

# Redesigning with Traits: the Nile Stream trait-based Library

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**Abstract.** Recently, traits have been proposed as a single inheritance backward compatible solution in which the composing entity has the control over the trait composition. Traits are fine-grained units used to compose classes, while avoiding many of the problems of multiple inheritance and mixin-based approaches.

To evaluate the expressiveness of traits, some hierarchies were *refactored*, showing code reuse. However, such large refactorings, while valuable, may not be facing all the problems, since the hierarchies were previously expressed within single inheritance and following certain patterns. We wanted to evaluate how traits enable reuse, and what problems could be encountered when building a library using traits *from scratch*, taking into account that traits are units of reuse. This paper presents our work on designing a new stream library named Nile. We present the reuse that we attained using traits, and the problems we encountered.

**Keywords.** Object-Oriented Programming, Inheritance, Refactoring, Traits, Code Reuse, Smalltalk

## 1 Introduction

Multiple inheritance has been the focus of a large amount of work and research efforts. Recently, traits proposed a solution in which the composite entity has the control and which can be flattened away, *i.e.*, traits do not affect the runtime semantics [1, 2]. Traits are fine-grained units that can be used to compose classes. Like any solution to multiple inheritance, the design of traits is the result of a set of trade-offs. Traits favor simplicity and fine-grained composition. Traits are meant for single inheritance languages. Trait composition conflicts are automatically detected but the composer has the control to resolve these conflicts explicitly. Traits claim to avoid many of the problems of multiple inheritance and mixin-based approaches that mainly favor linearization where conflicts never arise explicitly and are solved implicitly by ordering.

Note that there exist different trait models. In the original trait model, *Stateless traits* [1, 2], traits only define methods, but not instance variables. *Stateful traits* [3] extends this model and lets traits also define state. *Freezable traits* [4] extend stateless traits with a visibility mechanism. In the context of this paper when we use *trait* we mean *Stateless trait*.

Previous research evaluated the usefulness of traits by refactoring the Smalltalk collection and stream libraries, which showed up to 12% gain in terms of code reuse [5]. Other research tried to semi-automatically identify traits in existing libraries [6]. While these are valuable results, they are all refactoring scenarios that investigated the applicability of traits using *existing* systems as input. Usability and reuse of traits when developing a *new* system has not been assessed. Implementing a stream library from scratch is an important experience to test the expressiveness of traits. By doing so we may face problems that may have been hidden in previous experiences and also face a large scheme of trait composition problems.

The goal of this paper is to experimentally verify the original claims of simplicity and reuse of traits in the context of a forward engineering scenario. More specifically, our experiments want to get answers to the following questions that quickly arise when using traits in practice:

- Trait granularity. We want to assess the granularity of traits that maximize their reusability and composition.
- Trait reusability. We want to understand how much code can be reused.
- Can we define traits as composable building units?
- Can we identify guideline to assess when trait composition should be preferred over inheritance?
- To what extent can we fix the problems identified in the current stream hierarchy?
- What trait limits and problems do we encounter?
- Does the use of trait imply an execution cost?

Our approach is based on designing and implementing a non-trivial library from scratch using traits. We decided to build a stream collection library (called *Nile*) that follows the ANSI Smalltalk standard [7] yet remains compatible with the current Smalltalk implementations. The choice for a stream library was motivated by a number of reasons:

- streams exhibit problems linked to the fact that they are naturally modeled using multiple inheritance. In presence of single inheritance the implementors are reduced to duplicated code and other tricks such as canceling methods;
- N. Schärli [5] and A. Lienhard [6] already refactored the Stream library using traits so we can compare with their results;
- streams are an important abstraction of computer language libraries;
- several constraints are imposed by the ANSI Smalltalk standard and the need to remain usable in existing Smalltalk dialects.

Nile is structured around three core traits and a set of libraries. During the definition of the libraries, the core traits proved to have a good granularity: it was easy to obtain each desired functionality composition using the adequate part of the core. Nile has 18% less methods and 15% less bytecodes than the corresponding Squeak collection-based stream library. Moreover, Nile has neither canceled method nor method implemented too high in the hierarchy. There are only three overrides compared to the fourteen of Squeak.

The contributions of the paper are: (1) the design of *Nile*, a new stream library made of composable units, (2) the assessment that traits are good building units for defining libraries and that they enable clean design and reuse through composability, and (3) the identification of problems when using the traits.

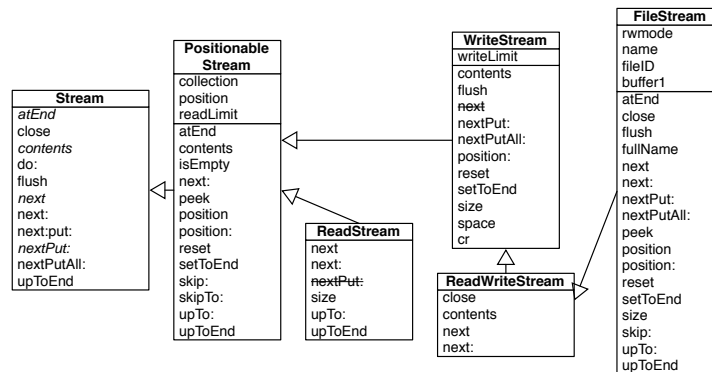
We start by presenting the existing Squeak Stream hierarchy limits and the ANSI Smalltalk standard protocols (Section 2). Section 3 presents an overview of Nile and the core of the library around its three most important traits. Section 4 and Section 5 detail the implementation of the collection-based and file-based stream libraries, respectively. Two other libraries will be presented in Section 6. Section 7 compares our approach with the one of N. Schärli [5]. It analyses the reuse offered by traits as well as performance issues and optimization solutions. Finally, Section 8 presents the problems we identify due to the use of traits.

## 2 Analyses

In this section, we analyze the existing stream hierarchy of Squeak the open-source Smalltalk [8]. We highlight the key problems and present the ANSI Smalltalk standard.

### 2.1 Analysis of the Squeak stream hierarchy

Squeak [8], like all Smalltalk environments, has its own implementation of a stream hierarchy. Figure 1 presents the core of this implementation, which is solely based on single inheritance and does not use traits. Note that most Smalltalk dialects reimplemented streams and therefore have similar yet different implementation. For example, even though Squeak and VisualWorks are both direct descendants from the original Smalltalk-80, their stream hierarchies are different since the one in VisualWorks was completely reimplemented.



**Fig. 1.** The Squeak core Stream hierarchy. Only the most important methods are shown.

The existing single-inheritance implementation has different problems that we detail.

