

2. Design by Contract

Design by Contract



Bertrand Meyer, *Object-Oriented Software Construction*, Prentice Hall, 1997.



Roadmap

- > Data abstraction and contracts
- > Stacks
- > Design by Contract
- > A Stack ADT
- > Assertions
- > Example: balancing parentheses



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What is Data Abstraction?

- > An Abstract Data Type (ADT):
 - encapsulates data and operations, and
 - hides the implementation behind a well-defined interface.
- > Encapsulation means *bundling together related entities*
- > Information hiding means *exposing an abstract interface and hiding the rest*

In object-oriented languages we can implement ADTs as classes

Why are ADTs important?

Communication – Declarative Programming

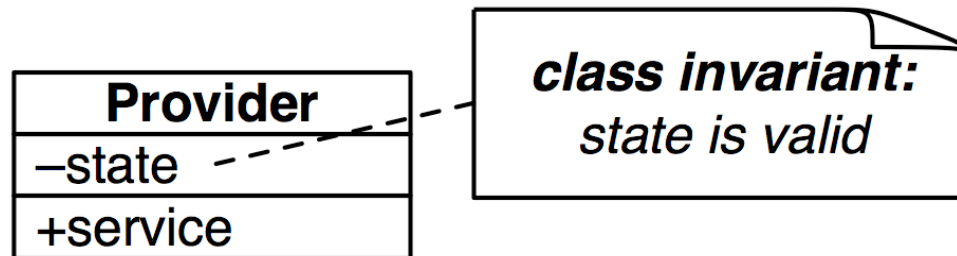
- > An ADT exports *what a client needs to know*, and nothing more!
- > By using ADTs, you communicate *what you want to do*, not how to do it!
- > ADTs allow you to *directly model your problem domain* rather than how you will use to the computer to do so.

Software Quality and Evolution

- > ADTs help you to *decompose a system into manageable parts*, each of which can be separately implemented and validated.
- > ADTs *protect clients from changes* in implementation.
- > ADTs encapsulate client/server *contracts*
- > Interfaces to ADTs can be *extended* without affecting clients.
- > New implementations of ADTs can be *transparently added* to a system.

Class Invariants

An invariant is a predicate that *must hold* at certain points in the execution of a program

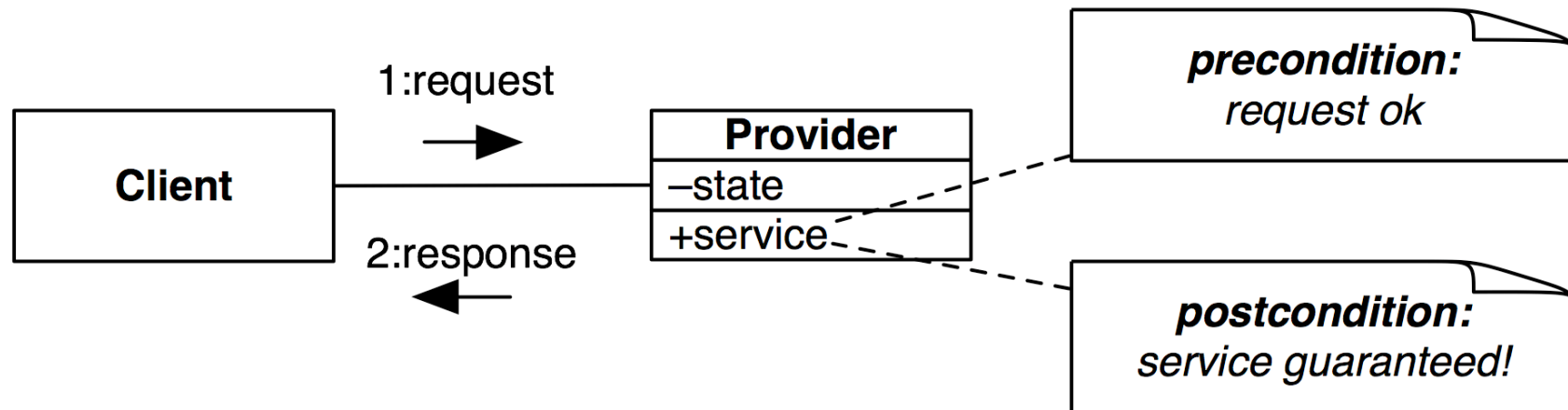


A class invariant characterizes the *valid states of instances*
It must hold:

