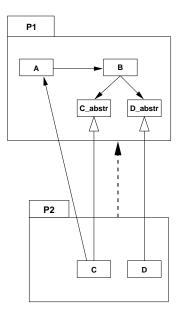


Figure 13.3: Dependencies reverted to fit the documented architecture

**Difficulties.** The new abstract class used by the source of the forbidden dependency B can be seen as an interface defined by B that supplier classes must implement. So this interface needn't contain all methods of the former target (C resp. D), but only the methods B needs.

In our example in the figures 13.2 and 13.3 there was only one class, class B, having a forbidden dependency on class C. If there is more than one class having a forbidden dependency the solution is a bit more complicated:

- Define only one interface per package that is used by all dependants on C.
- For dependants in different packages create one abstract class per package and let class C implement all of them or move the abstract class(es) into a package which every affected package may depend on.

**Language Specific Issues.** The transformation for reversing dependencies is only generally applicable for languages that allow an inheritance relation to be added to a pure abstract class (interface in JAVA ). This is the case for  $C^{++}$  and JAVA and also for SMALLTALK, because in SMALLTALK there is no need to inherit from a pure abstract class.

### **Tools**

Tools support is available for the following tasks of the detection section:

- Produce static structure graphs from source code. The tool set GOOSE contains too parsers, RETRIEVER and TABLEGEN with different advantages which can generate design information from C++ code in a format readable further tools.
- Visualizing a graph.

You can use VCG, a graph layout tool, with a great variety of hierarchical layouts or Graphlet offering a set of layout algorithms with quite different aproaches. Unfortunately Graphlet (Version 2.8) has problems with printing the graphs.

- Setting up a new graph can be done with a graph editor like Graphlet.
- Finding cycles in a graph. Execute the command

reView strongComponents < graph.gml | printCycles

of the tool set GOOSE with graph.gml replaced by your graph file.

- Computing the difference between two graphs.
- The tool set GOOSE lets you convert graphs in a relational ASCII format. Filter off any other information besides first three collumns containing the type, source and target of the relation in these ASCII files with the Unix command cut -f1-3. Then use the Unix command comm -23 followed by the two files on which the difference should be computed.

# **Known Uses**

In one of the FAMOOS case studies, there was an architecture defined with a base line framework and different products on top of this framework. When analysing the code, a class of the base line framework was found to inherit from several product classes. This kind of dependency was forbidden by the architecture definition. To repair this, an interface class was introduced from which the product classes inherited. This way, the framework was no longer dependent on the products, which made the system easier to change and decreased compile times.

A further example for an successful architecture clean up is the change in the event model of the Java Development Kit from the Version 1.0 to 1.1 (...cite). In this case the OBSERVER pattern was applied<sup>1</sup>. The observer pattern [GHJV95] is a special form of the general principle to introduce an abstract interface to decouple classes.

<sup>&</sup>lt;sup>1</sup>In the JDK, the Observer is called Listener.

# TRANSFORMING INHERITANCE INTO COM-POSITION

### Author: Benedikt Schulz

# Intent

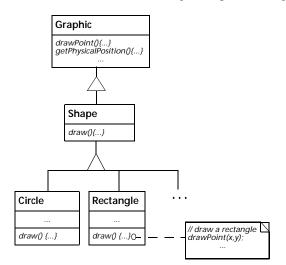
Improve the flexibility and comprehensibility of your design by transforming an inheritance relationship into a component relationship and delegating of a set of methods to this component.

# Motivation

The following example occurred in a project which aimed at visualising hydraulic data of river parts. The data was visualised in a two-dimensional diagram which changed over time. The user of the system got the impression of seeing a film because of this animation.

The most crucial part in the system concerning efficiency was the subsystem which was responsible for drawing lines on the screen: For every new frame of the animation the complete set of lines representing the data had to be redrawn.

**Initial Situation.** In the first version of the system drawing lines was handled by the GDI subsystem of the Win32s operating system. This was pretty efficient until a new requirement came into play. The customers wanted to be able to change properties of the lines like colour, thickness, style, etc. The GDI subsystem was not able to draw lines with customisable thickness in an efficient way however: The system was showing rather a slide show than a film. The initial design is depicted in Figure 13.4.



### Figure 13.4: Initial Situation

Some experiments with a new technology called DirectDraw (that is also a subsystem of the operating system) revealed its superiority and thus the project manager decided to replace GDI with DirectDraw.

This led to serious problems: Since the class responsible for drawing lines was using functionality of GDI by inheritance it was not possible just to replace it by DirectDraw. DirectDraw had a different interface and so the implementation of a lot of methods which were responsible for drawing lines had to be changed.

**Final Situation.** To avoid similar problems in the future the project manager decided not only to change the the drawing system but additionally to introduce a flexible new design which should allow for easy exchange of different drawing systems.

The new design got its flexibility mainly from one change: Instead of relying on inheritance to reuse functionality, a component relationship together with the concept of delegation was used. This means that a Shape-object no longer "knows" (directly or via inheritance) how to draw points but it rather "knows" an object which "knows" how to draw the points. Since objects can even be changed during run-time of the system the flexibility of the system was significantly improved. The final design is depicted in Figure 13.5 where new or changed entities are marked grey.