*Methods implemented too high in the hierarchy.* A common technique to avoid duplicating code consists in implementing a method in the topmost common superclass of all classes which need this method. Even if efficient, this technique pollutes the interface of classes which do not want this method. For example, `Stream` defines `nextPutAll:` which calls `nextPut:`

```
Stream>>nextPutAll: aCollection
  "Append the elements of aCollection to the sequence of objects
  accessible by the receiver. Answer aCollection."

  aCollection do: [:v | self nextPut: v].
  ^ aCollection.
```

The method `nextPutAll:` writes all elements of the parameter `aCollection` to the stream by iterating over the collection and calling `nextPut:` for each element. The method `nextPut:` is abstract and must be implemented in subclasses, and even if `Stream` defines methods to write to the stream, some subclasses are used for read-only purposes, like `ReadStream`. Those classes must then cancel explicitly the methods they don't want.<sup>3</sup> This approach, even if it was probably the best available solution at the time of the first implementation, has some drawbacks. First of all the class `Stream` and its subclasses are polluted with a number of methods that are not available in the end. This complicates the task of understanding the hierarchy and extending it. It also makes it more difficult to add new subclasses. To add a new subclass, a developer must analyze all of the methods implemented in the superclasses and cancel all unwanted ones.

*Unused superclass state.* The class `FileStream` is a subclass of `ReadWriteStream` and an indirect subclass of `PositionableStream` which is explicitly made to stream over collections (see Figure 1). Then, the instance variables `collection`, `position` and `readLimit` inherited from the `PositionableStream` and `writeLimit` inherited from `WriteStream` are not used for `FileStream` and all its subclasses.

*Simulating multiple inheritance by copying.* `ReadWriteStream` is conceptually both a `ReadStream` and a `WriteStream`. However, Smalltalk is a single inheritance-based language, so `ReadWriteStream` can only subclass one of these. The behaviour from the other one has to be copied, leading to code duplication and all of its related maintenance problems.

The designers of the Squeak stream hierarchy decided to subclass `WriteStream` to implement `ReadWriteStream`, and then copy the methods related to reading from `ReadStream`.

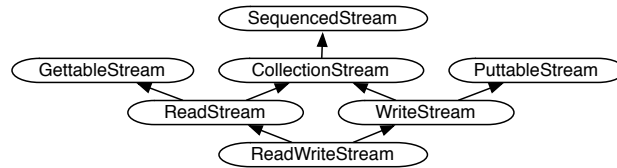
One of the copied methods is `next`, which reads and returns the next element in the stream. This leads to a strange situation where `next` is being canceled

<sup>3</sup> In Smalltalk, canceling a method is done by reimplementing the method in the subclass and calling `shouldNotImplement` from it.

out in `WriteStream` (because it should not be doing any reading), only to be reintroduced by `ReadWriteStream`. The reason for this particular situation is due to the combination of `next` defined too high in the hierarchy and single inheritance.

*Reimplementation.* In Figure 1, one can see that `next:` is implemented five times. Not a single implementation uses `super` which means that each class completely reimplements the method logic instead of specializing it. But this statement should be tempered because often in the Squeak stream hierarchy, methods override other methods to improve speed execution: this is because in subclasses, the methods have more knowledge and, thus, can do a faster job. However, a method reimplemented in nearly all of the classes in a hierarchy suggests inheritance hierarchy anomalies.

## 2.2 The ANSI Smalltalk standard



**Fig. 2.** The ANSI Smalltalk standard stream protocol hierarchy.

Figure 2 shows that even if Smalltalk is a single inheritance language, the ANSI Smalltalk standard [7] defines the different protocols using multiple inheritance. In the standard, streams are based on the notion of sequence values. Each stream has past and future sequence values. The ANSI Smalltalk standard defines a decomposition of stream behavior around three main protocols: `GettableStream`, `SequencedStream` and `PuttableStream`. Table 1 and Table 2 summarize the protocol contents.

The ANSI Smalltalk standard provides a useful starting point for an implementation even if a lot of useful methods are not described. We therefore chose to adopt it for Nile.

*About `GettableStream`>>`peekFor:`.* The standard proposes a definition of `peekFor:` that most Smalltalk implementations do not follow. The ANSI Smalltalk standard is equivalent to an equality test between the peeked object and the parameter:

```

GettableStream>>peekFor: anObject
  ^ self peek = anObject
  
```

SequencedStream	
close	Disassociate a stream from its backing store.
contents	Returns a collection containing the receiver's past and future sequence values in order.
isEmpty	Returns a boolean indicating whether there are any sequence values in the receiver.
position	Returns the number of sequence values in the receiver's past sequence values.
position:	Sets the number of sequence values in the receiver's past sequence values to be the parameter.
reset	Resets the position of the receiver to be at the beginning of the stream of values.
setToEnd	Set the position of the stream to its end.

PuttableStream	
flush	Upon return, if the receiver is a write-back stream, the state of the stream backing store must be consistent with the current state of the receiver.
nextPut:	Writes the argument to the stream.
nextPutAll:	Enumerate the argument, adding each element to the receiver.

**Table 1.** The `SequencedStream` and `PuttableStream` protocols defined by the ANSI Smalltalk standard.

GettableStream	
atEnd	Returns true if and only if the receiver has no future sequence values available for reading.
do:	Evaluates the argument with each receiver future sequence value.
next	The first object is removed from the receiver's future sequence values and appended to the end of the receiver's past sequence values. The object is returned.
next:	Does next a certain amount of time and returns a collection of the objects returned by next.
nextMatchFor:	Reads the next object from the stream and returns true if and only if the object is equivalent to the argument.
peek	Returns the next object in the receiver's future sequence values without advancing the receiver's position.
peekFor:	Peek at the next object in the stream and returns true if and only if it matches the argument.
skip:	Skip a given amount of object in the receiver's future sequence values.
skipTo:	Sets the stream just after the next occurrence of the argument and returns true if it's found before the end of the stream.
upTo:	Returns a collection of all the objects in the receiver up to, but not including the next occurrence of the argument.

**Table 2.** `GettableStream` protocol defined by the ANSI Smalltalk standard.

Most Smalltalk implementations (including Dolphin, GemStone, Squeak, VisualAge, VisualSmalltalk, VisualWorks, Smalltalk-X and GNU Smalltalk) do not only test the equality but also increment the position in case of equality as shown by the following implementation.

```
peekFor: anObject
    "Answer false and do not move over the next element if it is not equal
    to the argument, anObject, or if the receiver is at the end. Answer
    true and increment the position, if the next element is equal to
    anObject."

    ^ (self atEnd not and: [self peek = anObject])
      ifTrue: [self next. true]
      ifFalse: [false]
```

This definition lets the following code parse '145', ' 145' and '-145' without problem:

```

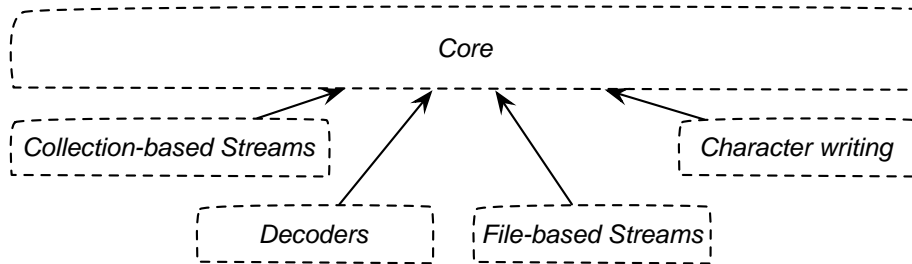
stream := ReadStream on: '- 145'.
negative := stream peekFor: '$-'.
stream peekFor: Character space.
number := stream upToEnd.