1. *after construction*
2. *before and after every public method*

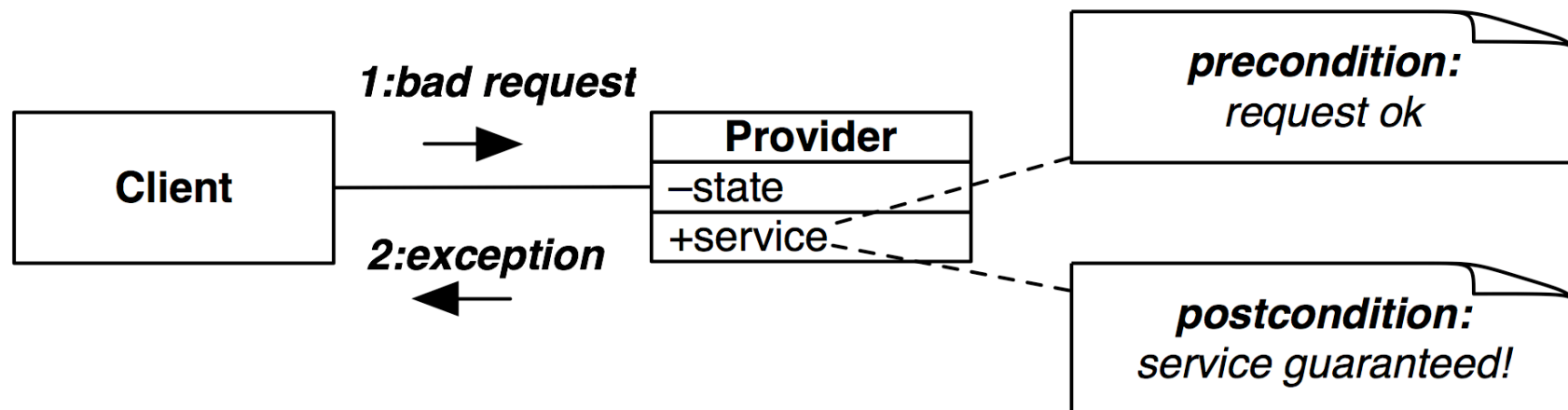
Contracts

A contract *binds the client* to pose valid requests, and *binds the provider* to correctly provide the service.



Contract violations

If either the client or the provider violates the contract, an *exception* is raised.



NB: The service does not need to implement any special logic to handle errors — it simply raises an exception!

Exceptions, failures and defects

- > An exception is the occurrence of an *abnormal condition during the execution of a software element*.
- > A failure is the *inability of a software element to satisfy its purpose*.
- > A defect (AKA “bug”) is the *presence in the software of some element not satisfying its specification*.

Disciplined Exceptions

- > There are only two reasonable ways to react to an exception:
 1. clean up the environment and *report failure* to the client (“organized panic”)
 2. attempt to *change the conditions* that led to failure and *retry*

A failed assertion often indicates presence of a software defect, so “organized panic” is usually the best policy.

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Stacks

A Stack is a classical data abstraction with many applications in computer programming.

Stacks support two mutating methods: push and pop.

Operation	Stack	isEmpty()	size()	top()
		TRUE	0	(error)
push(6)		FALSE	1	6
push(7)		FALSE	2	7
push(3)		FALSE	3	3
pop()		FALSE	2	7
push(2)		FALSE	3	2
pop()		FALSE	2	7

Stack pre- and postconditions

Stacks should respect the following contract:

<i>service</i>	<i>pre</i>	<i>post</i>
<code>isEmpty()</code>	-	<i>no state change</i>
<code>size()</code>	-	<i>no state change</i>
<code>push(Object item)</code>	-	not empty, size == old size + 1, top == item
<code>top()</code>	not empty	<i>no state change</i>
<code>pop()</code>	not empty	size == old size - 1

Stack invariant

- > The only thing we can say about the Stack class invariant is that the size is always ≥ 0
 - we don't know anything yet about its state!

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Design by Contract

When you design a class, each service S provided must specify a clear contract.

