

11. Program Transformation

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Roadmap



- > Program Transformation
- > Refactoring
- > Aspect-Oriented Programming

Links

- > **Program Transformation:**
 - <http://swerl.tudelft.nl/bin/view/Pt>
 - <http://www.program-transformation.org/>
- > **Stratego:**
 - <http://strategoxt.org/>
- > **TXL:**
 - <http://www.txl.ca/>
- > **Refactoring:**
 - <http://www.ibm.com/developerworks/library/os-ecref/>
 - <http://recoder.sourceforge.net/wiki/>
 - <http://www.refactory.com/RefactoringBrowser/>
- > **AOP:**
 - <http://www.eclipse.org/aspectj/>

Roadmap



- > **Program Transformation**
 - Introduction
 - Stratego/XT
 - TXL
- > Refactoring
- > Aspect-Oriented Programming

Thanks to Eelco Visser and Martin Bravenboer for their kind permission to reuse and adapt selected material from their Program Transformation course.
<http://sweri.tudelft.nl/bin/view/Pt>

What is “program transformation”?

- > *Program Transformation* is the process of transforming one program to another.
- > Near synonyms:
 - Metaprogramming
 - Generative programming
 - Program synthesis
 - Program refinement
 - Program calculation

Applications of program transformation

> Translation

- *Migration*
- *Synthesis*
 - Refinement
 - Compilation
- *Reverse Engineering*
 - Decompilation
 - Architecture Extraction
 - Visualization
- *Program Analysis*
 - Control flow
 - Data flow

Refinement — transform high-level spec down to an implementation that fulfils requirements

Renovation — reengineering

Translation – compilation

<pre>function fact(n : int) : int = if n < 1 then 1 else (n * fact(n - 1))</pre>	<i>Tiger</i> ⇒	<pre>fact:subu \$sp, \$sp, 20 sw \$fp, 8(\$sp) addiu \$fp, \$sp, 20 sw \$s2, -8(\$fp) sw \$ra, -4(\$fp) sw \$a0, 0(\$fp) move \$s2, \$a1 li \$t0, 1 bge \$s2, \$t0, c_0 li \$v0, 1 b d_0 c_0: lw \$a0, (\$fp) li \$t0, 1 subu \$a1, \$s2, \$t0 jal fact_a_0 mul \$v0, \$s2, \$v0 d_0: lw \$s2, -8(\$fp) lw \$ra, -4(\$fp) lw \$fp, 8(\$sp) addiu \$sp, \$sp, 20 jr \$ra</pre>
---	-------------------	---

MIPS

<http://www.cs.uu.nl/docs/vakken/pt/slides/PT05-ProgramTransformation.pdf>

Translation – migration from procedural to OO

```
type tree = {key: int, children: treelist}
type treelist = {hd: tree, tl: treelist}
function treeSize(t : tree) : int =
  if t = nil then 0 else 1 + listSize(t.children)
function listSize(ts : treelist) =
  if ts = nil then 0 else listSize(t.tl)
```

Tiger



```
class Tree {
  Int key;
  TreeList children;
  public Int size() {
    return 1 + children.size
  }
}
class TreeList { ... }
```

Java

<http://www.cs.uu.nl/docs/vakken/pt/slides/PT05-ProgramTransformation.pdf>

Rephrasing – desugaring regular expressions

```
Exp := Id
     | Id "(" {Exp ","}* ")"
     | Exp "+" Exp
     | ...
```

EBNF

⇒

```
Exp := Id
     | Id "(" Exps ")"
     | Exp "+" Exp
     | ...
Exps :=
     | Expp
Expp := Exp
     | Expp "," Exp
```

BNF

Rephrasing – partial evaluation

```
function power(x : int, n : int) : int =  
  if n = 0 then 1  
  else if even(n) then square(power(x, n/2))  
  else (x * power(x, n - 1))
```

