

# Parsing F# with PetitParser



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# F# specification



expr  $\leftarrow$  expr expr

expr op expr

ident

[op] const

op  $\leftarrow$  '+'

300 pages

230 production rules

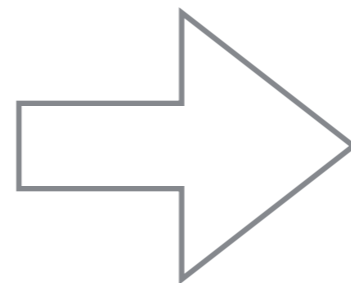
context-free form

# Left recursion

F# spec

```
expr ← expr expr  
      expr op expr  
      ident  
      [op] const
```

op ← '+'



“It’s dead Jim”

**Left recursion in grammars  
cause VM crash!**

# Eliminating left recursion

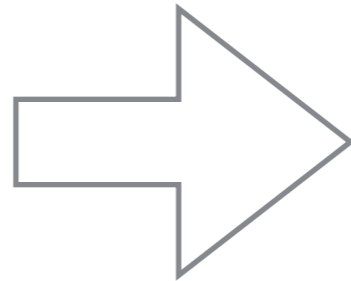
expr  $\leftarrow$  expr expr

expr op expr

ident

[op] const

op  $\leftarrow$  '+'



expr  $\leftarrow$  **exprTerm** expr

**exprTerm** op expr

**exprTerm**

**exprTerm**  $\leftarrow$  ident

[op] const

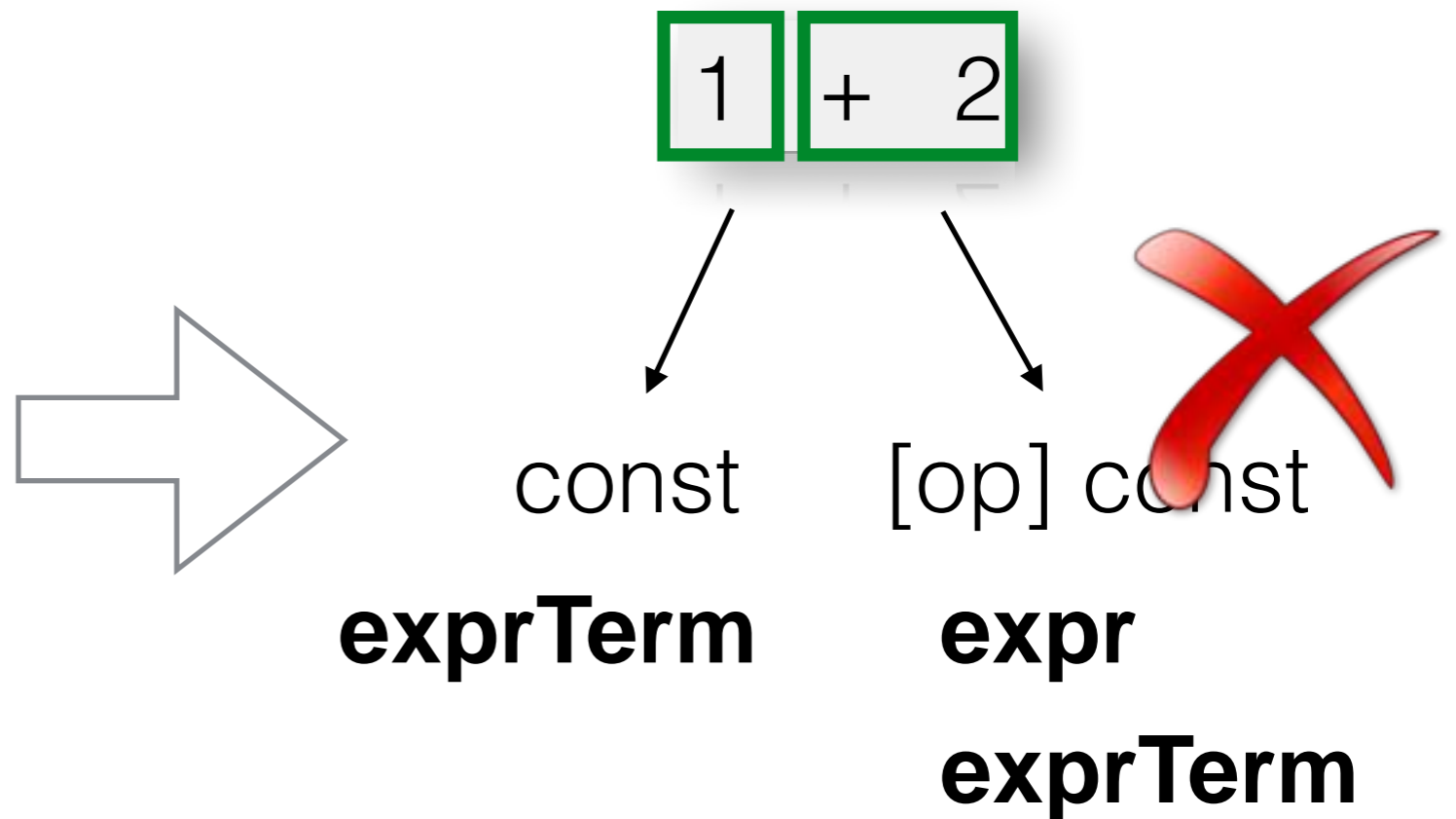
op  $\leftarrow$  '+'