“If you promise to call S with the precondition satisfied, then I , in return, promise to deliver a final state in which the post-condition is satisfied.”

Consequence:

—if the precondition does not hold, *the object is not required to provide anything!* (in practice, an exception is raised)

In other words ...

Design by Contract = *Don't accept anybody else's garbage!*

Pre- and Post-conditions

The pre-condition binds clients:

- it defines what the ADT *requires* for a call to the operation to be legitimate
- it may involve *initial state and arguments*
- example: *stack is not empty*

The post-condition, in return, binds the provider:

- it defines the conditions that the ADT *ensures* on return
- it may only involve the *initial and final states, the arguments and the result*
- example: *size = old size + 1*

Benefits and Obligations

A contract provides *benefits and obligations* for both clients and providers:

	<i>Obligations</i>	<i>Benefits</i>
<i>Client</i>	Only call pop() on a non-empty stack!	Stack size decreases by 1. Top element is removed.
<i>Provider</i>	Decrement the size. Remove the top element.	No need to handle case when stack is empty!

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StackInterface

Interfaces let us *abstract* from concrete implementations:

```
public interface StackInterface {  
    public boolean isEmpty();  
    public int size();  
    public void push(Object item);  
    public Object top() ;  
    public void pop();  
}
```

-  How can clients accept multiple implementations of an ADT?
- ✓ *Make them depend only on an interface or an abstract class.*

Interfaces in Java

Interfaces *reduce coupling* between objects and their clients:

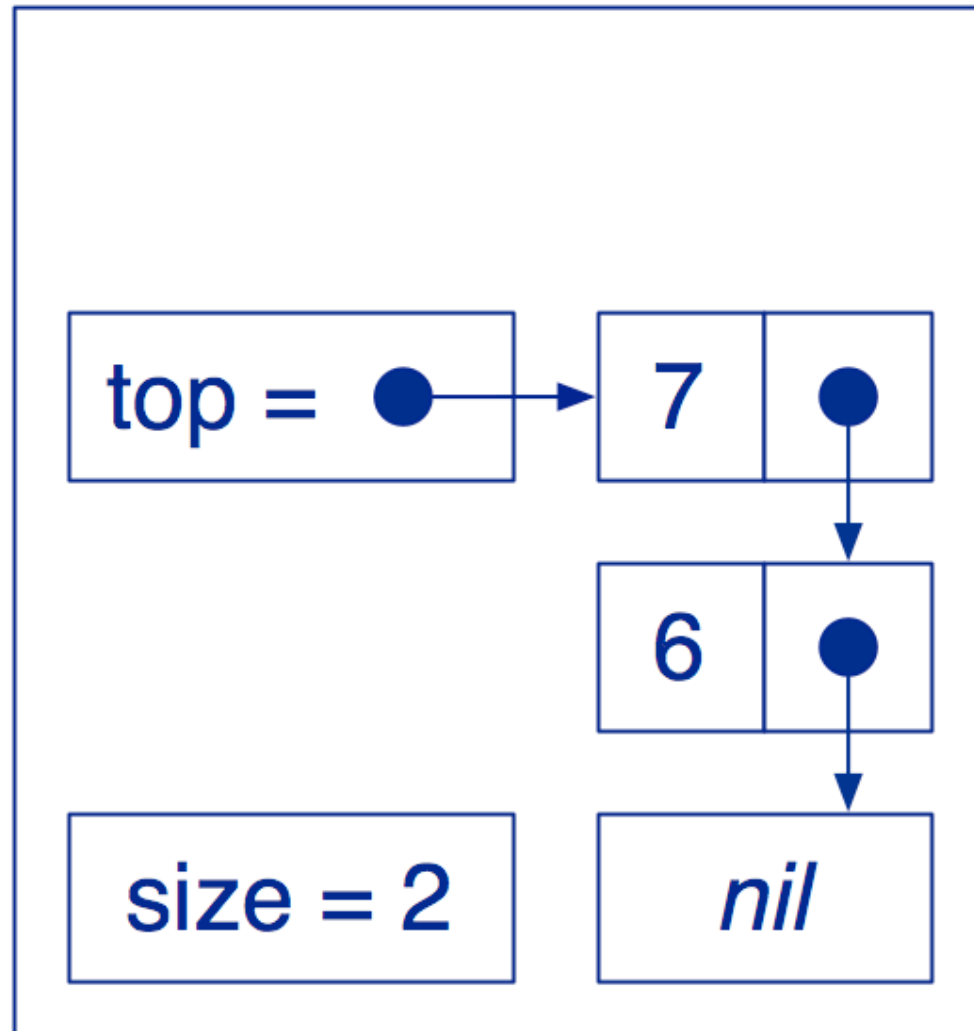
- > A class can implement multiple interfaces
 - ... but can only extend one parent class
- > Clients should depend on an interface, not an implementation
 - ... so implementations don't need to extend a specific class