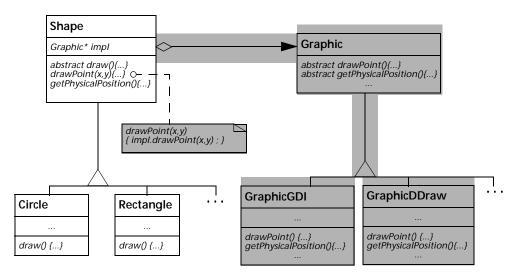


Figure 13.5: Target Structure

Some weeks after the redesign of the system it was revealed that the DirectDraw subsystem was not automatically installed on all systems running Win32s. But since the system could check whether DirectDraw was installed or not during run-time and since the drawing system was made exchangeable during run-time this new fact did not lead to any problems.

In the end the target structure is an instance of the Bridge design pattern [GHJV95]. (It was not possible to use a Singleton Graphic acting as a facade to the libraries, because Graphic is not stateless and can have different states for different Shape-objects.) The Transforming Inheritance into Composition pattern is nevertheless not equivalent to the Bridge design pattern, because it not only describes "good" target structures but rather the process of applying the Bridge design pattern to an existing object-oriented legacy system.

# Applicability

Transforming Inheritance into Compositionis applicable whenever you recognise during the reviewing of your legacy system that *you should have used* one of the following design patterns *but you have not used* them:

- Bridge [GHJV95],
- Strategy [GHJV95] or
- State [GHJV95][DA96].

All of these design patterns make use of the Objectifier design pattern [Zim95] and the technique of delegation.

The application of this pattern is difficult if the inheritance relationship is deeply nested in the hierarchy because breaking the hierarchy means that all the methods which were inherited (and this can be a large number) have to be delegated. Therefore the inheritance relationship is *not* removed in a variant of the Transforming Inheritance into Composition reengineering pattern which will be discussed later.

The reengineering pattern should not be used in the following cases:

- Inheritance *is* the appropriate modelling technique for the problem (e.g., if there is a *is-a* relationship between two classes).
- Introducing delegation would be too expensive with respect to efficiency. This has to be considered especially when the delegation takes place within a loop which is processed a lot of times.
- In statically typed languages: Clients use the two classes related via inheritance polymorphically and you do not want to change these clients.

**Symptoms.** The application of this pattern can improve your design if you encounter one of the following problems:

- For a certain problem you should have used the *Bridge*, *Strategy* or *State* design pattern but in the system you are reengineering these design patterns have not been used. You know how to use the respective design patterns when you are building a new system but you do not know how to apply them to an existing design.
  - You want to be able to change the implementation of an abstraction in a more flexible way, maybe even at run-time (*Bridge* design pattern). The actual design does not allow for this kind of flexibility.
  - You want to extend the class system with new classes which share the same interface but differ in their behaviour (*Strategy* design pattern). The actual design does not allow for this kind of flexibility.
  - You have a lot of conditional statements in your code because the behaviour of an object depends strongly on its current state. You want to get rid of these conditionals (*State* design pattern).
- The inheritance relationship was established mainly for code reuse. The code which was the reason for using inheritance now has to be changed and so you want to remove the inheritance relationship because it is no longer appropriate. You do not know how to do this without changing the functionality of the system.

**Reengineering Goals.** The goal of the Transforming Inheritance into Compositionreengineering pattern is to help software engineers to apply a design pattern relying on the Objectifier design pattern and delegation to an existing design. In particular the pattern aims at

- increasing run-time flexibility. This is achieved because after the application of the reengineering pattern you will be able to change the component during run-time.
- increasing static flexibility (configurability). This is achieved because after the application of the pattern you will be able extend the component class hierarchy independently from the abstraction.
- increasing comprehensibility. This is achieved because the reengineering pattern can remove inheritance for code reuse which is hard to understand from your system.

**Related Reengineering Patterns.** The Transforming Inheritance into Compositionreengineering pattern is related to all design patterns which rely on the Objectifier design pattern [Zim95] and delegation like

- Bridge
- Strategy
- State

## Structure

The problem structure is depicted in Figure 13.6. Transforming Inheritance into Compositionleads you to the target structure depicted in Figure 13.7

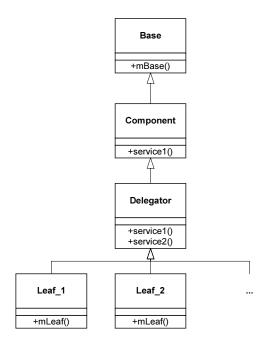


Figure 13.6: Problem Structure for the reengineering pattern

### Participants.

- **Base** is the root of the inheritance tree.
- **Component** (Graphic) is the class which gets cut out from the inheritance hierarchy to serve as a provider of certain services. The inheritance relationship to Base may remain in existence.
- **Delegator** (Shape) is the class which uses services from Component by inheritance in Figure 13.6. After application of the reengineering pattern in Figure 13.7 Delegator will make use of these services by delegation.
- Leaf\_1, Leaf\_2, ... (Circle, Rectangle, ...) are the leaves of the inheritance hierarchy
- Component\_A, Component\_B, ... (GraphicGDI, GraphicDDraw, ...) are the subclasses of Component implementing the services of their super-class in different ways..