# Ordered choice in PEG

```
expr ← exprTerm expr
      exprTerm op expr
      exprTerm
      [op] const

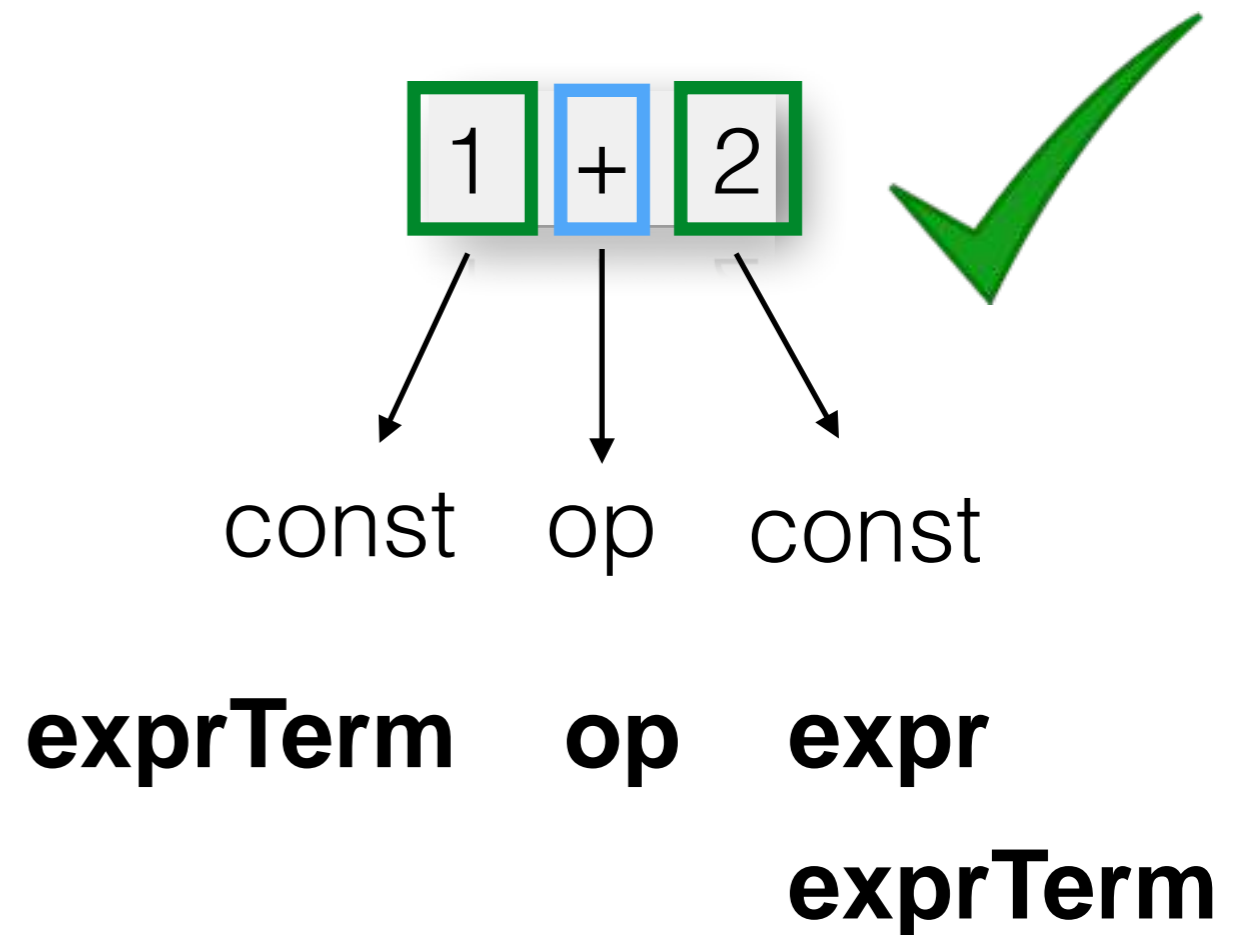
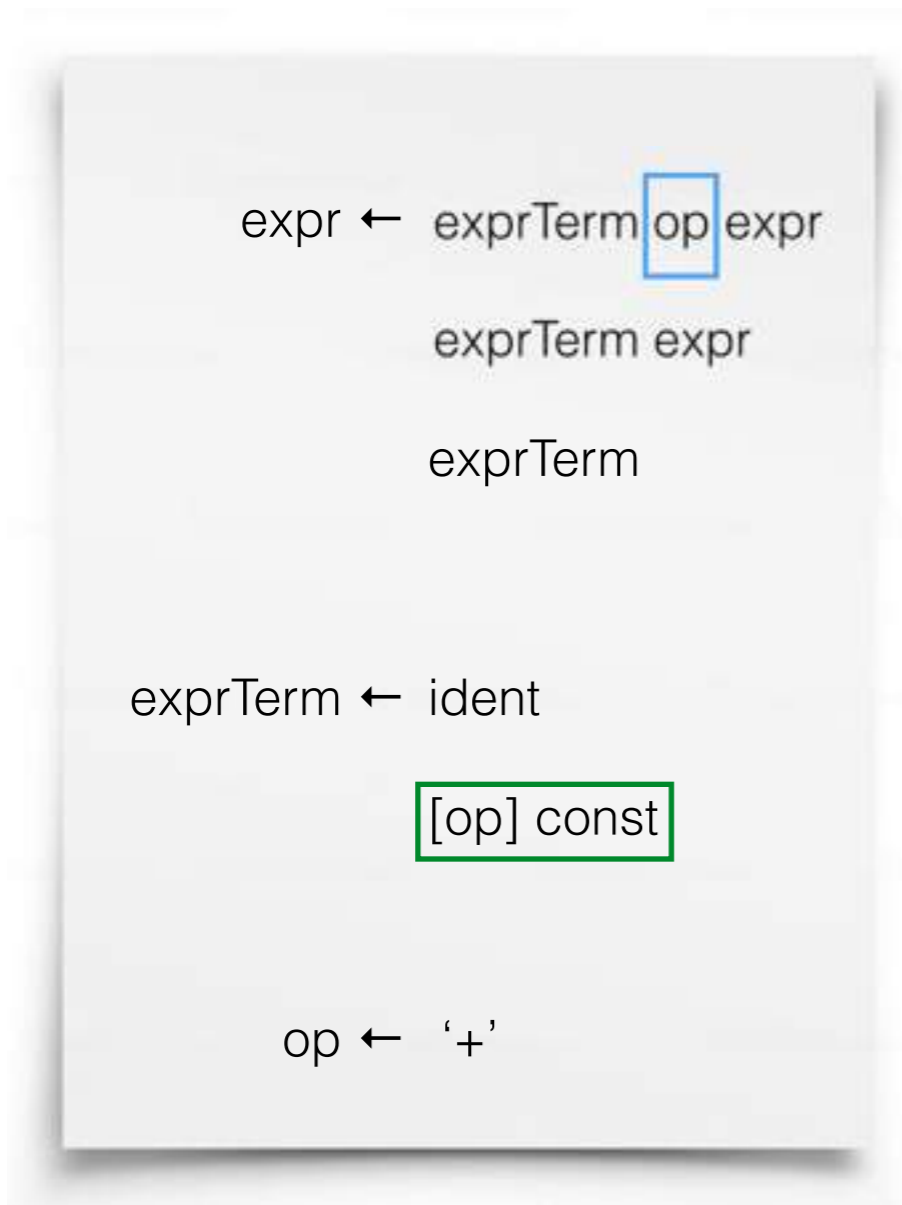
exprTerm ← ident
          [op] const

op ← '+'
```