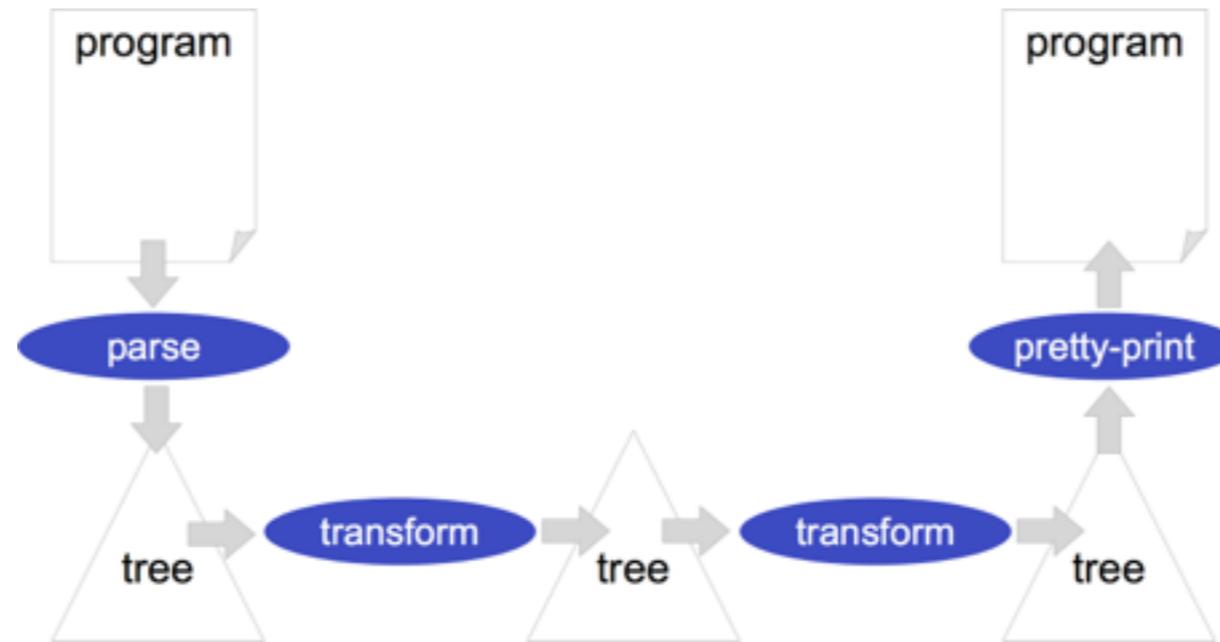
Tiger

↓ n = 5

Tiger

```
function power5(x : int) : int =  
  x * square(square(x))
```

Transformation pipeline



<http://losser.st-lab.cs.uu.nl/~mbravenb/PT05-Infrastructure.pdf>

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This general scheme applies to Stratego, TXL and various other systems. Transformation systems and languages may support or automate different parts of this pipeline.

If the source language is fixed, then a fixed parser and pretty-printer may be used.

If the source and target languages are arbitrary, then there should be support to specify grammars and automatically generate parsers and pretty-printers.

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Stratego/XT

> **Stratego**

—A language for specifying program transformations

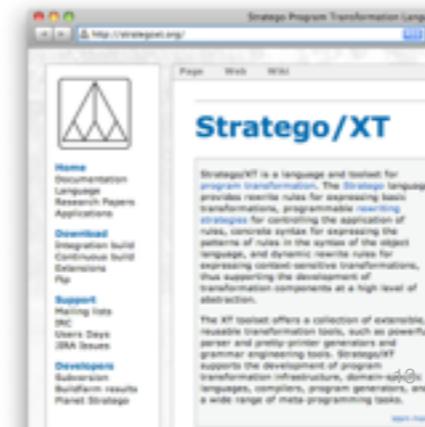
- *term rewriting rules*
- *programmable rewriting strategies*
- *pattern-matching against syntax of object language*
- *context-sensitive transformations*

> **XT**

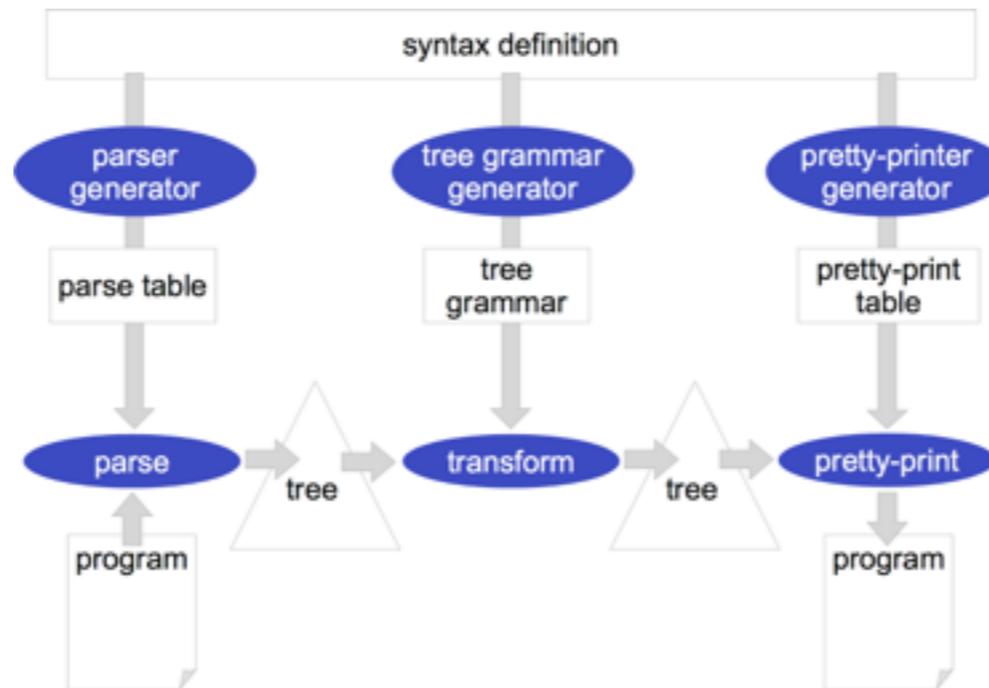
—A collection of transformation tools

- *parser and pretty printer generators*
- *grammar engineering tools*

<http://strategoxt.org/>



Stratego/XT



<http://losser.st-lab.cs.uu.nl/~mbravenb/PT05-Infrastructure.pdf>

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Parser and basic pretty-printer 100% generated.
Language specific support for transformations generated.

Parsing

Rules translate
terms to terms

*Stratego parses any
context-free language
using Scannerless
Generalized LR Parsing*

```
module Exp
exports
  context-free start-symbols Exp
  sorts Id IntConst Exp

lexical syntax
[\ \t\n] -> LAYOUT
[a-zA-Z]+ -> Id
[0-9]+ -> IntConst

context-free syntax
Id -> Exp {cons("Var")}
IntConst -> Exp {cons("Int")}

("(" Exp ")") -> Exp {bracket}

Exp "*" Exp -> Exp {left, cons("Mul")}
Exp "/" Exp -> Exp {left, cons("Div")}
Exp "%" Exp -> Exp {left, cons("Mod")}

Exp "+" Exp -> Exp {left, cons("Plus")}
Exp "-" Exp -> Exp {left, cons("Minus")}

context-free priorities
{left:
  Exp "*" Exp -> Exp
  Exp "/" Exp -> Exp
  Exp "%" Exp -> Exp
}
> {left:
  Exp "+" Exp -> Exp
  Exp "-" Exp -> Exp
}
```

File: Exp.sdf

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See the Makefile for the steps needed to run this.