```

*Regarding the name of *SequencedStream*.* The name *SequencedStream* is not well chosen, since this protocol provides absolute positioning in the stream. A name evoking this would have been better.

### 3 Nile overview and core

Nile is designed around a core of traits offering base functionality reflecting the ANSI Smalltalk standard. The core consists of only three traits and it is then used in several libraries that we discuss in detail throughout the paper. File-based streams and collection-based streams are among the most prominent libraries. Other libraries we discuss are support for writing character constants and decoders (streams that can be chained). Figure 3 presents an overview of Nile.



**Fig. 3.** Overview of Nile: the core and its different libraries.

We designed Nile around three independent traits, reflecting the ANSI Smalltalk standard: *TPositionableStream*, *TGettableStream* and *TPuttableStream*. They are shown in Figure 4.

*TGettableStream.* The trait *TGettableStream* is meant for all streams used to read elements of any kind. The trait requires 4 methods: *atEnd*, *next*, *peek* and *outputCollectionClass*. The method *peek* returns the following element without moving the stream whereas *next* reads and returns the following element and moves the stream. The method *TGettableStream>>outputCollectionClass* is used to determine the type of collection which is used when returning collection of elements as with *next:* and *upTo:*.

<b>TGettableStream</b>		<b>TPositionableStream</b>		<b>TPutableStream</b>	
do:	atEnd	atEnd	position	nextPutAll:	nextPut:
nextMatchFor:	next	atStart	setPosition:	next:put:	
next:	peek	close	size	print:	
peekFor:	outputCollectionClass	isEmpty		flush	
skip:		position:			
skipTo:		reset			
upTo:		setToEnd			
upToEnd					
upToElementSatisfying:					

**Fig. 4.** The Nile core traits.

*TPositionableStream.* The trait `TPositionableStream` allows for the creation of streams that are positioned in absolute manner. It corresponds to the ANSI Smalltalk standard `SequencedStream` protocol; we thought the name `TPositionableStream` made more sense. The only required methods are `size` and two accessors for a `position` variable. We decided to implement the bound verification of the method `position:` in the trait itself: the parameter must be between zero and the stream size. This means that two methods have to be implemented: a pure accessor, named `setPosition:` here, and the real public accessor named `position:` which verifies its parameter value.

*TPutableStream.* This trait is the simplest of the Nile library. It provides `nextPutAll:`, `next:put:`, `print:` and `flush` and requires `nextPut:`. By default, `flush` does nothing. It is used for ensuring that everything has been written. Buffer-based streams should have their own implementations.

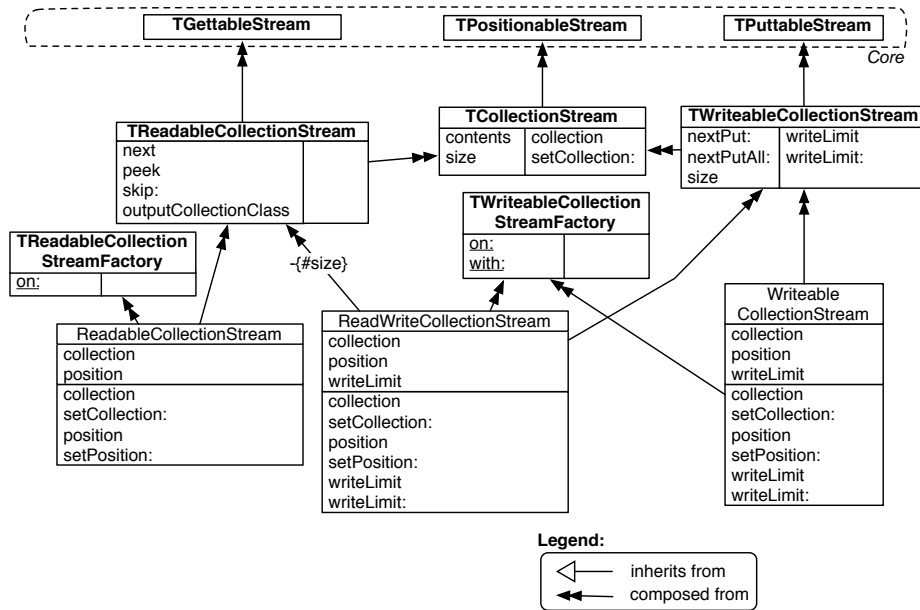
## 4 Collection-based streams

To support streaming over collections we implemented a set of dedicated traits and what we call *trait factories* that define their creation protocols. Note that, in contrast to the default Squeak implementation and like in VisualWorks, our implementation actually works with any sequenceable collection, not just `Arrays` and `Strings`.

### 4.1 The traits

The traits `TCollectionStream`, `TReadableCollectionStream` and `TWritableCollectionStream` implement the collection-based functionalities (as shown in Figure 5 — Note that in the figures traits have their name in bold whereas classes not). They provide all necessary methods required by the core traits, while only requiring 4 new accessors.





**Fig. 5.** The collection-based stream library. We use a UML-based notation to represent traits: methods on the right are *required* and methods on the left are *provided*.

*TReadableCollectionStream.* The trait `TReadableCollectionStream` helps creating classes which streams over readable collections. It implements the required methods of `TGettableStream`: `next`, `outputCollectionClass`, and `peek`. It also redefines `skip`: for efficiency reasons. The required method `TGettableStream>>atEnd` is provided by `TPositionableStream` and thus, does not require further work.

*TCollectionStream.* This trait is inspired by the ANSI Smalltalk standard. It is used for every stream that needs to read from or write to a collection. This trait defines `contents` and `size` in terms of two new methods: `collection` and `setCollection`. The former must return the internal collection and the latter provides a setter for this collection. The method `size` returns the size of the collection.

*TWritableCollectionStream.* The trait `TWritableCollectionStream` depends on a new instance variable accessible through two accessors `writeLimit` and `writeLimit`. This variable allows the internal collection to be bigger than the number of characters in the stream. This is a common technique used to avoid creation of a new collection each time an object is written to the stream. The `TWritableCollectionStream>>size` returns the value of `writeLimit` and `nextPut`: writes its parameter at the right position in the collection. The trait also reimplements `nextPutAll`: for efficiency reasons.