**If the first alternative succeeds,  
the second alternative is ignored.**

# Ordered choice in PEG



# Performance

expr ← exprTerm ~~op~~ expr

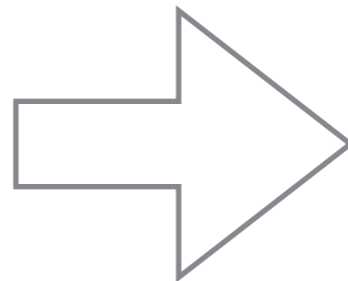
exprTerm ~~expr~~

exprTerm ✓

exprTerm ← ident

[op] const

op ← '+'



1

- **packrat parsing not enabled by default**
- **two possibilities:**
  - refactor grammar
  - enable packrat (parser memoization)

# Grammar transformation

expr  $\leftarrow$  expr expr

expr op expr

ident

[op] const

op  $\leftarrow$  '+'



expr  $\leftarrow$  exprTerm op expr

exprTerm expr

exprTerm

exprTerm  $\leftarrow$  ident

[op] const

op  $\leftarrow$  '+'



# ..in the wild (1)

```
expr :
  const
  ( expr )
  begin expr end
  Long-ident-or-op
  expr '.' Long-ident-or-op
  expr expr
  expr(expr)
  expr<types>
  expr infix-op expr
  prefix-op expr
  expr.[expr]
  expr.[slice-range]
  expr.[slice-range, slice-range]
  expr <- expr
  expr , ... , expr
  new type expr
  { new base-call object-members interface-impls }
  { field-initializers }
  { expr with field-initializers }
  [ expr ; ... ; expr ]
  [| expr ; ... ; expr |]
  expr { comp-or-range-expr }
  [ comp-or-range-expr ]
  [| comp-or-range-expr |]
  lazy expr
  null
  expr : type
  expr :> type
  expr :? type
  expr :?> type
  upcast expr
  downcast expr
  let function-defn in expr
  let value-defn in expr
  let rec function-or-value-defns in expr
  use ident = expr in expr
  . . .
```

```
. . .
  expr '.' Long-ident-or-op
  expr expr
  expr(expr)
  expr<types>
  expr infix-op expr
  prefix-op expr
  expr.[expr]
  expr.[slice-range]
  expr.[slice-range, slice-range]
  expr <- expr
  expr , ... , expr
  new type expr
  { new base-call object-members interface-impls }
  { field-initializers }
  { expr with field-initializers }
  [ expr ; ... ; expr ]
  [| expr ; ... ; expr |]
  expr { comp-or-range-expr }
  [ comp-or-range-expr ]
  [| comp-or-range-expr |]
  lazy expr
  null
  expr : type
  expr :> type
  expr :? type
  expr :?> type
  upcast expr
  downcast expr
```