GLR parsing essentially does a parallel, breadth-first LR parse to handle ambiguity.

http://en.wikipedia.org/wiki/GLR_parser

Testing

```
testsuite Exp
```

```
topsort Exp
```

File: Exp.testsuite

```
test eg1 parse
```

```
"1 + 2 * (3 + 4) * 3 - 1"
```

```
->
```

```
Minus(
```

```
  Plus(
```

```
    Int("1")
```

```
  , Mul(
```

```
    Mul(Int("2"), Plus(Int("3"), Int("4")))
```

```
  , Int("3")
```

```
  )
```

```
)
```

```
, Int("1")
```

```
)
```

Running tests

```
pack-sdf -i Exp.sdf -o Exp.def
including ./Exp.sdf
```

Pack the definitions

```
sdf2table -i Exp.def -o Exp.tbl -m Exp
SdfChecker:error: Main module not defined
--- Main
```

Generate the parse table

```
parse-unit -i Exp.testsuite -p Exp.tbl
```

Run the tests

```
-----
executing testsuite Exp with 1 tests
-----
```

```
* OK   : test 1 (egl parse)
-----
```

```
results testsuite Exp
successes : 1
failures  : 0
-----
```

Interpretation example

```
module ExpEval
  imports libstratego-lib
  imports Exp

  rules
    convert : Int(x) -> <string-to-int>(x)
    eval : Plus(m,n) -> <add>(m,n)
    eval : Minus(m,n) -> <subt>(m,n)
    eval : Mul(m,n) -> <mul>(m,n)
    eval : Div(m,n) -> <div>(m,n)
    eval : Mod(m,n) -> <mod>(m,n)

  strategies
    main = io-wrap(innermost(convert <+ eval))
```

File: ExpEval.str

File: ultimate-question.txt

```
1 + 2 * (3 + 4) * 3 - 1
```

Stratego separates the specification of *rules* (transformations) from *strategies* (traversals). In principle, both are reusable.

Strategies

A *strategy* determines how a set of rewrite rules will be used to traverse and transform a term.

- innermost
- top down
- bottom up
- repeat
- ...

Running the transformation

```
sdf2rtg -i Exp.def -o Exp.rtg -m Exp
SdfChecker:error: Main module not defined
--- Main

rtg2sig -i Exp.rtg -o Exp.str
Generate signature

strc -i ExpEval.str -la stratego-lib
[ strc | info ] Compiling 'ExpEval.str'
[ strc | info ] Front-end succeeded      : [user/system] = [0.56s/0.05s]
[ strc | info ] Optimization succeeded -O 2 : [user/system] = [0.00s/0.00s]
[ strc | info ] Back-end succeeded       : [user/system] = [0.16s/0.01s]
gcc -I /usr/local/strategoxt/include -I /usr/local/strategoxt/include -I /usr/local/strategoxt/
include -Wall -Wno-unused-label -Wno-unused-variable -Wno-unused-function -Wno-unused-parameter -
DSIZEOF_VOID_P=4 -DSIZEOF_LONG=4 -DSIZEOF_INT=4 -c ExpEval.c -fno-common -DPIC -o .libs/ExpEval.o
gcc -I /usr/local/strategoxt/include -I /usr/local/strategoxt/include -I /usr/local/strategoxt/
include -Wall -Wno-unused-label -Wno-unused-variable -Wno-unused-function -Wno-unused-parameter -
DSIZEOF_VOID_P=4 -DSIZEOF_LONG=4 -DSIZEOF_INT=4 -c ExpEval.c -o ExpEval.o >/dev/null 2>&1
gcc .libs/ExpEval.o -o ExpEval -bind_at_load -L/usr/local/strategoxt/lib /usr/local/strategoxt/lib/
libstratego-lib.dylib /usr/local/strategoxt/lib/libstratego-lib-native.dylib /usr/local/strategoxt/
lib/libstratego-runtime.dylib -lm /usr/local/strategoxt/lib/libATerm.dylib
[ strc | info ] C compilation succeeded : [user/system] = [0.31s/0.36s]
[ strc | info ] Compilation succeeded   : [user/system] = [1.03s/0.42s]

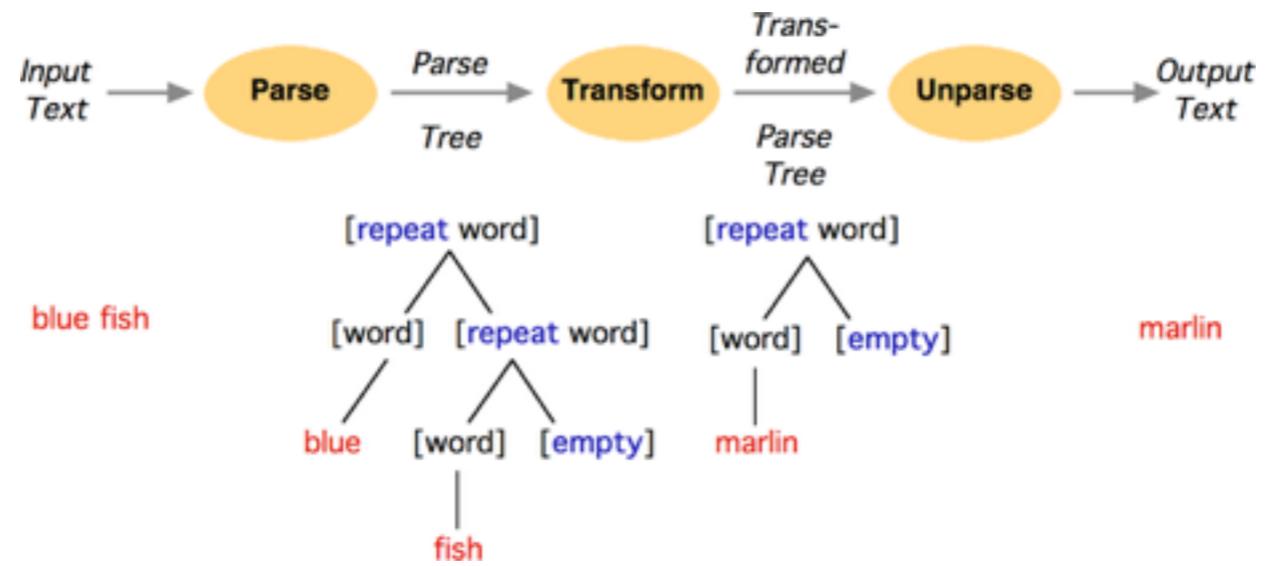
sglri -p Exp.tbl -i ultimate-question.txt | ./ExpEval
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Parse and transform
```