## 4.2 Trait factories

The ANSI Smalltalk standard defines `ReadStreamFactory>>on:` and `WriteStreamFactory>>with:` to create new streams. Basically there are three places where the stream instance creation methods can be defined. The most two natural ones are on the traits `TReadableCollectionStream` and `TWritableCollectionStream` or directly in the classes. Each solution has advantages and disadvantages. Adding the instance creation methods in the two traits helps their reuse. However, this forces all classes interested in these traits to have those same methods, even if they don't need them. If the instance creation methods are implemented in the classes, there will be duplication amongst the different classes.

We chose a third solution and implement the instance creation methods in separate traits. We named those traits “factories” because they support new stream creation.

We developed two factories: `TReadableCollectionStreamFactory` and `TWritableCollectionStreamFactory`. The former implements `on:` and the latter implements `on:` and `with:`. Even if the ANSI Smalltalk standard does not define `on:` for writable streams, we decided to implement it following the Squeak and VisualWorks implementations.

## 4.3 Classes

Traits alone are not enough to create a library. Classes are required to compose and create new instances. The original Squeak hierarchy provides three classes for collection-based streams: `ReadStream`, `WriteStream` and `ReadWriteStream`. Our implementation has equivalent classes with more explicit names: `ReadableCollectionStream`, `WritableCollectionStream` and `ReadWriteCollectionStream`.

Those classes have nothing more to do than declaring the use of already defined traits, declaring some instance variables and implementing the required accessors.

The only difficulty arises with `ReadWriteCollectionStream` which has a conflict with the method `size`. The method `size` is implemented in both `TReadableCollectionStream`, obtained from `TCollectionStream`, and `TWritableCollectionStream`. The first implementation reflects the size of the collection whereas the other takes care of the variable `writeLimit` and the efficient implementation in `TWritableCollectionStream`. That's why `ReadWriteCollectionStream` has to use the implementation of `TWritableCollectionStream`. To do this, the class removes the implementation of `size` coming from `TReadableCollectionStream`. This can be seen in Figure 5 on the arrow going from `ReadWriteCollectionStream` to `TReadableCollectionStream`<sup>4</sup>.

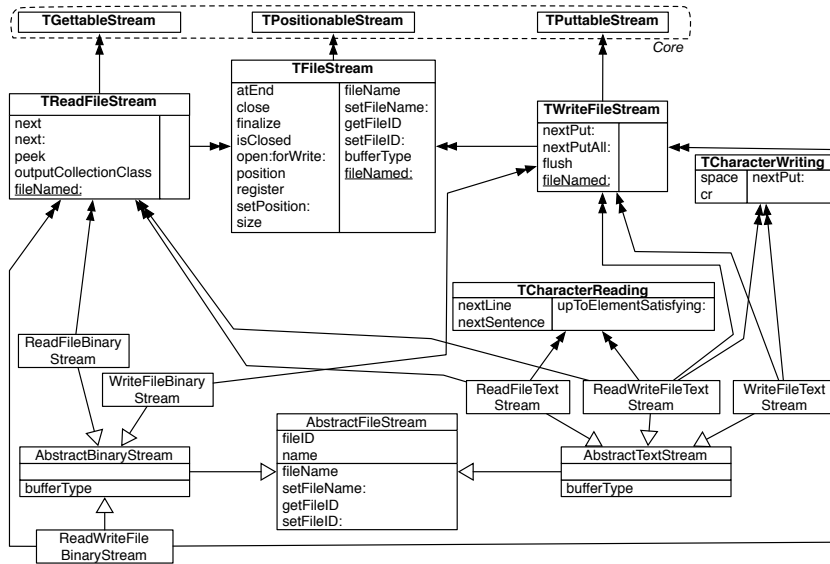


Fig. 6. FileStream implementation.

## 5 File-based streams

Nile includes a file-based stream library, shown in Figure 6. As with other file-based streams, it allows one to work with both binary and text files, supporting three access modes for each (read, write, and readwrite).

Each kind of file access is represented by a different class: the developer must explicitly choose the class based on what she wants to do with the file: reading, writing or both, in a binary or a text file. That way, the user has only the methods she can send in the interface of the stream. Note that this is a library design choice and it does not impact the way we decompose the behavior into traits.

Each file-based stream should be positionable, that’s why the trait `TPositionableStream` is used by `TFileStream`. `TFileStream` is the common trait for all file-based streams. It implements base functionalities for file access and requires four accessors, a `bufferType` method and an instance creation method `fileNamed`. The private method `bufferType` is used to differentiate binary from text files.

The traits `TReadFileStream` and `TWriteFileStream` use the reusable traits `TGettableStream` and `TPutableStream` from the core, respectively. They implement the required methods of these traits. Having implemented the reading and writing methods in separate traits instead of classes really helps here. This way, our

<sup>4</sup> The trait model gives the composer the possibility to remove methods through the minus (-) operator.

file-based streams only get the desired methods, not all methods like in the Squeak hierarchy.

At the very bottom of Figure 6, we defined the abstract classes `AbstractFileStream`, `AbstractBinaryStream` and `AbstractTextStream` to factorize instance variables definition and accessors. These abstract classes allow the definition of the six concrete classes with no more work.

Text-based streams use the traits `TCharacterReading` and `TCharacterWriting` depending on the type of file access. Even if simple, these two traits help defining methods only where they are needed and in all places where they are needed.

## 6 Other libraries

In this section we show how traits support reuse by presenting two libraries. We first present how character-related writing methods can be factored out in a trait. And then we describe the trait `TDecoder` that implements stream composition. Note that Nile offers several other libraries which are summarized later in Table 3.

### 6.1 Writing characters

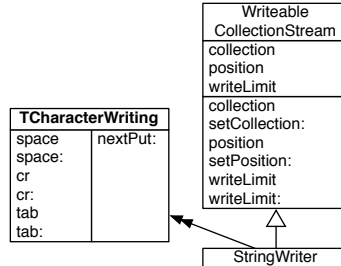


Fig. 7. Writing characters to a string.

In the Squeak hierarchy, the class `WriteStream` contains methods like `space`, `cr` and `tab` to write specific characters. These methods are only useful in case the user wants to write characters in her stream. If she wants to write binary data then those methods are useless and even pollute the interface of the stream. That's why we chose to implement the character-writing methods in a specific trait `TCharacterWriting`. Another advantage of using a specific trait is that Nile is then able to give those methods to any class which can write characters such as `StringWriter` in Figure 7 (a collection-based write stream which is writing characters) or `WriteFileTextStream` and `ReadWriteFileTextStream` in Figure 6.