Define an interface for any ADT that will have more than one implementation

Stacks as Linked Lists

A Stack can easily be implemented by a linked data structure:

```
stack = new Stack();  
stack.push(6);  
stack.push(7);  
stack.push(3);  
stack.pop();
```




LinkStack Cells

We can define the Cells of the linked list as an *inner class* within LinkStack:

```
public class LinkStack implements StackInterface {
    private Cell top;
    private class Cell {
        Object item;
        Cell next;
        Cell(Object item, Cell next) {
            this.item = item;
            this.next = next;
        }
    }
    ...
}
```

Private vs Public instance variables

-  When should instance variables be public?
- ✓ *Always make instance variables private or protected.*

The Cell class is a special case, since its instances are strictly private to LinkStack!

LinkStack ADT

The constructor must construct a *valid initial state*:

```
public class LinkStack implements StackInterface {
    ...
    private int size;
    public LinkStack() {
        // Establishes the class invariant.
        top = null;
        size = 0;
    }
    ...
}
```

Class Invariants

A class invariant is any condition that expresses the *valid states* for objects of that class:

- > it must be *established* by every constructor
- > every public method
 - may *assume* it holds when the method starts
 - must *re-establish* it when it finishes

Stack instances must satisfy the following invariant:

- > $\text{size} \geq 0$
- > ...

LinkStack Class Invariant

A valid `LinkStack` instance has an integer `size`, and a `top` that points to a sequence of linked `Cells`, such that:

- `size` is always ≥ 0
- When `size` is zero, `top` points nowhere (`== null`)
- When `size` > 0 , `top` points to a `Cell` containing the top item

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Assertions

- > An assertion is a declaration of a *boolean expression* that the programmer believes *must hold* at some point in a program.
 - Assertions should not affect the logic of the program
 - If an assertion fails, an *exception* is raised

```
x = y*y;  
assert x >= 0;
```

Assertions

Assertions have four principle applications:

1. **Help in writing correct software**
 - formalizing invariants, and pre- and post-conditions
2. **Documentation aid**
 - specifying contracts
3. **Debugging tool**
 - testing assertions at run-time
4. **Support for software fault tolerance**
 - detecting and handling failures at run-time