p. 272, F# 3.0 language specification

# ..in the wild (2)

Operator or expression	Associativity
<code>f&lt;types&gt;</code>	Left
<code>f(x)</code>	Left
<code>.</code>	Left
<code>prefix-op</code>	Left
<code>"  rule"</code>	Right
<code>"f x"</code> <code>"lazy x"</code> <code>"assert x"</code>	Left
<code>**OP</code>	Right
<code>*OP /OP %OP</code>	Left
<code>-OP +OP</code>	Left
<code>:?</code>	Not associative
<code>::</code>	Right
<code>^OP</code>	Right
<code>! =OP &lt;OP &gt;OP =  OP &amp;OP \$</code>	Left
<code>:&gt; :?&gt;</code>	Right

....

Rules of precedence, p. 35, F# 3.0 language specification

# ..in the wild (3)

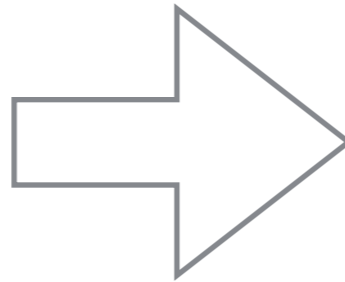
```
....  
additional-constr-defn :  
    attributesopt accessopt new pat as-defn = additional-constr-expr  
!  
additional-constr-expr :  
! stmt ';' additional-constr-expr  
  additional-constr-expr then expr  
  if expr then additional-constr-expr else additional-constr-expr  
! let val-decls in additional-constr-expr  
  additional-constr-init-expr  
  
additional-constr-init-expr :  
! '{' class-inherits-decl field-initializers '  
  new type expr  
....
```

p. 279, F# 3.0 language specification

# Results (1): tool for left recursion detection



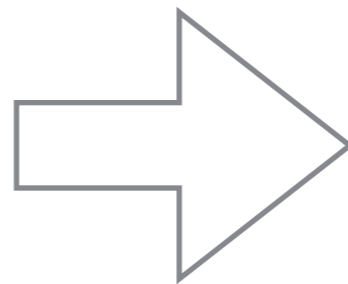
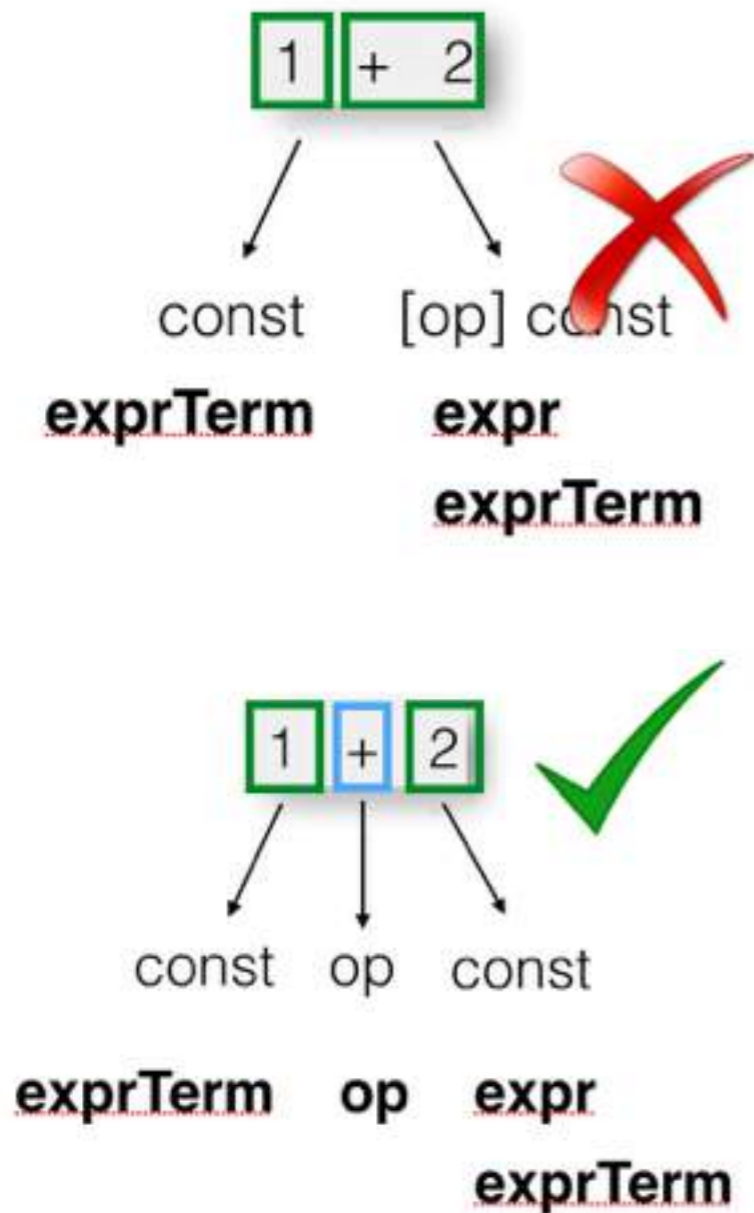
“It’s dead Jim”



```
petitGrammar hasLeftRecursion  
if True: [ petitGrammar getLeftRecursion ].
```