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The TXL paradigm: *parse, transform, unparse*



<http://www.txl.ca/docs/TXLintro.pdf>

TXL programs

Base grammar

defines tokens and non-terminals

Grammar
overrides

extend and modify types from grammar

Transformation
rules

rooted set of rules and functions

Expression example

File: Question.Txl

```
% Part I. Syntax specification
define program
  [expression]
end define

define expression
  [expression] + [term]
  | [expression] - [term]
  | [term]
end define

define term
  [term] * [primary]
  | [term] / [primary]
  | [primary]
end define

define primary
  [number]
  | ( [expression] )
end define
```

```
% Part 2. Transformation rules
rule main
  replace [expression]
    E [expression]
  construct NewE [expression]
    E [resolveAddition]
      [resolveSubtraction]
      [resolveMultiplication]
      [resolveDivision]
      [resolveBracketedExpressions]
  where not
    NewE [= E]
  by
    NewE
end rule

rule resolveAddition
  replace [expression]
    N1 [number] + N2 [number]
  by
    N1 [+ N2]
end rule
...

rule resolveBracketedExpressions
  replace [primary]
    ( N [number] )
  by
    N
end rule
```

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NB: TXL reverses the usual BNF convention and puts non-terminals in square brackets while interpreting everything else (except special chars) as terminals.

The default lexical scanner can be modified, but is usually fine for first experiments.

Running the example

File: Ultimate.Question

```
1 + 2 * (3 + 4) * 3 - 1
```

```
txl Ultimate.Question  
TXL v10.5d (1.7.08) (c)1988-2008 Queen's University at Kingston  
Compiling Question.Txl ...  
Parsing Ultimate.Question ...  
Transforming ...  
42
```

Example: TIL — a tiny imperative language

```
// Find all factors of a given input number
var n;
write "Input n please";
read n;
write "The factors of n are";
var f;
f := 2;
while n != 1 do
  while (n / f) * f = n do
    write f;
    n := n / f;
  end
  f := f + 1;
end
```

File: factors.til

<http://www.program-transformation.org/Sts/TILChairmarks>

TIL Grammar

```
% Keywords of TIL
keys
  var if then else while
  do for read write
end keys

% Compound tokens
compounds
  := !=
end compounds

% Commenting convention
comments
  //
end comments
```

File: TIL.Grm

All TXL parsers are also pretty-printers if the grammar includes formatting cues

```
define program
  [statement*]
end define

define statement
  [declaration]
  | [assignment_statement]
  | [if_statement]
  | [while_statement]
  | [for_statement]
  | [read_statement]
  | [write_statement]
end define

% Untyped variables
define declaration
  'var [id] ;           [NL]
end define

define assignment_statement
  [id] := [expression] ;   [NL]
end define

define if_statement
  'if [expression] 'then   [IN][NL]
  [statement*]             [EX]
  [opt else_statement]
  'end                     [NL]
end define
...
```

Pretty-printing TIL

```
include "TIL.Grm"  
function main  
  match [program]  
    _ [program]  
end function
```

File: TILparser.Txtl

```
txl factors.til TILparser.Txtl
```

```
var n;  
write "Input n please";  
read n;  
write "The factors of n are";  
var f;  
f := 2;  
while n != 1 do  
  while (n / f) * f = n do  
    write f;  
    n := n / f;  
  end  
  f := f + 1;  
end
```

Generating statistics

```
include "TIL.Grm"

function main
  replace [program]
    Program [program]

  % Count each kind of statement we're interested in
  % by extracting all of each kind from the program

  construct Statements [statement*]
    _ [^ Program]
  construct StatementCount [number]
    _ [length Statements] [putp "Total: %"]

  construct Declarations [declaration*]
    _ [^ Program]
  construct DeclarationsCount [number]
    _ [length Declarations] [putp "Declarations: %"]
  ...
  by
    % nothing
end function
```

File: TILstats.Txl

```
Total: 11
Declarations: 2
Assignments: 3
Ifs: 0
Whiles: 2
Fors: 0
Reads: 1
Writes: 3
```

Tracing

```
include "TIL.Grm"
...
redefine statement
  ...
  | [traced_statement]
end redefine

define traced_statement
  [statement] [attr 'TRACED]
end define

rule main
replace [repeat statement]
  S [statement]
  Rest [repeat statement]
...
  by
    'write QuotedS;      'TRACED
    S                    'TRACED
    Rest
end rule
...