Assertions in Java

`assert` is a keyword in Java since version 1.4

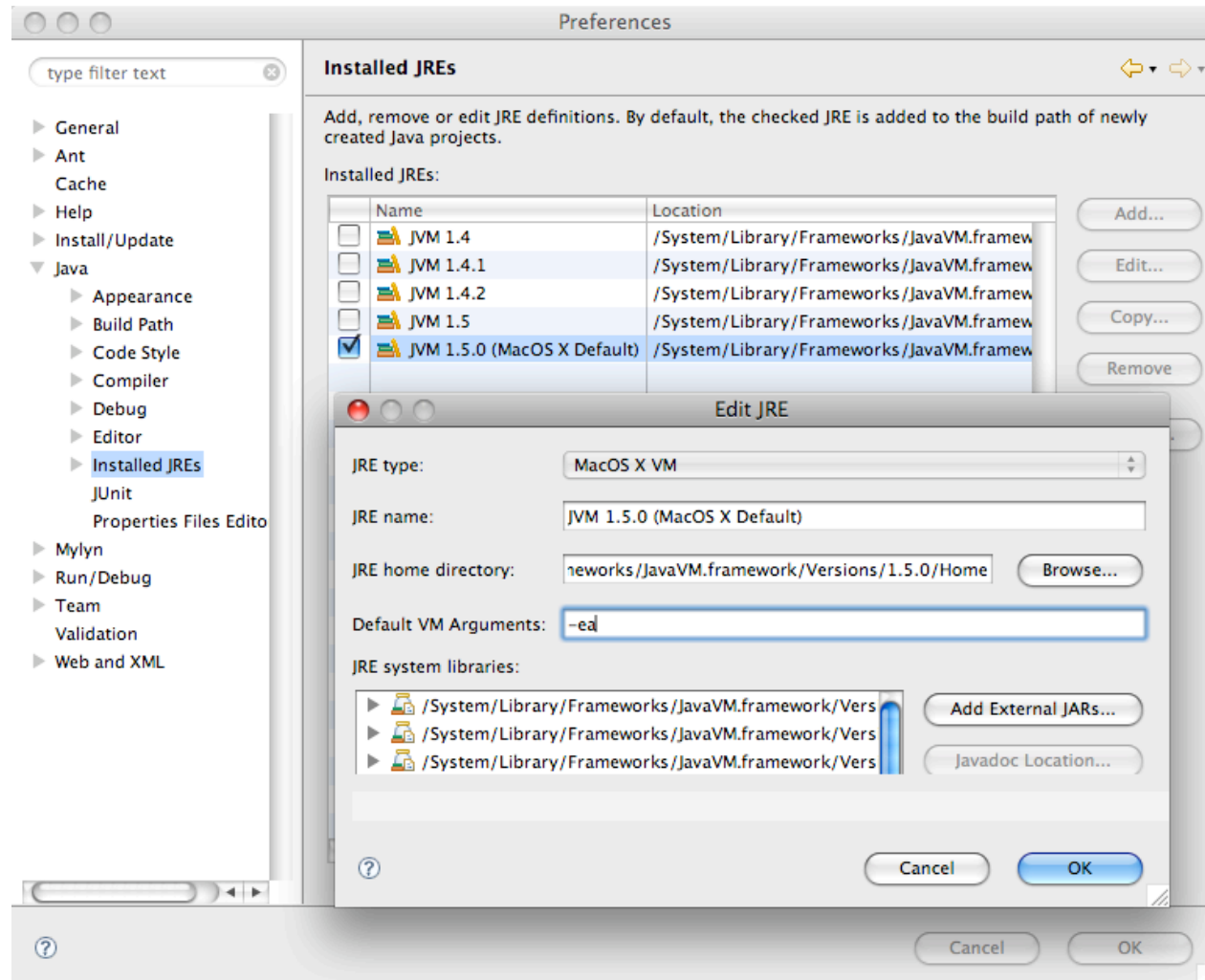
```
assert expression;
```

will raise an `AssertionError` if *expression* is false.

— *NB*: Throwable *Exceptions* must be declared; *Errors* need not be!

✓ *Be sure to enable exceptions in eclipse! (And set the vm flag -enableassertions [-ea])*

Enabling assertions in eclipse




Checking pre-conditions

Assert pre-conditions to inform clients when *they* violate the contract.

```
public Object top() {  
    assert(!this.isEmpty()); // pre-condition  
    return top.item;  
}
```

NB: This is all you have to do!

-  When should you check pre-conditions to methods?
- ✓ *Always check pre-conditions, raising exceptions if they fail.*

Checking class invariants

Every class has its own invariant:

```
protected boolean invariant() {  
    return (size >= 0) &&  
        ( (size == 0 && this.top == null)  
          || (size > 0 && this.top != null));  
}
```

Why protected and not private?

Checking post-conditions

Assert post-conditions and invariants to inform yourself when *you* violate the contract.

```
public void push(Object item) {
    top = new Cell(item, top);
    size++;
    assert !this.isEmpty();           // post-condition
    assert this.top() == item;       // post-condition
    assert invariant();
}
```

NB: This is all you have to do!

 When should you check post-conditions?

✓ *Check them whenever the implementation is non-trivial.*

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Example: Balancing Parentheses

Problem:

- > Determine whether an expression containing parentheses (), brackets [] and braces { } is correctly balanced.