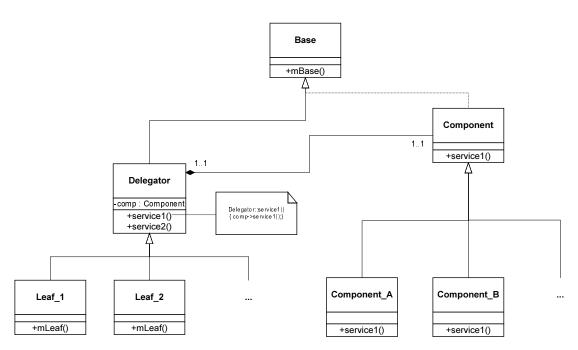


Figure 13.7: Target Structure for the reengineering pattern

### Collaborations.

- Delegator makes use of service1 (drawPoint) provided by Component. This is done
  - in the problem structure by executing inherited methods from Component whereas
  - in the target structure the execution of these methods is *delegated* to Component.

#### **Consequences.**

- Positive benefits
  - Transforming Inheritance into Compositionsolves an important and basic reengineering problem and the application of the reengineering pattern allows for the introduction of several known design patterns [GHJV95].
  - Since abstraction and implementation are separated, changing the implementation does not require recompilation but only rebinding of the system.
  - The implementors of servicel can be designed to form a separate inheritance tree. (This is suggested by the class ComponentA in Figure 13.7.) This is impossible before the application of the reengineering pattern.
- Negative liabilities
  - The execution of service1 provided by Delegator will take longer in the target structure because it has to be delegated. This may be critical if service1 is needed a lot of times.
  - The target structure is slightly more difficult to implement since the attribute of Delegator named comp has to be initialised whenever a new instance of Delegator is created and destroyed whenever that instance is deleted.

## Process

The process mainly relies on the idea of combining the approach of considering design patterns as operators [Zim97] (rather than building blocks) and the refactoring approach presented in [JO93]. This idea is presented and discussed in detail in [SGMZ98a].

### **Detection.**

Since violations against flexibility issues can only be detected if you know where flexibility is needed and which kind of flexibility (e.g., run-time flexibility, configurability) is needed, algorithmic detection is difficult. However, you can

- ask people who designed and implemented the system if there is a case where they wanted to be able to change the implementation of an interface at run-time and this was not possible.
- look for methods with a large amount of conditional statements. The behaviour of an object may depend strongly on its internal state (Type Check Elimination within a Provider Hierarchy).
- look for two classes, one inheriting from the other, which are never used polymorphically. This means that a variable declared as super-class is never used for an instance of the subclass.

**Recipe.** In this section we show how to apply tTransforming Inheritance into Composition and what kind of reengineering operations have to be applied. If we name entities (like classes, methods and attributes) we refer to the participants of the problem structure depicted in Figure 13.6 and the target structure depicted in Figure 13.7.

- 1. Create a new attribute comp of Component in the class Delegator. Change the constructor method of Delegator so that it initialises the attribute comp with a new instance of Component. If you plan to add several subclasses of Component later on (you should do so!) than add a new formal argument to the constructor method of Delegator which will serve as an indicator of which concrete subclass of Component to use.
- 2. Copy all the signatures of the methods from Component which are visible to Delegator to Delegator. For each added method add an implementation which delegates the execution of the method to the corresponding method of Component. For an example, see the implementation of Delegator:service1() in Figure 13.7.
- 3. Remove the inheritance relationship between Component and Delegator. Caution: In statically typed languages you will not be able to use an instance of Delegator polymorphically as an instance of Component after this step. In particular it is not possible any more to cast instances of Delegator to Component.

### Difficulties.

If you decide to introduce an additional formal parameter to the constructor of Delegator then every piece of code that creates an instance of Delegator has to be changed. In languages which support default values for formal parameters this problem can be resolved by defining an appropriate default value (e.g., Component if this class is not made abstract).

If there is no way to avoid polymorphism between Delegator and Component but you still have strong reasons to apply Transforming Inheritance into Composition and you are using a statically typed language, you can omit removing the inheritance relationship between Component and Delegator. You should be aware of the fact, that you might have the following problem: The class Component has two parts:

- One part of the methods represents set of utility services. You made Delegator inherit from Component because you wanted to be able to use these services without re-implementing them.
- The other part of the methods represents the *real* interface of Delegator. You made Delegator inherit from Component because you wanted to establish an *is-a* relationship between Delegator and Component to be able to use instances of both classes polymorphically.

In this case consider splitting the Component class into two separate classes.

#### Language Specific Issues.

- In C<sup>++</sup> you should implement the attribute comp as a pointer. Otherwise you will not be able to use polymorphism for the inheritance tree with root Component.
- In dynamically typed languages like SMALLTALK it is not necessary that two classes are related via an inheritance link to use them polymorphically. This means, for example, that you can still use instances of Component and Delegator together in one container object.

### Discussion

Since the detection of the problem structure is far away from being an algorithmic, tool supported process, you should not explicitly look for this problem structure. But since software development is an iterative process you will find the problem structure while trying to extend or modify your system. Once you have found the problem structure in your code, you should strongly consider the application of Transforming Inheritance into Composition.

The relevance of this reengineering pattern is high: In a lot of companies which were early adopters of the object-oriented paradigm, the maturity of the software engineers concerning object-oriented technology was low. This resulted in an overuse of inheritance, mainly for code reuse. These software defects can be removed by the application of the reengineering pattern.