```

File: TILtrace.Txl

```
write "Trace: var n;";
var n;
write "Trace: write \"Input n please\";";
write "Input n please";
write "Trace: read n;";
read n;
...

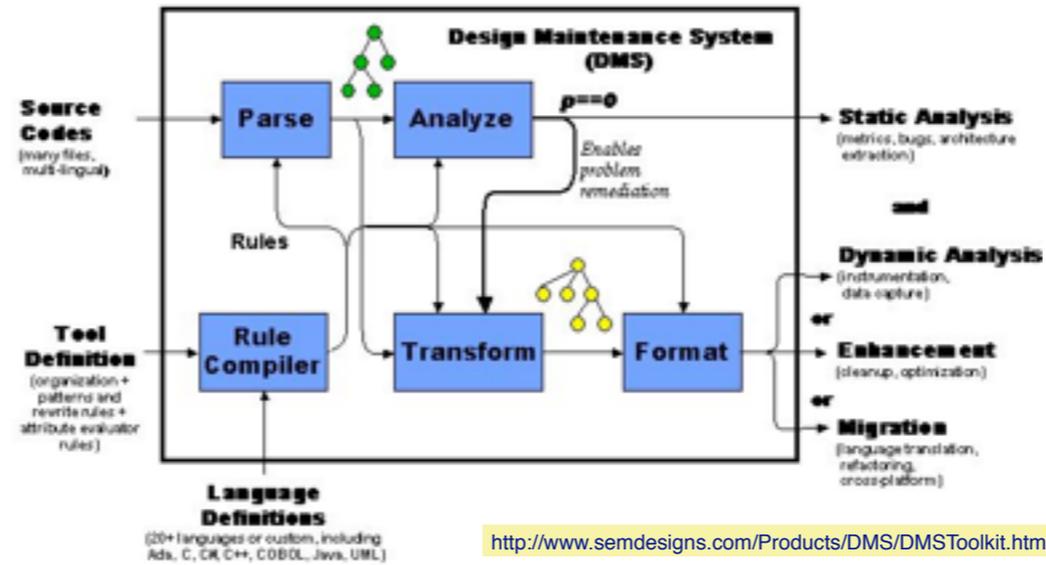
```

TXL vs Stratego

<i>Stratego</i>	<i>TXL</i>
Scannerless GLR parsing	Agile parsing (top-down + bottom-up)
Reusable, generic traversal strategies	Fixed traversals
Separates rewrite rules from traversal strategies	Traversals part of rewrite rules
	

Commercial systems

“The DMS Software Reengineering Toolkit is a set of tools for automating customized source program analysis, modification or translation or generation of software systems, containing arbitrary mixtures of languages.”



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See also <http://www.semdesigns.com/Products/DMS/DMSComparison.html> for a comparison to other approaches

Roadmap



- > Program Transformation
- > **Refactoring**
 - Refactoring Engine and Code Critics
 - Eclipse refactoring plugins
- > Aspect-Oriented Programming

What is Refactoring?

- > The process of *changing a software system* in such a way that it *does not alter the external behaviour* of the code, yet *improves its internal structure*.

— Fowler, et al., Refactoring, 1999.

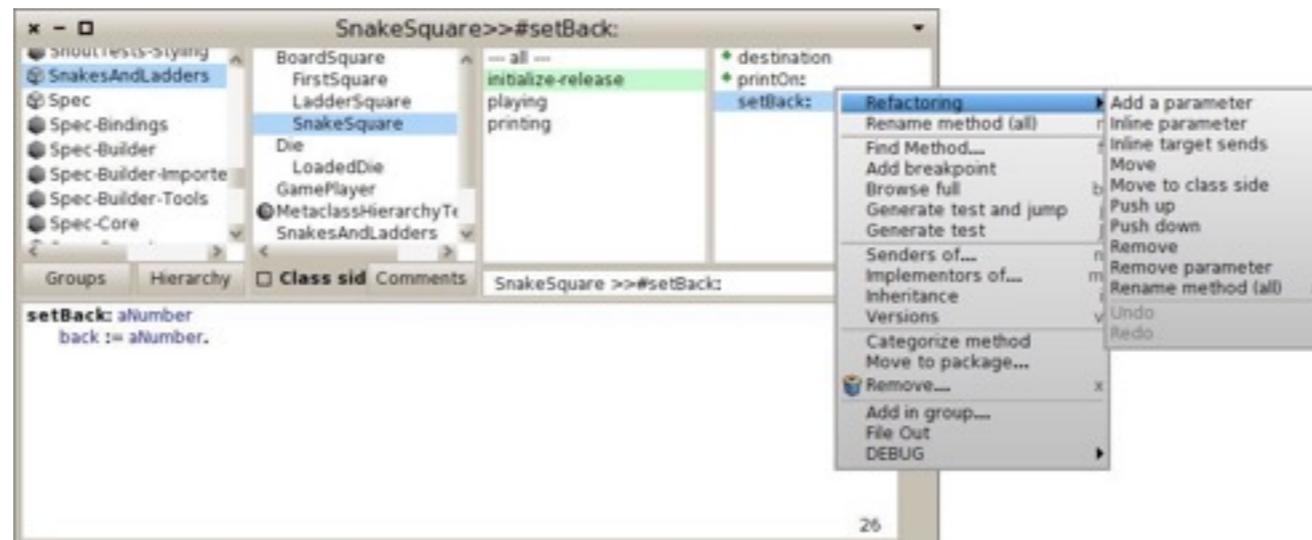
Rename Method — manual steps

- > Do it yourself approach:
 - Check that no method with the new name already exists in any subclass or superclass.
 - Browse all the implementers (method definitions)
 - Browse all the senders (method invocations)
 - Edit and rename all implementers
 - Edit and rename all senders
 - Remove all implementers
 - Test
- > Automated refactoring is better !