Examples:

- > balanced:

```
if (a.b()) { c[d].e(); }  
else { f[g][h].i(); }
```

- > not balanced:

```
((a+b()))
```

A simple algorithm

Approach:

- > when you read a *left* parenthesis, *push the matching parenthesis* on a stack
- > when you read a *right* parenthesis, *compare it* to the value on top of the stack
 - if they *match*, you *pop and continue*
 - if they *mismatch*, the expression is *not balanced*
- > if the *stack is empty* at the end, the whole expression is *balanced*, otherwise not

Using a Stack to match parentheses

Sample input: “([{ }]]”

<i>Input</i>	<i>Case</i>	<i>Op</i>	<i>Stack</i>
(left	push))
[left	push])]
{	left	push })]]
}	match	pop)]
]	match	pop)
]	mismatch	^false)

The ParenMatch class

A ParenMatch object *uses a stack* to check if parentheses in a text String are balanced:

```
public class ParenMatch {
    private String line;
    private StackInterface stack;

    public ParenMatch (String aLine, StackInterface aStack)
    {
        line = aLine;
        stack = aStack;
    }
}
```

A declarative algorithm

We implement our algorithm at a high level of abstraction:

```
public boolean parenMatch() {
    for (int i=0; i<line.length(); i++) {
        char c = line.charAt(i);
        if (isLeftParen(c)) { // expect matching right paren later
            stack.push(matchingRightParen(c)); // Autoboxed to Character
        } else {
            if (isRightParen(c)) {
                // empty stack => missing left paren
                if (stack.isEmpty()) { return false; }
                if (stack.top().equals(c)) { // Autoboxed
                    stack.pop();
                } else { return false; } // mismatched paren
            }
        }
    }
    return stack.isEmpty(); // not empty => missing right paren
}
```

Ugly, procedural version

```
public boolean parenMatch() {
    char[] chars = new char[1000]; // ugly magic number
    int pos = 0;
    for (int i=0; i<line.length(); i++) {
        char c = line.charAt(i);
        switch (c) { // what is going on here?
            case '{' : chars[pos++] = '}'; break;
            case '(' : chars[pos++] = ')'; break;
            case '[' : chars[pos++] = ']'; break;
            case ']' : case ')' : case '}' :
                if (pos == 0) { return false; }
                if (chars[pos-1] == c) { pos--; }
                else { return false; }
                break;
            default : break;
        }
    }
    return pos == 0; // what is this?
}
```

Helper methods

The helper methods are trivial to implement, and their details only get in the way of the main algorithm.

```
private boolean isLeftParen(char c) {  
    return (c == '(') || (c == '[') || (c == '{');  
}  
  
private boolean isRightParen(char c) {  
    return (c == ')') || (c == ']') || (c == '}');  
}
```

Running parenMatch

```
public static void parenTestLoop(StackInterface stack) {
    BufferedReader in =
        new BufferedReader(new InputStreamReader(System.in));
    String line;
    try {
        System.out.println("Please enter parenthesized expressions to test");
        System.out.println("(empty line to stop)");
        do {
            line = in.readLine();
            System.out.println(new ParenMatch(line, stack).reportMatch());
        } while(line != null && line.length() > 0);
        System.out.println("bye!");
    } catch (IOException err) {
    } catch (AssertionException err) {
        err.printStackTrace();
    }
}
```








Running ParenMatch.main ...

```
Please enter parenthesized expressions to test
(empty line to stop)
(hello) (world)
"(hello) (world)" is balanced
()
"()" is balanced
static public void main(String args[]) {
"static public void main(String args[]) {" is not balanced
()
"()" is not balanced
}
"}" is balanced






"" is balanced
bye!
```

Which contract has been violated?

What you should know!

-  *What is an abstract data type?*
-  *What is the difference between encapsulation and information hiding?*
-  *How are contracts formalized by pre- and post-conditions?*
-  *What is a class invariant and how can it be specified?*
-  *What are assertions useful for?*
-  *What situations may cause an exception to be raised?*
-  *How can helper methods make an implementation more declarative?*

Can you answer these questions?

-  *When should you call `super()` in a constructor?*
-  *When should you use an inner class?*
-  *What happens when you `pop()` an empty `java.util.Stack`?
*Is this good or bad?**
-  *What impact do assertions have on performance?*
-  *Can you implement the missing `LinkStack` methods?*

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