The concept of delegation and the Objectifier design pattern [Zim95] are the fundamentals of this reengineering pattern and the resulting target structure is closely related to the Bridge, Strategy and State design patterns [GHJV95]. A good understanding of these design patterns helps to use the reengineering pattern.

## Tools

The detection of pairs of classes which are never used polymorphically can be done with the tool-set *Goose* [BC98][Ciu99]. *Goose* can not only detect missing polymorphism but a lot of other design defects which occur in object-oriented systems.

Since the application of the reengineering pattern relies on the application of refactorings [Opd92] you can use every tool which supports this technique, such as the *Refactoring Browser* [RBJ97b] for SMALLTALK, which is the most advanced refactoring tool. The Refactoring Browser is described and available for free at http://st-www.cs.uiuc.edu/~brant/Refactory/.

For a subset of  $C^{++}$  we implemented a prototype to support refactorings. This tool is called *RefaC*++ and described in [Moh98]. *RefaC*++ can perform a subset of the refactorings presented in [Opd92] and can also apply the Bridge design pattern automatically.

# **Known Uses**

Transforming Inheritance into Compositionhas been applied in the following known cases:

- The reengineering pattern was applied with success in the project described in the motivation section. It was possible to increase the flexibility of the system so that the new requirement (DirectDraw not available on every Win32s installation) could be fulfilled without problems.
- We are analysing and flexibilising a graphical information system for a German middle-sized enterprise. We found several design flaws which have been corrected by applying this reengineering pattern.
- [RJ96] describes how frameworks evolve. In the White-box Framework design pattern [RJ96] the engineer is encouraged to use inheritance for reuse because it is easier to understand and to reuse. In later stages of the framework development inheritance has to be replaced by polymorphic composition.

# DISTRIBUTE RESPONSIBILITIES

### Author(s): Holger Bär and Oliver Ciupke

## Intent

Distribute the responsibilities equally among the classes of an object-oriented system to prevent large, hardly maintainable and reusable classes.

# Applicability

A responsibility is a description of a service offered by a class. It is fulfilled by a set of publicly accessible methods.

**Symptoms.** If the responsibilities aren't distributed among the classes, there will be one or more classes incorporating a lot of responsibilities. Such classes, called *multiple responsible classes* (MRC) from now on, result in the following symptoms.

- If you ask for the responsibilities of a MRC, you get long and unclear answers.
- The MRC is used by other classes for different purposes (low level MRC).
- The MRC uses a lot of classes (high level MRC or *manager class*).
- Functional enhancements somewhere in the system often require changes in one of the high level MRCs.
- A MRC is mostly large in lines of code and number of methods, because many responsibilities result in many methods resulting in many lines of code for concrete classes.
- High level MRCs can hardly be reused because too many design decisions of the specific application are coded into them.
- Maintenance work on MRCs is hard, because there is no boundary between the different responsibilities, so that it's unclear where to change the class for a certain maintenance action and which the effects of the change are.

### **Reengineering Goals.**

- Understandability: classes with many responsibilities are hard to understand, because the responsibilities are mixed together, i. e one can't identify the individual responsibilities and understand their implementation and collaboration with other classes in isolation.
- Flexibility: classes with few responsibilities allow fine grained adoptions by subclassing or replacing a class.
- Reusability: classes are normally reused as a whole. Therefore it's unlikely that MRCs are reused because the particular combination of responsibilities needed in its original application is unlikely to occur in another one.

#### **Related Reengineering Patterns.**

- large classes
- large methods
- structural programming Classic structural programming principles applied to an OO-language often lead to one central manager class operating on several dumb data classes.

# Motivation

**Initial Situation.** The UML diagram in Figure 13.8 shows a manager class, AccountManager, together with two passive classes, AccountData and BarChart. The responsibilities of the manager class are

- 1. to process user input in OnCalculateSummary,
- 2. to do the summary calculation using Get...Transaction methods to query the transaction data from class AccountData,
- 3. and to present the results with the help of class Screen.

So the manager class has three responsibilities and implements nearly the whole functionality. Major design decisions like how the summary is calculated and presented and the reactions on the user input are hard coded in this class thus making it hardly reusable.

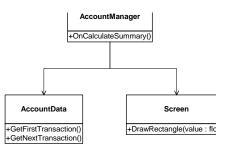


Figure 13.8: Example of a manager class with two passive classes.

**Final Situation.** We can distribute the three responsibilities of the manager class among three classes: UserInteraction, Account and BarChart. In this design all classes besides UserInteraction have a high potential for reuse.

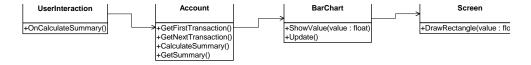


Figure 13.9: The improved example with distributed responsibilities.

# Structure

The structure of the problem and the target structure differ both between high level MRCs and low level MRCs.

### High level MRC problem structure.

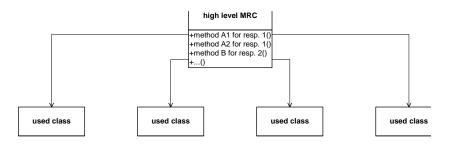


Figure 13.10: Problem structure of a high level MRC.

Participants. The high level MRC shows a broad interface with a set of methods per responsibility.