“offline” and “online” algorithm for left recursion detection in PetitParser grammars

# Results (2): AST extractor for validation



F# compiler services  
AST extractor

Generates AST for validation

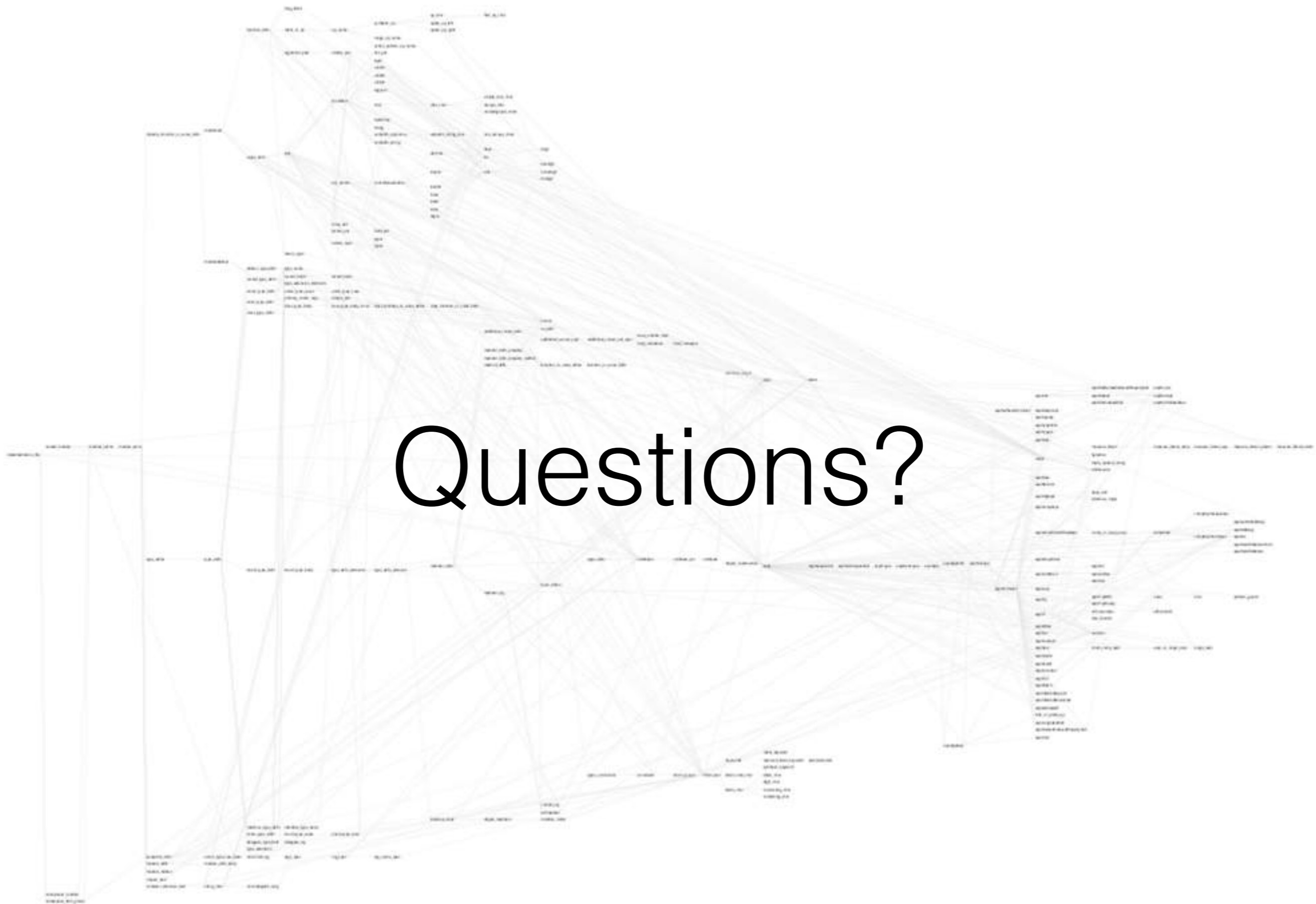
# Results (3): transformed, validated F# syntax rules