Rename Method

- > Rename Method (method, new name)
- > Preconditions
 - No method with the new name already exists in any subclass or superclass.
 - No methods with same signature as method outside the inheritance hierarchy of method
- > PostConditions
 - method has new name
 - relevant methods in the inheritance hierarchy have new name
 - invocations of changed method are updated to new name
- > Other Considerations
 - Typed/Dynamically Typed Languages => Scope of the renaming

The Refactoring Browser



Typical Refactorings

Class Refactorings	Method Refactorings	Attribute Refactorings
add (sub)class to hierarchy	add method to class	add variable to class
rename class	rename method	rename variable
remove class	remove method	remove variable
	push method down	push variable down
	push method up	pull variable up
	add parameter to method	create accessors
	move method to component	abstract variable
	extract code in new method	

*Bill Opdyke, "Refactoring Object-Oriented Frameworks,"
Ph.D. thesis, University of Illinois, 1992.*

*Don Roberts, "Practical Analysis for Refactoring,"
Ph.D. thesis, University of Illinois, 1999.*

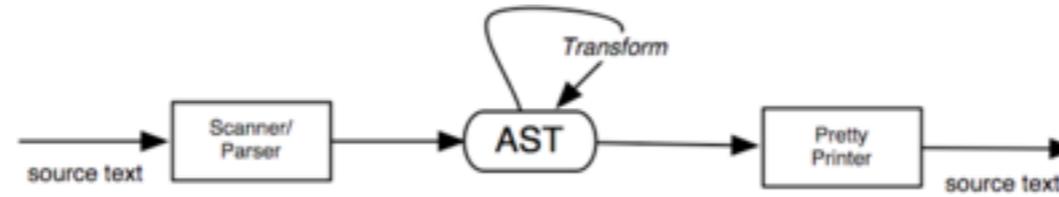
Code Critic – search for common errors

The screenshot shows the Code Critic application interface. The top window, titled "Critics Browser", displays a tree view of rules. The "Selected Rules" section is expanded to show a list of rules, including "Coding Idiom Violation" and "Eliminate guarding clauses". The "Eliminate guarding clauses" rule is selected, and its details are shown in the right pane. The details pane shows a list of rules, with "LinkedList>>#do: (Collections-Sequenceable)" selected. Below the panes are buttons for "Run new configuration", "Save Critics", "Browse", "Transform", and "False Positive".

```
do: aBlock

| aLink |
aLink := firstLink.
[ aLink isNil ]
whileFalse: [
  aBlock value: aLink value.
  aLink := aLink nextLink ]
[aLink == nil] whileFalse:
[aBlock value: aLink value.
aLink := aLink nextLink]
```

Refactoring Engine – matching trees



NB: All metavariables start with `

Syntax	Type
`	recurse
@	list
.	statement
#	literal

``@object halt	recursively match send of halt
`.Statements	match list of statements
Class `@message: `@args	match all sends to Class

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The first ` is for all meta-variables.

Rewrite rules

The screenshot shows the IDE interface for the `RBEqualNilRule` class. The top-left pane displays a project tree with the following structure:

- Refactoring-Core-Conditions
- Refactoring-Core-Model
- Refactoring-Core-Refactorings
- Refactoring-Core-Support
- Refactoring-Critics
 - Refactoring-Critics-BlockRules
 - Refactoring-Critics-ParseTreeR
 - Refactoring-Critics-Transforma
 - Refactoring-Critics-Unused
- Refactoring-Environment
- Refactoring-Pharo-Platform

The top-right pane shows the class hierarchy for `RBEqualNilRule`, listing various subclasses such as `RBAIfAnyNoneSatisfyRule`, `RBAIfAssignmentNilTrueRule`, `RBAIfAbsentRule`, `RBBetweenAndRule`, `RBCascadedNextPutAllsRule`, `RBDetectIfNoneRule`, `RBEqualNilRule`, `RBGuardClauseRule`, `RBMinMaxRule`, `RBNotEliminationRule`, `RBSuperSendsRule`, and `RBTranslateLiteralInMenusRule`.

The bottom pane shows the source code for the `initialize` method:

```
initialize
  super initialize.
  self rewriteRule
    replace: '@object = nil' with: '@object isNil';
    replace: '@object == nil' with: '@object isNil';
    replace: '@object ~= nil' with: '@object notNil';
    replace: '@object ~ nil' with: '@object notNil'
```

The bottom-right corner of the IDE window displays the page number 197.

Roadmap



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- > **Refactoring**
 - Refactoring Engine and Code Critics
 - **Eclipse refactoring plugins**
- > Aspect-Oriented Programming

A workbench action delegate

When the workbench action proxy is triggered by the user, it delegates to an instance of this class.

```
package astexampleplugin.actions;
...
import org.eclipse.ui.IWorkbenchWindowActionDelegate;

public class ChangeAction implements IWorkbenchWindowActionDelegate {
    ...
    public void run( IAction action ) {
        for ( ICompilationUnit cu : this.classes ) {
            try {
                ...
                parser.setSource( cu );
                ...
                CompilationUnit ast = (CompilationUnit)parser.createAST( null );
                ...
                StackVisitor visitor = new StackVisitor( ast.getAST() );
                ast.accept( visitor );
                ...
            } catch ...
        }
    }
    ...
}
```

http://help.eclipse.org/ganymede/index.jsp?topic=/org.eclipse.jdt.doc.isv/guide/jdt_api_manip.htm

A field renaming visitor

```
package astexampleplugin.ast;
...
import org.eclipse.jdt.core.dom.ASTVisitor;

public class StackVisitor extends ASTVisitor {

    private static final String PREFIX = "_";
    ...
    public boolean visit(FieldDeclaration field){
        ...
    }

    public boolean visit(FieldAccess fieldAccess){
        String oldName = fieldAccess.getName().toString();
        String newName = this.fields.get( oldName );
        if(newName == null){
            newName = PREFIX + oldName;
            this.fields.put( oldName , newName );
        }
        fieldAccess.setName( this.ast.newSimpleName( newName ) );
        return true;
    }
}
```

The visitor simply implements the visit method for field declarations and accesses, and prepends an underscore.