Collaborations. High level MRCs often use many other classes to fulfill their numerous responsibilities.

Low level MRC problem structure.

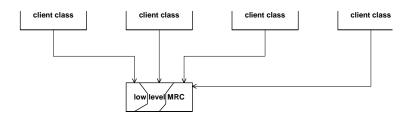


Figure 13.11: Problem structure of a low level MRC.

Collaborations. The various clients use different responsibilities of the low level MRC.

#### High level MRC target structure.

For high level MRCs there is no general target structure. The goal is to distribute the responsibilities. Good candidates for receiving responsibilities are the used classes. But sometimes it's necessary to define a new class like BarChart of the motivating example. The manager class itself will be reduced in size (lines of code, methods) or will disappear completely.

### Low level MRC target structure.

**Participants.** The low level MRC has been split into several classes according to the responsibilities of the MRC and the parts used by the client classes.

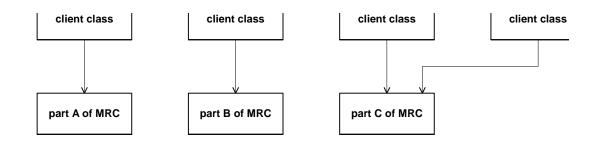


Figure 13.12: Target structure of a low level MRC.

**Collaborations.** The client classes of the low level MRC now only use those parts of the MRC they actually need.

### **Consequences.**

- The responsibilities of the high level MRC are distributed moving the code closer to the data it works on.
- In case the manager class disappears completely there is no central control any more. The design has made a step towards autonomous interacting objects.
- The remainder of the high level MRC and the parts of the low level MRC are smaller, easier to understand and exhibit more potential for reuse.
- Instead of many dependencies on one low level MRC the application of this pattern leads to a set of classes each with a lower number of dependents.
- The smaller "part classes" of the low level MRC are more stable than the original class simply because they encapsulate less design decisions. So together with the previous topic the compilation times after changes to the system will be reduced.
- In both cases the number of classes may increase.
- The distribution of responsibilities may affect the efficiency of the system.

## **Process**

#### **Detection.**

- MRCs are normally the largest classes in a system, program or package both in lines of code and in number of methods.
- To find high level MRCs search for classes with manager, man, driver, initiator and so on in their name.
- Classes that use many other classes are also good candidates for high level MRCs. They can be found by looking for classes with high values for coupling metrics like the CBO metric [CK94].
- Low level MRCs are used for different purposes. So the implementation of their responsibilities are likely to not communicate with each other. Therefore these classes will often exhibit low cohesion. There are numerous cohesion metrics, e. g. the TCC metric [BK95].

The conjunction of the size and coupling criteria and optionally the name criterion should produce satisfying results for the detection of high level MRCs as should the conjunction of the size and cohesion criteria for low level MRCs.

### Recipe.

- Search for candidate MRCs as described in the detection section above. The next steps depend on whether you have detected a high level or a low level MRC.
- High level MRCs

Try to distribute the responsibilities of the manager class to other classes. Good candidates for receiving responsibilities are the classes used by the manager class. It may be necessary to define a new class like **BarChart** of the motivating example. The manager class itself will be reduced in size (lines of code, methods) or will disappear completely.

- Low level MRCs
  - 1. Determine the parts of the low level MRC.

There are two ways to determine the parts — one considers the use of the class by its clients, the other one examines the class' internal structure:

- (a) Analyse the use of the class by its clients. Note for each type of client the features of the MRC (methods and public attributes) it uses. Find a partition of the feature set so that each client uses only one or few parts.
- (b) Although cohesion metrics indicate whether a class could and should be split, they do not directly indicate where to split the class. You can get good suggestions for splitting by computing the minimum cut on the undirected graph containing all methods and attributes of the class as nodes and all method calls and variable accesses within the class as edges<sup>2</sup>. The minimum cut algorithm computes such a partition of the graph in two sets of nodes that the number of edges from one set to the other is minimised. Splitting a class according to this partition leads to two classes with minimal communication between them. Of course this splitting step can be applied to the two sets recursively until the resulting classes are small enough.
- 2. You may use one of the mentioned partitions of the MRC or a combination of both to split the MRC. In cases where there is no optimal partition (e.g. client uses more than one part or there is more than zero communication between the parts) the partition often needs some manual fine tuning to end up with a set of reasonable classes.
- 3. Split the MRC according to the partition and reorganise its context.

### Difficulties.

- The detection of MRCs is not very precise, especially for low level MRCs. The key point in detecting a low level MRC is to recognise that it is used by different types of clients for different purposes.
- Determining the partition of the MRC can hardly be done fully automatically because the parts must be reasonable classes.

### **Tools**

Several prototype tools have been developed within the FAMOOS project which can help to detect and to solve this problem.

<sup>&</sup>lt;sup>2</sup>All edges have weight 1; parallel edges are allowed for modelling multiple calls and accesses

- A visualisation of the static system structure at the right level of abstraction, e.g. with the tools within the tool set GOOSE can help detecting central classes or subsystems.
- TABLEGEN computes the TCC cohesion metric and also coupling metrics.
- The computation of minimum cuts to determine the partition of a MRC can be done with REVIEW, a tool also developed within the FAMOOS project.
- GOOSE' relational representation of design information enables to search for classes with a high out-degree of usage. The classes with the highest out-degrees are good candidates for being high level MRCs.