- implemented all expression rules and in total 170 out of 230 specified rules
- validated them with 8000 LOC verbose F# code

# Conclusions

The F# CFG to PEG conversion requires quite some manpower for:

- a) left recursion elimination
  - left recursion detection tool
- b) ordered choice / priority modification
  - F# AST extractor for validation
- c) performance issues
- d) slight differences between F# implementation and specification







# F#

- open source, cross-platform compiler
- 3.1 stable release
- 3.0 specification final
- originated from ML and has been influenced by OCaml, C#, Python, Haskell, Scala and Erlang



## verbose syntax

```
1 #light "off"
2
3 let rec fibonacci n =
4     if (n = 1 || n = 2) then
5         1
6     else
7         let result = fibonacci(n-1)
8             + fibonacci(n-2) in
9         result
10 in
11
12 for i in 1 .. 10 do
13     printfn "%d: %d" i (fibonacci i)
14 done
```

## lightweight syntax

```
1 #light "on"
2
3 let rec fibonacci n =
4     if (n = 1 || n = 2) then
5         1
6     else
7         let result = fibonacci(n-1)
8             + fibonacci(n-2)
9         result
10
11 for i in 1 .. 10 do
12     printfn "%d: %d" i (fibonacci i)
```

# eliminating left-recursion

1. For each nonterminal  $A_i$ :
  1. Repeat until an iteration leaves the grammar unchanged:
    1. For each rule  $A_i \rightarrow \alpha_i$ ,  $\alpha_i$  being a sequence of terminals and nonterminals:
      1. If  $\alpha_i$  begins with a nonterminal  $A_j$  and  $j < i$ :
        1. Let  $\beta_i$  be  $\alpha_i$  without its leading  $A_j$ .
        2. Remove the rule  $A_i \rightarrow \alpha_i$ .
        3. For each rule  $A_j \rightarrow \alpha_j$ :
          1. Add the rule  $A_i \rightarrow \alpha_j \beta_i$ .
      2. Remove direct left recursion for  $A_i$  as described above.