## 6.2 Decoders

Developers often want to chain several streams. They want to use them like pipes that are connected together. For example, a developer may want a stream to read from a file and another stream which decompresses the first one on-the-fly. We generalized a mechanism which was already available in Squeak for classes like `ZipWriteStream` and have implemented a trait to support the composition of such decoders. We first present a scenario for such decoders and then describe our implementation.

A decoder is a `GettableStream` which reads its data from another `GettableStream` called its input stream. This way decoders can be chained. The decoder can do whatever it wants with the contents of its input stream: for example, it can ignore some elements, it can convert characters to numbers, it can compress or decompress...

*Selective number reading.* Imagine you have a string, or a file, containing space separated numbers. We can get all even numbers as presented in the code below. Here the developer composes three elementary streams which are subclasses of `Decoder` which uses the trait `TDecoder`.

```
| stream |
stream := ReadableCollectionStream on: '123 12 142 25'.
stream := NumberReader inputStream: stream.
stream := SelectStream selectBlock: [:each | each even] inputStream: stream.

stream peek.    ==> 12
stream next.    ==> 12
stream atEnd.   ==> false
stream next.    ==> 142
stream atEnd.   ==> true
```

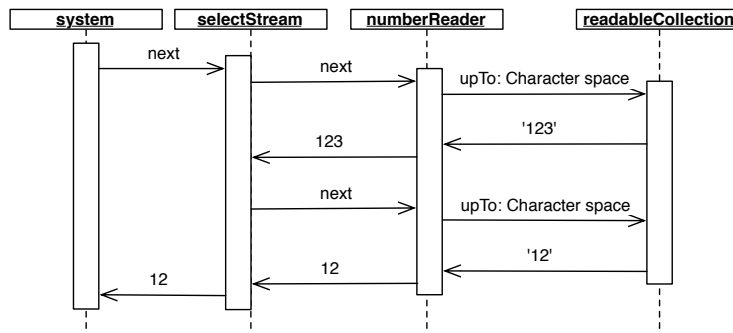
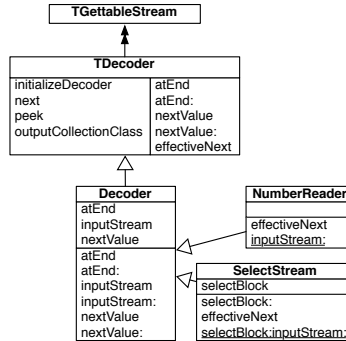


Fig. 8. Chaining streams

Figure 8 illustrates the stream connection. `NumberReader` transforms a character based stream in a number-based stream. `SelectStream` ignores all elements in the input stream for which the select block does not answer `true`.



**Fig. 9.** The decoder and two possible clients.

*The trait TDecoder.* Figure 9 shows the decoder hierarchy. A decoder is basically a `GettableStream`, that's why `TDecoder` uses the trait `TGettableStream`. We chose to implement the decoding methods in a trait to let developers incorporate its functionalities into their own hierarchies.

`TDecoder` provides implementations for all required methods of `TGettableStream` (see Figure 4) but `atEnd` and it requires four accessors (including `atEnd`) and the method `effectiveNext`. This method `effectiveNext` is where all the work happens. It should read its input stream and return a new element. The method `TDecoder>>next` calls `effectiveNext` and catches `StreamAtEndErrors` for setting the `atEnd` variable.

Factoring the Nile core in traits proved again to be useful. If we had implemented it using single inheritance we would have been forced to choose a superclass between class `Stream`, which provides writing methods we don't want, or class `ReadStream` which only streams over collections, which is not what we want to do with decoders.

## 7 Discussions

This section compares Nile with other stream implementations, analyzes its performance and discusses where traits did and did not help us.

## 7.1 Comparison with previous work

There is no previous work building a library from scratch using traits. However, Schärli *et al.* [5] were the first to refactor the collection and stream hierarchies using traits.

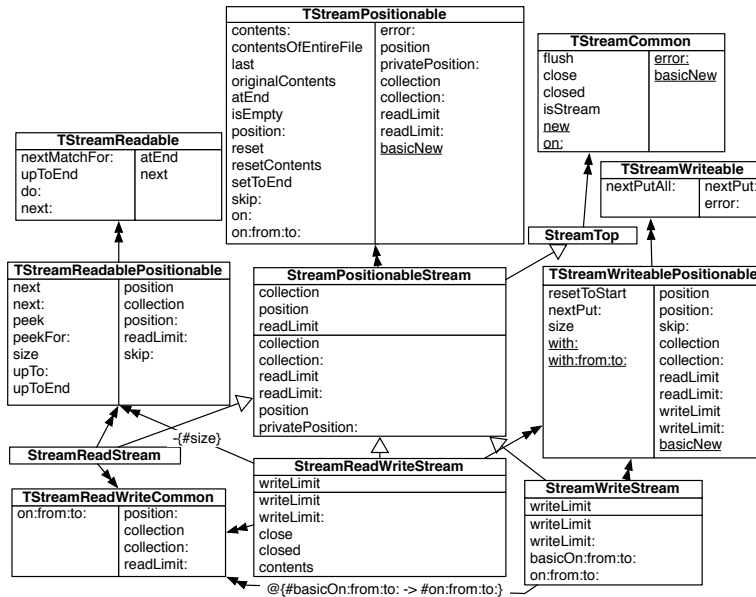


Fig. 10. Schärli's refactored stream hierarchy.

Figure 10 shows Schärli's stream hierarchy. Their work is a refactoring, where they took the original Squeak stream hierarchy and extracted the existing behavior into traits. This was a valuable experience that showed how a non-trivial implementation could be replaced with a cleaner implementation that was backwards compatible. While valuable, the backward compatibility constraint forces the result to be linked to the original implementation. Therefore it exhibits a number of problems:

- The positioning methods for a stream have to be based on collections because the methods `position:`, `atEnd` and `setToEnd` are all defined in the trait `TStreamPositionable` which depends on `collection` and `collection:`. Therefore it can not be used with files for example.
- The method `TStreamReadablePositionable>>peek` is dependent of the existence of methods `collection` and `position` but it shouldn't be.
- The granularity of the traits is big which hampers their reuse. For example, if we would like to have a `skip:` method, which is provided by `TStreamPositionable`,

we would get many more methods and, worse, we have to provide many collection-related methods.

## 7.2 Nile Analysis

*The factories.* Having implemented the factories in two separate traits complicates the hierarchy. Another solution would have been to define `ReadWriteCollectionStream` as a subclass of `WritableCollectionStream` to inherit both instance creation methods directly. However we believe that having explicit traits is better, since they are potentially reusable.