# **Known Uses**

During restructuring in one of the FAMOOS case studies, there was a big class found which incorporated responsibilities for several different products. This class was split into several pieces to make the program more flexible with respect to frequent changes [Rit98].

# USE TYPE INFERENCE

### **Author: Markus Bauer**

## Intent

It is hard to understand the structure and the workings of a software system written in a dynamically typed language because of the lack of type declarations. Therefore add type annotations to the program code which document the system and which can additionally be used by sophisticated reengineering tools.

# Applicability

Apply this pattern when reengineering systems that are written in *Smalltalk* or in a similar, dynamically typed programming language, where you have only limited knowledge about the system. Typical situations could be:

- You have to maintain and/or modify the software system, but you have only limited knowledge about its inner workings. You are interested to learn, which types of objects of the system are manipulated by some code your are working on, but this is difficult since you do not have type declarations in your source code that provide you with that information.
- You want to support a reengineering task by some tools, but these tools rely on type information for the system's variables and methods. Most reengineering tools rely on such type information. Examples include (but are not limited to) the *Smalltalk Refactoring Browser* [RBJ97b]<sup>3</sup> or tools that calculate software product metrics (like those described in [CK94]).
- You want to reengineer or rewrite the system using a statically typed programming language, but to achieve this, you need appropriate type declarations for the system's variables and methods.

# Problem

In dynamically typed systems, the lack of static type information (i.e. the lack of type declarations for variables and method signatures) makes some reengineering tasks difficult or impossible, since such type information usually represents prominent parts of a system's semantics.

# Motivation

Consider some code fragments<sup>4</sup> for a dynamically typed application that manipulates drawings. Such an application might have a class Container for storing some objects. Figure 13.13 shows a method add that is used to add objects to the container.

For reengineering purposes we might be interested in an answer to the following question: What kind of objects can be stored in the container, that is, of what types are the objects, that are passed as arguments to the add-Method?

<sup>&</sup>lt;sup>3</sup>The current implementation of the Refactoring Browser does not infer precise types for the system's entities though, it relies on (unprecise) heuristics instead.

 $<sup>^{4}</sup>$ We present code examples in a syntax close to Java. Since we deal with dynamically typed code, we just omit Java's type declarations.

```
add: anObject
  contents add: anObject.
  anObject draw.
  "..."
```

Figure 13.13: Method add in class Container.

### Forces

- To learn about the types of objects that are manipulated by some code you are looking at, you might consider manually tracing the execution of your code and guess what's going on in your system, but for larger systems, this is an infeasible and error-prone task.
- You could also try to capture that information by looking at method and variable names, but in many legacy systems naming conventions do not exist or do not provide enough information about the object types and the manipulations that are made with them (see our example above). Even worse, you can't be sure that the names do not lead you to wrong conclusions.
- To migrate from a dynamically typed language to a statically typed language, you could apply approaches that do not rely on type information, like those proposed for the translation of Smalltalk applications to Java in [EK98]. These approaches simulate Smalltalk's dynamic type system in Java. The resulting code, however, is not authentic Java code and is hard to understand and maintain. Such code additionally has the usual shortcomings of untyped code: it is not type safe.

## Solution

Find out what types the variables and method parameters have and put this information into the source code, using type annotations or comments.

In more detail:

- 1. Perform a program analysis of your dynamically typed object oriented legacy system.
- 2. Use the results of the program analysis to determine type information for the program's variables, including global and local variables, parameters and return values of methods. Based on this type information, add type annotations to the program's source code.
- 3. Use these type annotations to understand how your legacy system works or as additional semantic information to more sophisticated reengineering tools.

This technique is called *type inference*, because you infer the type of an object at a certain place in the code by tracing its way from its creation to the current place.

If we can enrich the code of our example application with type annotations (see figure 13.14) by using the techniques described below we can easily find an answer to the question we asked above: Our Container holds points, lines, splines, ..., so it has obviously something to do with some geometrical shapes that make up a drawing.

We learn from this example that type annotations like those given in figure 13.14 make code much easier to understand and that they contain valuable information about the inner workings of a system.

```
add: anObject
  " {Container} × {Point, Line, Spline,...} → {} "
  contents add: anObject.
  anObject draw
  "..."
```

Figure 13.14: Method add annotated with type information.

# Process

Type inference usually can't be done manually for reasonable large and complex applications. Therefore, we have to automate the task of computing type information for variables and method signatures.

To implement a tool or other means to get the information, we observe that during the runtime of the system, type information propagates through the system's expressions and statements: Upon creation, each object has a certain type assigned to it, and this type information is spread to all expressions and statements (including variable and method parameter expressions), that do some operations with the object. Thus, to infer types for the variables and methods of the system, we need to inspect object creations and the data flow through the system.

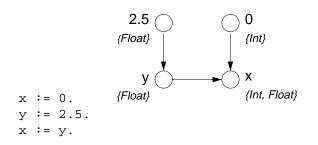
Basically we can do this in two ways: We either can execute the application and collect the type information we are interested in during its runtime (*dynamic type inference*) or we can use static program analysis techniques (*static type inference*) to analyse the applications source code and compute how the type information flows through the application's expressions. We will cover both approaches in some more detail below.