Renaming fields

The screenshot illustrates the 'Rename Fields' refactoring process in an IDE. On the left, a project tree shows the file 'Machine.java' selected. A context menu is open over it, with 'Rename Fields' highlighted. A secondary 'Rename' dialog box is also visible. The main editor window shows the code for the 'Machine' class in two states: the original code and the code after the refactoring.

Original Code:

```
public class Machine {
    private Hashtable<String,Integer> store; // current this.values of variables
    private StringBuffer output; // print stream so far
    private int value; // result of current expression
    private Vector<Integer> vlist; // list of expressions computed

    public Machine() {
        this.store = new Hashtable<String,Integer>();
        this.output = new StringBuffer();
        this.setValue(0);
        this.vlist = new Vector<Integer>();
    }
}
```

Refactored Code:

```
public class Machine {
    private Hashtable<String,Integer> _store; // current this.values of variables
    private StringBuffer _output; // print stream so far
    private int _value; // result of current expression
    private Vector<Integer> _vlist; // list of expressions computed

    public Machine() {
        this._store = new Hashtable<String,Integer>();
        this._output = new StringBuffer();
        this._setValue(0);
        this._vlist = new Vector<Integer>();
    }
}
```

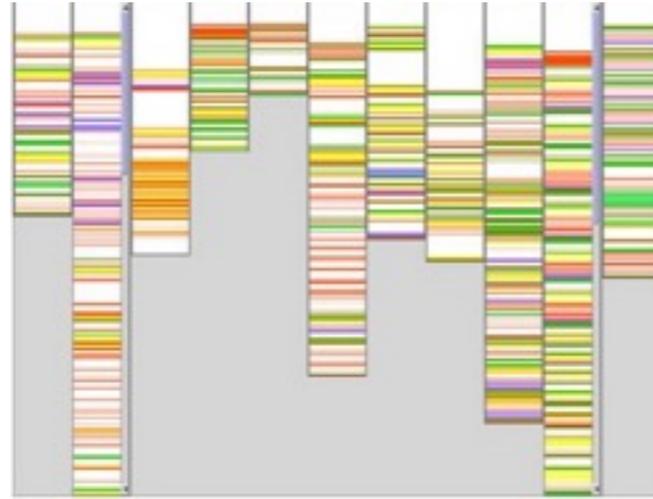
A large blue arrow points from the original code to the refactored code, indicating the transformation. The refactored code uses underscore-prefixed names for all fields.

Roadmap



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- > **Aspect-Oriented Programming**

Problem: cross-cutting concerns



Certain features (like logging, persistence and security), cannot usually be encapsulated as classes. They *cross-cut* code of the system.

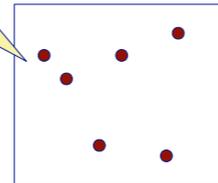
Aspect-Oriented Programming

AOP improves modularity by supporting the separation of cross-cutting concerns.

An *aspect* packages cross-cutting concerns



A *pointcut* specifies a set of *join points* in the target system to be affected



Weaving is the process of applying the aspect to the target system

Canonical example – logging

```
package tjp;

public class Demo {
    static Demo d;
    public static void main(String[] args){
        new Demo().go();
    }
    void go(){
        d = new Demo();
        d.foo(1,d);
        System.out.println(d.bar(new Integer(3)));
    }
    void foo(int i, Object o){
        System.out.println("Demo.foo(" + i + ", " + o + ")\n");
    }
    String bar (Integer j){
        System.out.println("Demo.bar(" + j + ")\n");
        return "Demo.bar(" + j + ")";
    }
}
```

<http://www.eclipse.org/aspectj/downloads.php>



```
Demo.foo(1, tjp.Demo@939b78e)
Demo.bar(3)
Demo.bar(3)
```

A logging aspect

Intercept execution within control flow of Demo.go()

Identify all methods within Demo

```
aspect GetInfo {  
    pointcut goCut(): cflow(this(Demo) && execution(void go()));  
    pointcut demoExecs(): within(Demo) && execution(* *(..));  
    Object around(): demoExecs() && !execution(* go()) && goCut() {  
        ...  
    }  
    ...  
}
```

Wrap all methods except Demo.go()

A logging aspect

```
aspect GetInfo {
  ...
  Object around(): demoExecs() && !execution(* go()) && goCut() {
    println("Intercepted message: " +
      thisJoinPointStaticPart.getSignature().getName());
    println("in class: " +
      thisJoinPointStaticPart.getSignature().getDeclaringClass().getName());
    printParameters(thisJoinPoint);
    println("Running original method:");
    Object result = proceed();
    println(" result: " + result );
    return result;
  }
  ...
}
```



```
Intercepted message: foo
in class: tjp.Demo
Arguments:
  0. i : int = 1
  1. o : java.lang.Object = tjp.Demo@c0b76fa
Running original method:

Demo.foo(1, tjp.Demo@c0b76fa)
  result: null
Intercepted message: bar
in class: tjp.Demo
Arguments:
  0. j : java.lang.Integer = 3
Running original method:

Demo.bar(3)
  result: Demo.bar(3)
Demo.bar(3)
```