*Using an abstract superclass.* Nile defines three concrete classes to stream over collections: `ReadableCollectionStream`, `WritableCollectionStream` and `ReadWriteCollectionStream`. They all define the same instance variables and the same instance creation methods. To simplify the implementation of these classes, we could have implemented an abstract superclass for all of these classes with two common instance variables `position` and `collection` and their accessors. This is what we chose for the file-based streams (see Figure 11).

*Classes vs. Traits.* One of the key questions when building a system with traits is to decide when to use classes and when to use traits. In certain situations as illustrated by the Squeak stream hierarchy (see Section 2), defining a class or inheriting from a class does not make sense since some of its state is not used or its behavior should be canceled. This is a clear indication for using traits.

Most of the time however the decision is not that easy to take and the designer has to assess whether potential clients may benefit from the traits, *i.e.*, if the defined behavior can be reused in another hierarchy. In a lot of situations this means that traits are favored, since the price to pay to use traits is very low compared with the benefits one gets.

*Reuse at Work.* Figure 11 offers an overview of the core and some libraries of Nile. The fact that we based our implementation on traits rather than on inheritance and that we completely rethought the stream hierarchy leads to several advantages.

With Nile comes some really reusable traits which can be plugged in any other hierarchy. For example, implementing socket-based streams would only require socket manipulation work whereas utility methods like `nextPutAll:`, `skip:`, `upToEnd` are offered to the developer. Using the trait `TGettableStream`, a developer can easily implement a `Random` class which is basically a stream over random numbers. Table 3 presents the current clients we implemented in Nile using traits as well as the number of implemented methods to get the desired behavior.

Table 4 presents how much our core traits are reused. It presents for each traits the number of clients, the number of required methods and the number of methods that the trait provides. We see a good ratio provided/required for most of the traits. The ratio may still improve if additional behavior based on the core functionality is introduced.



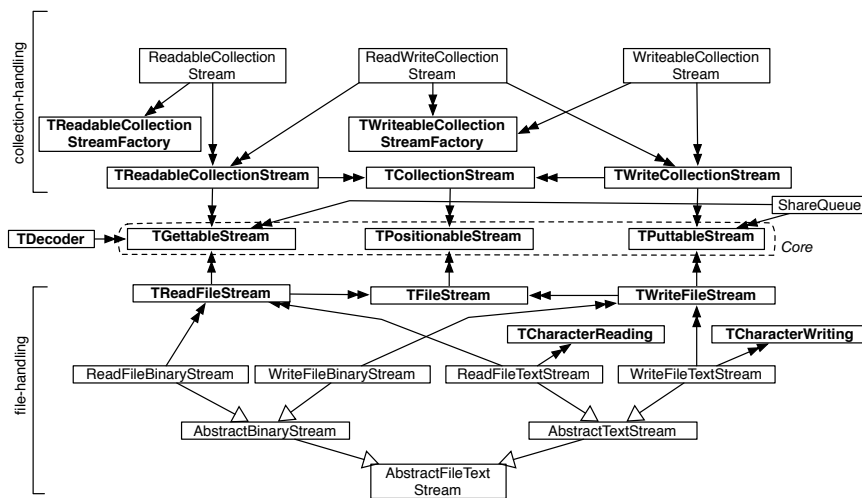


Fig. 11. An overview of Nile first clients

client name	superclass and trait used	met.	description
Random	TGettableStream	4	generate random numbers.
LinkedListStream	TGettableStream	5	stream over linked elements.
History	TPuttableStream	7	manage do and undo of command objects.
	TReadableCollectionStream		
SharedQueue	TWriteableCollectionStream	5	concurrent access on a queue.
	TGettableStream		
StringReader	TPuttableStream	0	add character-based reading methods.
	ReadableCollectionStream		
StringWriter	TCharacterReading	0	add character-based writing methods.
	WriteableCollectionStream		
CompositionStream	TCharacterWriting	1	multiplexer for input streams.
	Decoder		
Tee	Decoder	1	fork the input stream (like the Unix <i>tee</i> command).
Buffer	Decoder	1	add a buffer to any kind of input stream.
NumberReader	Decoder	1	read numbers from a character based input stream.
SelectStream	Decoder	1	select elements from an input stream.
PipeEntry	TGettableStream	7	allow data to be manually put into a pipe.
	TPuttableStream		

Table 3. Nile clients

Table 5 presents some metrics which compares the same functionalities in the Squeak implementation and in Nile for the collection-based streams. The first two metrics show that Nile uses a lot of traits and only a few classes. This is because Nile is designed to have fine grained and reusable units. The next two (number of methods and number of bytes) are more interesting and show that the amount of code is really smaller in Nile than in Squeak. Nile has 18% less methods and 15% less bytecodes than the corresponding Squeak collection-based stream library. Finally, we can deduce from the last metrics that the design of

Trait	client	classes	required	met.	provided	met.	$\frac{\text{provided}}{\text{required}}$
TGettableStream	22	4	11	275%			
TPositionableStream	20	3	9	300%			
TPuttableStream	13	1	4	400%			
TReadableCollectionStream	6	4	26	650%			
TCollectionStream	12	4	11	275%			
TWritableCollectionStream	6	6	23	383%			
TCharacterReading	3	2	1	50%			
TCharacterWriting	3	1	8	800%			
TByteReading	3	3	6	200%			
TByteWriting	3	2	5	250%			
TDecoder	7	6	14	233%			

**Table 4.** Nile-trait reusability.

	Squeak	Nile	$\frac{\text{Squeak}-\text{Nile}}{\text{Squeak}}$
Number of Classes And Traits	5	13	-160%
Number of Classes	5	4	20%
Number of Methods	53	43	18%
Number of Bytes	1725	1459	15%
Number of Cancelled Methods	2	0	100%
Number of Reimplemented Methods	14	3	78%
Number of Methods Implemented Too High	10	0	100%

**Table 5.** Some metrics for the collection-based streams

Nile is better: there is no cancelled method nor method implemented too high and there are only four methods reimplemented for speed reason compared to the fourteen of the Squeak version.

*About Trait Composition.* During trait composition it is possible that required methods of a trait are fulfilled by the provided methods of another traits. When this happens the developer does not have to do any extra work and benefits from the composition result. We can see this at work for the method `atEnd` that is required in `TGettableStream` and provided by `TPositionableStream`. The trait `TReadableCollectionStream` doesn't have any work to get the implementation of `atEnd`. However, such a situation is rare and based on the decomposition of traits using a compatible behavior and vocabulary.

However, it is sometimes better or necessary to override a method coming from a trait. It is because the new implementation have more knowledge than the overridden one and thus can do a better job. For example, the method `TReadableCollectionStream>>skip:` overrides the method `TGettableStream>>skip:`. The new method is more efficient because the stream is positioned directly, needing only a small bound computation:

```
TReadableCollectionStream>>skip: amount
"Moves relatively in the stream. Go forward amount elements. Go backward if amount is
negative. Overrides TGettableStream>>skip: for efficiency and backward possibility."
```

self position: ((self position + amount) min: self size max: 0)

Moreover, `skip:` is now able to go backward if `amount` is negative, which was not the case in the implementation of `TGettableStream`.