**Dynamic type inference.** With dynamic type inference, we modify the application or its runtime environment, to have it record the runtime type information for us.

- Determine the most common execution paths through your program, that is, determine the most common usage scenarios of your legacy system. In some cases you might be able to use already existing testing scenarios for this. In other cases, determining these common usage scenarios might be difficult, especially if you don't know much about the system.
- Instrument the code with instructions that record the data flow through your system and that collect the runtime types of the system's variables. [RBFDD98] describes how to modify the runtime libraries of a Smalltalk environment to achieve this with only minor changes to the application's code.
- 3. Run the system and have it execute the most common usage scenarios you collected in step 1.
- 4. Use the recorded runtime type information to put type annotations into the source code.

**Static type inference.** With static type inference, we need a tool that reads in the complete source code of the application and analyses it to construct a data flow graph. This is done by representing the application's expressions as nodes in the graph, and by modelling the dependencies between them as edges. The dependencies that are taken into account to construct the data flow graph are given by the following rules:

1. An *assignment* var := expr generates a data flow from the right hand side expression expr to the variable var on the left hand side.





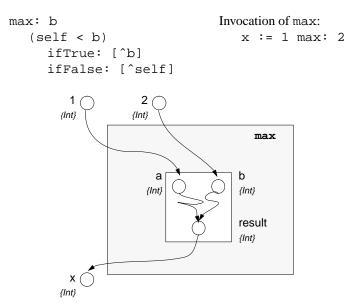


Figure 13.16: Data flow across method boundaries

- 2. A *variable access* generates a data flow from the variable being accessed to the surrounding expression.
- 3. A *method invocation* generates a data flow from the actual argument expressions to the formal arguments of the invoked method, and from the result of the invoked method to the invoking expression.

A data flow graph for a short piece of code is shown in figure 13.15.

For each node the tool then tries to compute the set of classes the corresponding expression can hold instances of. It starts by determining type information for the program's literal expressions and object creation statements (which are represented as source nodes in the graph) and moves that information along the edges through the graph. Each node then carries the union set of all type information of its predecessors. In figure 13.15, for example, the node for x carries the type information {*Int, Float*}, since it depends on the type information of the nodes for y and 0.

Some subtle problems arise, whenever method invocations cause data flows across method boundaries (as given by rule 3). Such a case is shown in figure 13.16.

There are some well proven techniques to allow for an analysis which keeps track of these inter-method data flows in an efficient and practicable way. One of these is Agesen's Cross Product Algorithm [Age95]<sup>5</sup>.

<sup>&</sup>lt;sup>5</sup>There are other algorithms that also allow the tracking of data flow across method boundaries, for example [PS91], [OPS92], [PC94], but Agesen's algorithm is superior to most of these, because it is easy to understand and computes precise type information

The basic idea is to create separate sub graphs for each method and link all those subgraphs together in an appropriate and efficient way.

After the graph has been complexly built up and all type information has been propagated through it, the type information associated with the graph's nodes can be used to annotate the source code of the application.

### Discussion

A problem of using type inference to reveal some information about a legacy system arises from the fact that we analyse the data flow through an application. To make our approach work, we have to analyse the complete source code of an executable application (including libraries), or, if we are using dynamic type inference, we have to execute an adapted version of the system. This might be a problem in some cases when parts of the source code are not available and/or a runnable version of the system cannot be produced. Furthermore, frameworks and class libraries cannot be analysed without application code using or instantiating them. Then, however, the inferred types are only valid in the specialised context of the particular application.

Static type inference algorithms usually have to overcome some difficulties: static analysis is complex and the results are often unprecise. Agesen's static type inference algorithm, as sketched above, addresses these difficulties in an appropriate way<sup>6</sup>. However, since the algorithm is very complicated it is difficult to implement it in a correct way and produce a reliable tool out of it. This is an issue, if you can't use one of the already existing tools (see for example [Li98]).

However, once a tool for performing such an analysis has been built, it can be used on other reengineering projects as well and then it quickly pays of its rather high development costs.

Dynamic type inference has serious limitations when being applied to larger systems: You have to ensure that the most important parts of the system are covered by the analysis in a sufficient way, which might not be feasible for larger systems if you do not have test cases or usage scenarios available.

## **Related Reengineering Patterns**

Type annotations document the inner workings of a legacy system. We can therefore see type inference as a technique to improve your knowledge about the legacy system. Thus, this pattern relates with all other reengineering patterns that describe *reverse engineering techniques*, i.e. analyses of the source code of legacy systems to extract additional semantic information and improve the understanding of the systems.

## **Known Uses**

ObjectShare has used type annotations (like those that can be computed by applying this pattern) to document large parts of the source code to the *Visualworks Smalltalk* environment. This emphasises that type annotations are of great help understanding source code.

The GOOSE tool set (and related tools) that support the reengineering of C++ applications by visualising software structures [Ciu97], checking design heuristics [BC98] and calculating software metrics [Mar97] can analyse Smalltalk applications after type inference is used and the source code is enriched with type annotations.

in a very efficient way[Age94].

<sup>&</sup>lt;sup>6</sup>A detailed discussion of the algorithm, especially regarding complexity and precision can be found in [Age95] or in [Bau98].