Making classes visitable with aspects

```
public class SumVisitor implements Visitor {
    int sum = 0;
    public void visit(Nil l) { }

    public void visit(Cons l) {
        sum = sum + l.head;
        l.tail.accept(this);
    }

    public static void main(String[] args) {
        List l = new Cons(5, new Cons(4,
            new Cons(3, new Nil())));
        SumVisitor sv = new SumVisitor();
        l.accept(sv);
        System.out.println("Sum = " + sv.sum);
    }
}

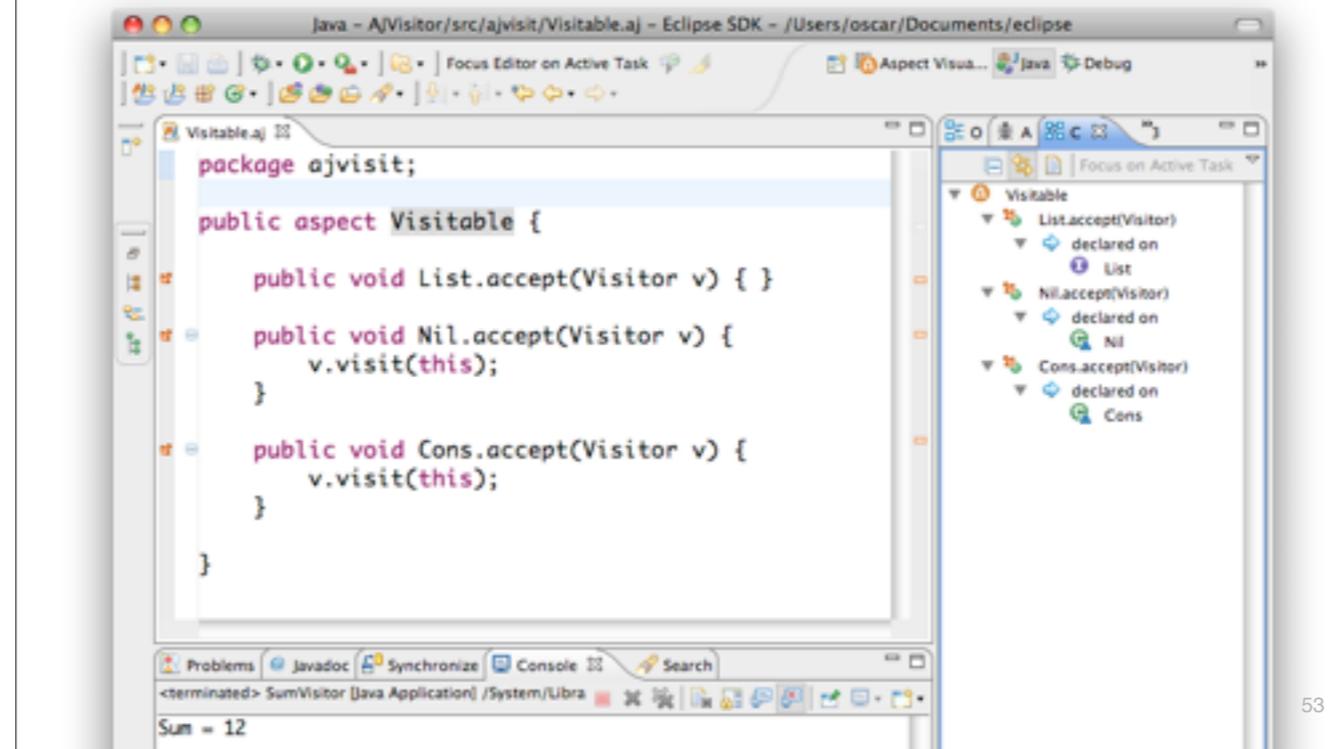
public interface Visitor {
    void visit(Nil l);
    void visit(Cons l);
}
```

We want to write this

```
public interface List {}
public class Nil implements List {}
public class Cons implements List {
    int head;
    List tail;
    Cons(int head, List tail) {
        this.head = head;
        this.tail = tail;
    }
}
```

But we are stuck with this ...

AspectJ



With aspects, who needs visitors?

This would be even cleaner

```
public class SumList {
    public static void main(String[] args) {
        List l = new Cons(5, new Cons(4, new Cons(3, new Nil())));
        System.out.println("Sum = " + l.sum());
    }
}
```

*The missing method
is just an aspect*

```
public aspect Summable {
    public int List.sum() {
        return 0;
    }
    public int Nil.sum() {
        return 0;
    }
    public int Cons.sum() {
        return head + tail.sum();
    }
}
```

Dunno why List.sum() needs a body – it should just be an interface signature.

What you should know!

- ✎ What are typical program transformations?*
- ✎ What is the typical architecture of a PT system?*
- ✎ What is the role of term rewriting in PT systems?*
- ✎ How does TXL differ from Stratego/XT?*
- ✎ How does the Refactoring Engine use metavariables to encode rewrite rules?*
- ✎ Why can't aspects be encapsulated as classes?*
- ✎ What is the difference between a pointcut and a join point?*

Can you answer these questions?

- ✎ How does program transformation differ from metaprogramming?*
- ✎ In what way is optimization a form of PT?*
- ✎ What special care should be taken when pretty-printing a transformed program?*
- ✎ How would you encode typical refactorings like “push method up” using a PT system like TXL?*
- ✎ How could you use a PT system to implement AOP?*



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