### 7.3 Performance optimization

One of the key challenge of Nile in terms of performance is to be able to iterate over any kind of collection while at the same time be as efficient as the squeak implementation for `Arrays` and `Strings`. We present our solution to this challenge.

Contrary to the Squeak class `WriteStream`, Nile's `TWriteableCollectionStream` is able to iterate over any kind of `SequenceableCollection`. In Squeak the method `WriteStream>>nextPutAll:` directly manipulates its internal collection using a primitive call to `replaceFrom:to:with:startingAt:` implemented in `String` and `Array`<sup>5</sup>, Nile has more work.

The idea is to propose a dedicated set of classes working specifically on `Array` and `String`. We first reimplemented the method `nextPutAll:` in `TWriteableCollectionStream` to take care of any kind of collection. This proved to be slow when iterating over `Arrays` and `Strings` compared to Squeak. Benchmarking shows that too much time was lost into calling methods. We have then implemented an optimized version (*i.e.*, using the primitive mentioned above) of `nextPutAll:` directly into the classes `ReadWriteArrayStream` and `WritableArrayStream` in which we are sure that the underlying collection is an `Array` (as shown on the left side of Figure 12).

*Accessor-use impact.* Traits cannot define state, which must be accessed via accessor methods. As Squeak does not have a JIT compiler, using accessors instead of direct instance variable access has a cost. Table 6 shows that using accessors in the context of stream on strings and arrays is 41% times slower than direct instance variable access. To optimize our library as much as possible we used direct accesses, *i.e.*, as shown on the left of Figure 12 we redefined `nextPutAll:`. However this has as impact that we have to duplicate the optimized implementation of `nextPutAll` into `WritableArrayStream` and `ReadWriteArrayStream`.

	execution (per second)	$\frac{nile}{squeak}$
Squeak implementation	126	
Nile with direct variable access	138	110%
Nile with accessors	81	64%

**Table 6.** Nile performances. Without accessors, Nile is faster than Squeak. But using them makes it slower.

<sup>5</sup> While `replaceFrom:to:with:startingAt:` is implemented for all kinds of `SequenceableCollections`, it does not work for `OrderedCollection`.

The right of Figure 12 presents another solution we implemented using an extra trait to share the optimized method for the two classes (*i.e.*, calling the primitive). However we discarded this solution since it is slower because traits forced us to use accessors.

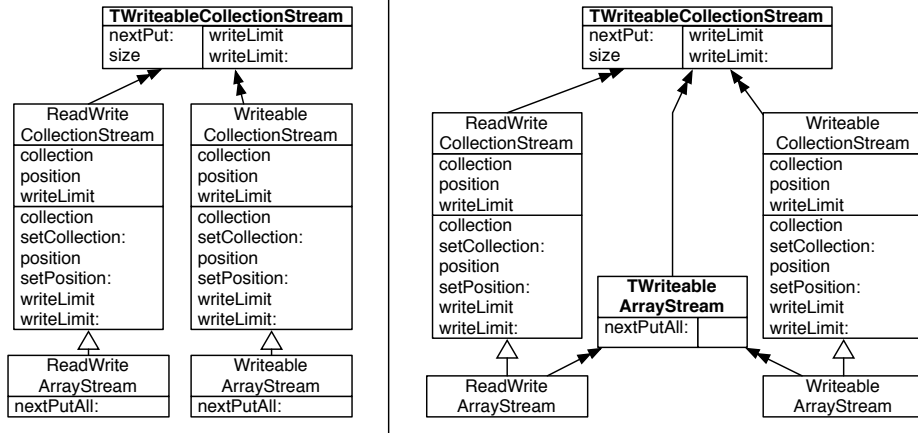


Fig. 12. Two solutions to optimize nextPutAll:

## 8 Problems with traits

Since one of the original goal of the work presented in this paper is to identify potential problems with traits, we now report the problems we faced while developing Nile. Note that some problems are not trait specific but due to Smalltalk's lack of visibility controls. In addition it should be noted that we did not encounter problem with aliasing in the context of recursive calls which is a known problem of traits.

### 8.1 Interface pollution

In this section we present some problems due to class interface extension.

*Required accessors.* With stateless traits, it is not possible to add state, *i.e.*, instance variables, to traits. Instead, the developer must add required accessors to its trait and the classes will implement those required accessors and the instance variable. This is a problem because the accessors are then part of the interface of the classes and this adds a burden to the class developers. However this would be solved if Smalltalk would have method access control. Stateful traits

[3] solve this problem by allowing traits to contain private state. For example, if we had used stateful traits the methods `TPositionableStream>>setPosition:` and `TWritableCollectionStream>>writeLimit` would not have been required.

However, developing Nile showed that stateful traits would not have been of great help. If we examine the trait `TCollectionStream` in Figure 5, we can see that implementing an instance variable `collection` here would have been interesting because classes would not have needed to define it. But, methods `collection` and `setCollection:` would still need to be in the interface because `TReadableCollectionStream>>outputCollectionClass` and `TReadableCollectionStreamFactory>>on:` need them.

We believe that stateful traits are not as interesting as what a first impression might tell.

*Lazily initialized variable.* There are basically three ways of initializing an instance variable giving it a first value: initializing lazily the variable in the accessor, using an `initialize` method, or initializing the variable in the instance creation method through an accessor.

Lazy initialization is a common programming pattern. Here is an example in Smalltalk which returns the value of the variable `checked` if it has been set or sets it to `false` and returns `false`:

```
checked
  ^ checked ifNil: [checked := false]
```

Now, imagine a trait needs a variable and a default value. Since traits can't contain state in their standard implementation, accessors must be required methods. But where do you lazily initialize the variable? Two solutions are possible: you can force users of the trait to initialize the variable or you can initialize in the trait and use another method for accessing the variable. Here is an example of the later possibility:

```
checked
  ^ self getChecked ifNil: [self checked: false. false].

checked: aBoolean
  self explicitRequirement.

getChecked
  self explicitRequirement.
```

This solution pollutes the trait interface with an unnecessary method `getChecked`. The other solution consists of letting the trait user initialize the variable. This solution does not pollute the interface but gives more responsibility to other developers and may produce code duplication or bugs.

The same problem appears when you want to do some checking before assigning to a variable as shown in `TPositionableStream>>position:` for example:

```

TPositionableStream>>position: newPosition
"Sets the number of elements before the position to be the parameter
newPosition. 0 for the start of the stream. Throws an error if the
parameter is lesser than 0 or greater than the size."

(self isInBounds: newPosition) iffFalse: [InvalidArgumentError signal].
self setPosition: newPosition.