The University of Stuttgart, Germany, has developed a tool called *Smalltalk Explorer* which is used to explore existing Smalltalk applications. It heavily relies on the type inference algorithm presented here. Type annotations are used to allow for an easy navigation through unknown Smalltalk code by documenting which classes are manipulating which other classes and by introducing hyperlinks between them[Li98].

The type inference algorithm is also used to facilitate a mostly automatic translation of dynamically typed Smalltalk applications into statically typed Java applications [Bau98]. Since most of Smalltalk's concepts can be mapped upon suitable Java concepts the most prominent issue is to infer appropriate static types for the resulting Java code. This is done by computing type annotations (as described above) and transforming them into type declarations. In more detail, to map a type annotation to a type declaration, a class must be found (or created by refactorings), that is a common abstraction to all classes included in the type annotation.

## Appendix A

# Glossary

#### **Class versioning.**

Each change to a class results in a new, distinct class definition; each object belongs to one class version, which is thus a snapshot of the class definition at a certain point in the life of a system [BH89].

#### Conversion.

The physical storage format of objects is transformed to match a different class definition [LH90].

#### Design Pattern.

A proven design that describes the core of a solution to a problem which occurs over and over again in OO software design, together with its range of applicability. The solution usually has developed and evolved over time [GHJV95].

#### Design recovery.

A subset of reverse engineering in which domain knowledge, external information, and deduction or fuzzy reasoning are added to the observations of the the subject system. The objective of design recovery is to identify meaningful higher-level abstractions beyond those obtained directly by examining the system itself [CCI90].

#### Filtering and screening.

Objects are wrapped with exception handlers that hide differences between different versions of the same class [SZ86].

#### Forward engineering.

The traditional process of moving from high-level abstractions and logical, implementation-independent designs to the physical implementation of a system [CCI90].

#### Framework.

A framework is an abstract object-oriented design together with a library of standard software components (abstract classes and templates as well as concrete components) that can be specialised, instantiated and combined to construct a number of systems with similar properties for a specific application domain [Cas96] [JF88].

#### Global reorganisation of hierarchies.

Algorithms put inheritance hierarchies in a normal form that minimises the number of classes and relationships and suppresses redundant definitions [Che96] [LBSL91].

#### Hierarchy maintenance.

Guidelines and semi-automatic algorithms are provided to reorganise inheritance hierarchies for improved reusability [Put94].

#### Incremental reorganisation of hierarchies.

Algorithms analyse and reorganise subclassing relationships to eliminate the need for redefinitions in inheritance relationships and to factor common functionality in a hierarchy [Cas92] [DDHL96].

#### Law of Demeter.

Reference and calling dependencies between methods and variables are changed to follow a modular programming style [LHR88].

#### Legacy System.

A system is called legacy if it exhibits the following properties: It is a production system carrying out business-critical tasks. It has been developed with older technology, or with older versions of an existing technology. It can no longer be easily adapted to meet changing requirements.

#### Method factorisation.

Code fragments common to several methods are extracted and put in separate methods to maximise code sharing [OJ90].

#### Metrics-based analysis.

A set of quality statistics on oo code is computed and then used to detect design problems and guide re-engineering activities.

#### **Object-oriented views.**

An abstraction layer is inserted between subsystems in an oo design to limit the scope of changes to single subsystems and to map different versions of subsystems to each other [Bra92] [RR95].

#### Pattern restructuring.

Modification primitives on oo programs structured according to specific kinds of patterns and architectures are formally specified to ensure the preservation of behaviour [HS96b].

#### Pattern-directed re-engineering.

Standard software structures serve to document and analyse oo applications, and serve as targets structures for re-engineering [Yel96] [Zim95].

#### **Redocumentation.**

A form of restructuring where the resulting semantically-equivalent representation is an alternative view intended for a human audience [CCI90].

#### **Reengineering.**

The examination and alteration of a subject system to reconstitute it in a new form and the subsequent implementation of the new form [CCI90].

#### **Refactoring.**

Frequent high-level re-engineering operations are identified and specified to preserve class behaviour across modifications [JO93] [OJ90].

#### **Restructuring.**

A transformation from one form of representation to another at the same relative level of abstraction. The new representation is meant to preserve the semantics and the external behaviour of the original [CCI90].

#### **Reverse engineering.**

The process of analysing a subject system with two goals in mind:

- 1. to identify the system's components and their interrelationships; and,
- 2. to create representations of the system in another form or at a higher level of abstraction [CCI90].

#### Schema evolution.

A set of operations specifies the kinds of high-level revisions to an oo design that occur during conceptual modelling [LM96].

#### Schema modification primitives.

A set of elementary modification operations, all specified so as to preserve basic integrity constraints, suffices to define any other more complex modification operation [BKKK87] [Zic92].

#### Software maintenance.

Modification of a software product after delivery to correct faults, to improve performance or other attributes, or to adapt the product to a changed environment [IEE83].

#### Tailoring and excuses.

Language mechanisms allow to redefine inherited properties and to accommodate exceptions in specialisation hierarchies [Bor88].

#### Transposed files.

By storing object variables in different tables, adding, updating and deleting variables is possible without reformatting records [ALC91].

#### Visual analysis of oo software.

Static and dynamic properties of oo programs are represented graphically to identify design and implementation problems [PHKV93] [SSC96].

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