```

This setter needs an additional method `setPosition:` which really modifies the variable and which is a required method of the trait.

*Initializing a trait.* In a class, when a developer wants to initialize a newly created object, he can use an `initialize` method:

```

initialize
  super initialize.
  color := Color transparent.

```

This can be done in a trait too provided that the developer uses an accessor instead of a direct reference to the variable. Problems arise when a class or a trait uses multiple traits, each defining its own `initialize` method. In this case, there will be conflicts and those conflict can only be resolved by aliasing. This brings lots of pollution in the class interface and require a lot of work.

Another solution would be to use a specific name for each method `initialize`. For example, if the trait `TPositionableStream` needs an `initialize` method, the developer can name it `initializePositionableStream`. Each user of the trait now needs to define its own `initialize` method which calls `initializePositionableStream`. We believe this is still clunky and requires too much work from the developer.

*Initializing in the instance creation method.* Instance creation methods can be used to initialize variables. This is what we did for Nile:

```

TWritableCollectionStreamFactory>>on: aCollection
  ^ self basicNew
    initialize;
    setCollection: aCollection;
    writeLimit: 0;
    reset;
    yourself

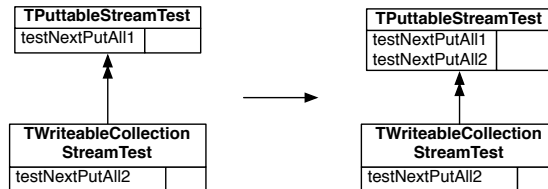
```

Smalltalk is made such that this requires that setters are available in the interface of the class. It also puts more responsibility on the instance creation method which now needs more knowledge over the class it instantiates.

## 8.2 Methods silently ignored

Sometimes, modifying a trait does not modify the users of this trait in the same way because of name overriding. Note that this problem is not trait specific but it is a problem of object-oriented programming as shown by Figure 13.

Figure 13 shows a part of our test hierarchy for Nile. The test hierarchy is very similar to the Nile hierarchy: for each model trait or class, there is a test trait or class. The method `nextPutAll:` is tested in two different places: in the methods `TPutableStreamTest>>testNextPutAll1` and `TWritableCollectionStreamTest>>testNextPutAll2`. If a tester adds a new test named `testNextPutAll2` in the trait `TPutableStreamTest`, then the test is silently ignored and will never be launched.



**Fig. 13.** If the tester implements `TPutableStreamTest>>testNextPutAll2`, the test will never be launched because `TWritableCollectionStreamTest>>testNextPutAll2` hides it.

## 9 Related work

We already compared our approach with the few work refactoring existing code using traits [5, 6]. We now present the approaches that automatically transform existing libraries using Formal Concept Analysis (FCA) or other techniques. FCA was used in different ways.

Godin [9] developed incremental FCA algorithms to infer implementation and interface hierarchies guaranteed to have no redundancy. To assess their solutions from a point of view of complexity and maintainability they propose a set of structural metrics. They analyze the Smalltalk Collection hierarchy. One important limitation is that they consider each method declaration as a different method and thus cannot identify code duplication. Moreover their approach serves rather as a help for program understanding than reengineering since the resulting hierarchies cannot be implemented in Smalltalk because of single inheritance.

In C, Snelting and Tip analyze a class hierarchy making the relationship between class members and variables explicit [10]. By analyzing the *usage* of the hierarchy by a set of client programs they are able to detect design anomalies such as class members that are redundant or that can be moved into a derived class. Taking into account a set of client programs, Streckenbach infer improved hierarchies in Java with FCA [11]. Their proposed refactoring can then be used for further manual refactoring. The tool proposes the reengineer to move methods up in the hierarchy to work around multiple inheritance situations generated by the generated lattice. The work of Streckenbach is based on the analysis of the

usage of the hierarchy by client programs. The resulting refactoring is behavior preserving (only) with respect to the analyzed client programs.

Lienhard *et al.* applied Formal Concept Analysis to semi-automatically identify traits [6]. We cannot really compare their resulting hierarchy with ours since the information about the respective traits is no longer available. However we can conclude that the resulting hierarchy was limited and resulted only from a refactoring effort and not from a new design.

Interfaces and specifications of the Smalltalk collection hierarchy are also analyzed by Cook [12]. He also takes method cancellation into account to detect protocols. By manual analysis and development of specifications of the Smalltalk collection hierarchy he proposes a better protocol hierarchy. Protocol hierarchies explicitly represent similarities between classes based on their provided methods. Thus, compared to our approach, protocol hierarchies present a *client* view of the library rather than one of the *implementor*.

Moore [13] proposes automatic refactoring of Self inheritance hierarchies. Moore focuses on factoring out common expressions in methods. In the resulting hierarchies none of the methods and none of the expressions that can be factored out are duplicated. Moore's factoring creates methods with meaningless names which is a problem if the code should be read. The approach is more optimizing method reuse than creating coherent composable groups of methods. Moore's analysis finds some of the same problems with inheritance that we have described in this paper, and also notes that sometimes it is necessary to manually move a method higher in the hierarchy to obtain maximal reuse.

Casais uses an automatic structuring algorithm to reorganize Eiffel class hierarchies using decomposition and factorization [14]. In his approach, he increases the number of classes in the new refactored class hierarchy. Dicky *et al.* propose a new algorithm to insert classes into a hierarchy that takes into account overridden and overloaded methods [15].

The key difference from our results is that all the work on hierarchy reorganization focuses on transforming hierarchies using inheritance as the only tool. In contrast, we are interested in exploring other mechanisms, such as explicit composition mechanisms like traits composition in the context of mixin-like languages. Another important difference is that we do rely on algorithms. This is important since we want to be able to use our result to compare it with the result of future approach extracting traits automatically, so the Nile library may serve as a reference point.

## 10 Conclusion

Traits are units of reuse that can be used to compose classes. This paper is an experience report. Even if other experiences have been made to test traits, they were always refactoring an existing hierarchy, moving methods from classes to traits. Our work however presents a brand new implementation. We started from the textual description from the ANSI Smalltalk standard and from existing implementations of stream libraries in Squeak and VisualWorks. Our result is a



completely new implementation, named Nile, of the stream hierarchy which does not share any code with previous implementations.

Our experience shows that traits are good building blocks which favor reuse across different hierarchies. In the present implementation of Nile we get up to 15% less code than the corresponding Squeak code. Core traits are reused by numerous clients. We also presented the problems we faced during the experience and believe that Nile can be used in the future as a reference point for comparing future trait enhancement.

This experience shows that well defined traits can naturally fit into lots of different clients which can benefit from methods offered by the trait for a relatively low cost.

### Acknowledgment

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